FOURIER BASED FAST MULTipoLE METHOD FOR THE HELMHOLTZ EQUATION

C. Cecka* and E. Darve
Stanford University and Army High Performance Computing and Research Center
Stanford, CA 94305

ABSTRACT

The multilevel fast multipole method (MLFMM) is an algorithm that has had great success in reducing the computational time required to find the solution to the Galerkin boundary integral form of the Helmholtz equation. We present a new formulation of the MLFMM using Fourier basis functions rather than spherical harmonics in order to accelerate and simplify the time-critical stages of the algorithm. With modifications to the transfer function in the precomputation stage of the MLFMM, the interpolation and anterpolation algorithms become straightforward applications of FFT interpolations only. Using spectral methods, constructive algorithms are derived to determine a near-optimal quadrature for a given level in the algorithm and an a-priori estimate of the integration error.

1. INTRODUCTION

Since its development in (Rokhlin, 1990; Coifman et al., 1993; Rokhlin, 1993; Engheta et al., 1992; Rahola, 1996), the MLFMM has proven to be a very effective tool for solving acoustic and electromagnetic problems. We consider only the scalar wave equation, although any results can be extended to the vector case as described in (Chew et al., 2001; Darve, 2000b). Applying the boundary integral method to the Helmholtz equation results in a dense linear system which can be solved by iterative methods such as GMRES or BCGSTAB. These methods require a dense matrix-vector product which, for a direct implementation, requires $O(N^2)$ operations to compute. The MLFMM uses an approximation of the dense matrix to perform the product in $O(N \log N)$ time. Computations taking advantage of this approximation are distributed through an oct-tree encompassing the domain of the scatterers to achieve the improved asymptotic running time.

There are a number of difficulties in implementing the MLFMM, each of which must be carefully considered and optimized to achieve the improved complexity. The largest complication is the quadrature sampling rate must increase with the size of the box in the oct-tree, requiring interpolation and anterpolation algorithms over the sphere to transform the data between quadratures. Local algorithms and interpolant matrix sparsifications are fast, but incur large errors. Fast transforms are available for spherical harmonics, but they are approximate, quite complicated to implement, and not always stable. We present the MLFMM and these choices in the next section.

2. THE MULTILEVEL FAST MULTipoLE METHOD

The MLFMM improves the asymptotic complexity of the matrix-vector multiplication

$$\sigma_i = \sum_{j\neq i} e^{i|\delta_x - \delta_j|} \psi_j := \sum_j M_{ij} \psi_j$$ (1)

for $i, j = 1, \ldots, N$ from $O(N^2)$ to $O(N \log^2 N)$. This improvement is based on the Gegenbauer series

$$\frac{e^{i|\delta|}}{|\delta|} = \frac{ik}{\pi} \sum_{n=0}^{\infty} (-1)^n (2n + 1) G_n(\delta)$$ (2)

$$G_n(\delta) = h^{(1)}_n(\kappa |\delta|) j_n(\kappa |\delta|) P_n(\hat{\delta} \cdot \hat{\delta}_0)$$

where $h^{(1)}_n$ is the spherical Hankel function of the first kind, $j_n$ is the spherical Bessel function of the first kind, and $P_n$ is the Legendre polynomial of order $n$. The series converges absolutely and uniformly for $|\delta| < |\delta_0|$ and has been studied extensively in (Carayol & Collino, 2004; Darve, 2000a).

Truncating the Gegenbauer series at $\ell$ and using the identity

$$\int_{S^2} e^{i\kappa \cdot \delta} P_n(\delta_0 \cdot \delta) \, dS(\delta) = 4\pi \rho^\ell j_n(\kappa |\delta|) P_n(\hat{\delta} \cdot \hat{\delta}_0)$$

where the integral is over the unit sphere $S^2 = \{\delta \in \mathbb{R}^3 : |\delta| = 1\}$, then

$$\frac{e^{i|\delta|}}{|\delta|} = \int_{S^2} e^{i\kappa \cdot \delta} T_{\ell, n}(\delta) \, dS(\delta) + \varepsilon_{\ell}$$ (3)
**Fourier Based Fast Multipole Method For The Helmholtz Equation**

**PERFORMING ORGANIZATION NAME(S) AND ADDRESS(ES)**
Stanford University and Army High Performance Computing and Research Center
Stanford, CA 94305

**ABSTRACT**
where we have defined the Gegenbauer truncation error $\varepsilon_{c}$ and the “transfer function”

$$T_{l,r_{0}}(s) = \frac{iK}{4\pi} \sum_{n=0}^{\ell} e^{i(2n+1)h_{n}^{(1)}(\kappa \cdot r_{0})} P_{n}(s \cdot \hat{r}_{0}).$$

(4)

Additionally, it will be useful to note the Jacobi-Anger series

$$e^{ic r \cdot s} = \sum_{n=0}^{\infty} i^{n} (2n+1) j_{n}(\kappa \cdot r) P_{n}(s \cdot \hat{r}).$$

(5)

Consider two disjoint clusters of points $x_{i} \mid i \in A$ and $x_{i} \mid i \in B$ with radii $r_{A} \geq r_{B} \geq 0$ and centers $c_{A}$ and $c_{B}$ respectively. If $|c_{A} - c_{B}| > r_{A}$, then the matrix-vector product (1) is accelerated by using the approximation

$$M \approx \begin{bmatrix} M_{AA} & M_{AB} \\ M_{BA} & M_{BB} \end{bmatrix}$$

where

$$M_{AA} = [M_{ij}]_{i,j \in A}$$

$$M_{AB} = \left[ \int_{S^{2}} e^{i\kappa \cdot r_{i,c}} T_{l,r_{0}}(s) e^{i\kappa \cdot r_{j}} dS(s) \right]_{i \in A, j \in B}$$

$$r_{ic} = x_{i} - c_{A}, \quad r_{0} = c_{A} - c_{B}, \quad r_{cj} = c_{B} - x_{j}$$

The improved asymptotic complexity of the FMM is achieved by constructing a tree of nodes $\{c_{p}\}$, typically an octree, which partitions the support of the source and field points at the finest level. First, the outgoing plane-wave expansions due to the source points (or basis functions) are accumulated at the finest level in the tree. These outgoing expansions are then aggregated upward through the tree via additional plane-wave expansions. Incoming plane-wave expansions are computed from the outgoing by multiplication of the transfer function. These incoming plane-waves are then disaggregated downward through the tree. At the finest level the integration is performed and accumulated with the near field contribution to determine the field value at the field points (or testing functions).

### 2.1 Spherical Quadrature in the MLFMM

Noting the addition formula

$$(2n+1)P_{n}(\hat{p} \cdot \hat{q}) = 4\pi \sum_{m=-n}^{n} Y_{n}^{m}(\hat{p}) Y_{n}^{m}(\hat{q})$$

the transfer function (4) and Jacobi-Anger series can be expressed as

$$T_{l,r_{0}}(s) = \sum_{n=0}^{\ell} \sum_{m=-n}^{n} l_{n}^{m} Y_{n}^{m}(s)$$

(6)

$$e^{i\kappa r \cdot s} = \sum_{n=0}^{\infty} \sum_{m=-n}^{n} e^{i\kappa r \cdot s}$$

(7)

Therefore, we wish to exactly integrate the spherical harmonics, $Y_{n}^{m}$. Below, we enumerate a number of choices for appropriate spherical quadratures that have previously been studied.

1. The simplest choice are sample points chosen uniformly in $\theta$ and $\phi$. However, this choice does not accurately integrate the spherical harmonics and requires approximately twice as many points as the Gauss-Legendre quadrature below (Darve, 2000b).

2. The most common choice of sample points are uniform points for $\theta$ and Gauss-Legendre points for $\phi$. With $N+1$ uniform points in the $\theta$ direction and $N+1$ Gauss-Legendre points in the $\phi$ direction, all $Y_{n}^{m}$, $-n \leq m \leq n$, $0 \leq n \leq N$ are integrated exactly (Darve, 2000b; Koc et al., 1999).

3. McLaren in (McLaren, 1963) developed optimal choices of samples for general functions on $S^{2}$ based on equally spaced points and derived from invariants of finite groups of rotations. He also proposes a method for constructing equally weighted integration formulas on sets of any desired number of points by taking the union of icosahedral configurations.

### 2.2 Interpolation in the MLFMM

The quadrature sampling rate depends on the spectral content of the translation operator, $e^{i\kappa \cdot r}$. This term’s spherical harmonic coefficients $e_{n}^{m}$ in (7) decreases super-exponentially roughly after the $\kappa \cdot r$ mode, which scales linearly with the box size in the tree. Therefore, as we go up the tree in the aggregation step and $|r|$ becomes larger, a larger quadrature and corresponding interpolation algorithm are required to resolve these higher modes. These modes must be resolved since they interact with the modes in the transfer function, which do not decay as $\ell$ increases.

Similarly, as we go down the tree in the disaggregation step, $|r|$ becomes smaller, preventing the higher modes of the incoming field approximation from contributing to the integral as a consequence of Parseval’s theorem. Thus, as the incoming field is disaggregated down the tree, a smaller quadrature can be used to resolve it. This makes the integration faster and is actually required to achieve an improved asymptotic running time at the cost of requiring anterpolation algorithms. Below, we enumerate a number of choices for interpolation and anterpolation algorithms that have previously been studied.

1. General interpolation algorithms for non-uniform grid points like Lagrange interpolation or B-
splines are fast and provide for simple error analysis. In (Koc et al., 1999) it is shown that the error induced from Lagrange interpolation decreases exponentially as the number of interpolation points is increased for a given function of finite bandwidth. Thus, there is a trade-off between error and speed. The introduction of error is unavoidable with local schemes and is dependent on the values of the function at the time of interpolation.

2. For a set of quadrature points \((\theta_k, \phi_k), k = 1, \ldots, K\) with respective weights \(\omega_k\) and corresponding function value \(f_k\), a direct spherical harmonic transform sends \(f_k\) to a new quadrature \((\theta'_k, \phi'_k), k' = 1, \ldots, K'\) via the linear transformation

\[
f'_k = \sum_{m,j \leq K} Y^m_j(\theta'_k, \phi'_k) \sum_k \omega_k Y^m_j(\theta_k, \phi_k) f_k
\]

This transform has nice properties analogous to those of the Fourier transform. A direct computation requires \(O(KK')\) operations which would result in an \(O(N^2)\) FMM. Fast spherical transforms (FST) have been developed in (Driscoll & Healy, 1994; Healy et al., 2003; Rokhlin & Tygert, 2006) and applied to the FMM in (Chowdhury & Jandhyala, 2006). Using the FST reduces the interpolation and anterpolation procedures to \(O(K \log^2 K)\), which results in a \(O(N \log^2 N)\) MLFMM. However, the accuracy and stability of these algorithms are questionable as they rely on the nonequispaced discrete Fourier transform and/or the fast Legendre transform.

3. Approximations of the direct spherical transform have also been investigated in (Jakob-Chien & Alpert, 1997; Darve, 2000b). The interpolation matrix

\[
A'_{k'k} = \sum_{m,j \leq K} Y^m_j(\theta'_k, \phi'_k) \sum_k \omega_k Y^m_j(\theta_k, \phi_k)
\]

can be sparsified in a number of ways to provide an interpolation/interpolation method that scales as \(O(K)\) with scalable relative error.

3. FOURIER BASED MLFMM

The Fourier based fast multipole method is based on the identity

\[
\int_{S^2} e^{i m \phi} T_{\ell,r_0}(s) \, dS(s) = \int_0^{2\pi} \int_0^{2\pi} e^{i m \phi} \mathcal{T}_{\ell,r_0}^m(s) \, d\theta d\phi
\]

where \(s = [\cos(\theta) \sin(\phi), \sin(\theta) \sin(\phi), \cos(\phi)]\) and \(T_{\ell,r_0}^m\) is the “modified transfer function”

\[
T_{\ell,r_0}^m(\theta, \phi) = \frac{1}{2} T_{\ell,r_0}(s) |\sin(\phi)|
\]

Because the integrand is continuous and periodic, this formulation allows the use of the Fourier basis functions \(\{e^{im\phi}\}\) which form an \(L^2\) orthonormal set over \([0, 2\pi] \times [0, 2\pi]\). Thus, we efficiently use two dimensional uniform quadratures, fast Fourier transforms in the interpolation and anterpolation steps, and spectral arguments in the error analysis. The result is a fast and easily implemented version of the MLFMM with strong control over the maximum absolute error, which we detail in the following sections.

3.1 Computing the Modified Transfer Function

Select a uniform quadrature with points \((\theta_i, \phi_j)\) defined by

\[
\theta_i = 2\pi \frac{i}{N_\theta}, \quad \phi_j = 2\pi \frac{j}{N_\phi}
\]

Noting that the plane wave and modified transfer function both have spherical symmetry since

\[
s(\theta, \phi) = s(\pi + \theta, 2\pi - \phi)
\]

then the computational and memory cost will be reduced if the quadrature is also constructed with spherical symmetry since only half of the quadrature points will be need computed and stored. Thus, we force \(N_\phi\) to be a multiple of 2 and \(N_\theta\) to be a multiple of 4.

By taking advantage of all of the symmetries of this quadrature the number of modified transfer functions that need to be precomputed is reduced from 316 per level to 34 saving a factor of 9.3 in memory and costing a negligible permutation of values in time.

A key step to computing \(T_{\ell,r_0}(s)\) is to remove the frequencies which a given quadrature cannot resolve. If \(T_{\ell,r_0}(s)\) were simply sampled, significant frequency aliasing (or folding) would occur unless we used an unreasonably large quadrature. This is due to the slow decay of the Fourier series of \(|\sin(\phi)|\),

\[
\mathcal{F}(m; |\sin(\phi)|) = \frac{(-1)^n + 1}{\pi(1 - m^2)} = \begin{cases} \frac{2}{\pi} \frac{1}{1 - m^2} & \text{if } m \text{ even} \\ 0 & \text{if } m \text{ odd} \end{cases}
\]

Since the spectrum of the plane-wave function

\[
e^{i m \phi} = e^{i \ell |r| \cos(\theta, r)} \sum_{n=-\infty}^{\infty} \ell^n J_n(\kappa |r|) e^{in(\theta, r)}
\]

decays very rapidly for \(n \geq \kappa |r|\), the high frequencies in \(T_{\ell,r_0}(s)\) will not contribute to the final integral as a result of Parseval’s theorem. By removing these frequencies from the modified transfer function, a smaller quadrature can be used without significantly affecting the total error.
Suppose we have chosen a quadrature characterized by \((N_\phi, N_\theta)\). Since \(T_{\ell,r_0}\) is band-limited with bandwidth \(2\ell + 1\) in the \(\theta\)-direction, we only need to treat the \(\phi\)-direction carefully. Noting that \(T_{\ell,r_0}\) is bandlimited in \(\phi\) with bandwidth \(2\ell + 1\), only frequencies \(|m| \leq N_\phi/2 + \ell \mid \sin(\phi)\) contribute to the \(N_\theta/2\) frequencies of \(T_{\ell,r_0}\) resolved by the quadrature. Therefore, to compute a low-pass \(T_{\ell,r_0}^s\) we follow the algorithm sketched below:

For each \(\theta_i, 0 \leq i < N_\theta/2\)
- \(T_k \leftarrow \frac{1}{2}T_{\ell,r_0}(\theta_i, \frac{2\pi k}{2\pi T})\), \(k = 0, \ldots, 2\ell\).
- \(\tilde{T}_k\) is Fourier series coefficient of \(T_\phi\).
- \(\tilde{s}_m \leftarrow \mathcal{F}[|m| \leq N_\phi/2 + \ell \mid \sin(\phi)]\).
- \(\tilde{T}_n\) convolution of Fourier series.
- \(\tilde{T}_n\) truncate to frequencies \(|n| \leq N_\phi/2\).
- \(T^s(\theta_i, \phi_j) \leftarrow \text{Inverse transform of } \tilde{T}_n^s\).

This algorithm yields the low-pass modified transfer function at \((\theta_i, \phi_j), 0 \leq i < N_\theta/2, 0 \leq j < N_\phi\) which can be unwrapped to the remaining points by using the spherical symmetry

\[
(\theta_i, \phi_j) = (\theta_{N_\phi/2+i}, \phi_{N_\phi-j}).
\]

Note that this method can also be simply expressed in terms of Fourier interpolations and interpolations. It is equivalent to making a Fourier interpolation of \(T_k\) to \((2\ell + 1) + (2(N_\phi/2 + \ell) + 1) - 1\) points, multiplying by a low-pass \([\sin(\phi)]\) with frequencies \(|m| \leq N_\phi/2 + \ell\), and performing a Fourier interpolation back to \(N_\phi\) points.

Finally, because sampling the transfer function at a single point is an \(O(\ell)\) operation, the algorithm as presented is \(O(\ell^2)\). The computation of the transfer function at all sample points can be accelerated to \(O(\ell^2)\) as in (Ergul & Gurel, 2006) by taking advantage of its symmetry about the \(r_0\) axis and using interpolation algorithms, but at the cost of introducing additional error.

### 3.2 Choice of Quadrature

The quadrature parameters can be constructively determined by deriving the maximum error they incur. The error in computing the integral with uniform quadrature is

\[
|\varepsilon_1| = \left| \int_0^{2\pi} \int_0^{2\pi} e^{i\pi r s}T_{\ell,r_0}^s(s) \ d\theta d\phi - \sum_{m=1}^{N_\theta} \sum_{n=1}^{N_\phi} \omega_{n,m}e^{in\pi n}T_{\ell,r_0}^s(s_{n,m}) \right|
\]

where \(s_{n,m} = [\cos(\theta_n)\sin(\phi_m), \sin(\theta_n)\sin(\phi_m), \cos(\phi_m)]\) and \(T_{\ell,r_0}^{s,L}(s_{n,m})\) is the low-pass modified transfer function described in Sec. 3.1,

\[
= \left| \int_0^{2\pi} \int_0^{2\pi} e^{i\pi r s} \left( T_{\ell,r_0}^s(s) + T_{\ell,r_0}^{\phi,H}(s) \right) d\theta d\phi - \sum_{m=1}^{N_\theta} \sum_{n=1}^{N_\phi} \omega_{n,m} \left( E_{\ell,r_0}^s(s_{n,m}) + E_{\ell,r_0}^{\phi,H}(s_{n,m}) \right) T_{\ell,r_0}^s(s_{n,m}) \right|
\]

where \(E_{\ell,r_0}^s\) consists of the low frequencies of \(e^{i\pi r s}\) which are resolved by the quadrature and \(E_{\ell,r_0}^{\phi,H}\) consists of the high frequencies not resolved by the quadrature. Since \(E_{\ell,r_0}^{s,L} = E_{\ell,r_0}^{s,H}\) is integrated exactly by a uniform quadrature,

\[
= \left| \int_0^{2\pi} \int_0^{2\pi} E_{\ell,r_0}^{\phi,H}(s) T_{\ell,r_0}^s(s) d\theta d\phi - \sum_{m=1}^{N_\theta} \sum_{n=1}^{N_\phi} \omega_{n,m} E_{\ell,r_0}^s(s_{n,m}) T_{\ell,r_0}^s(s_{n,m}) \right|
\]

where we have denoted the aliased high frequencies of \(e^{i\pi r s}\) as \(E_{\ell,r_0}^{\phi,H}(s)\),

\[
= \left| \int_0^{2\pi} \int_0^{2\pi} \left( E_{\ell,r_0}^s(s) - E_{\ell,r_0}^{\phi,H}(s) \right) T_{\ell,r_0}^s(s) d\theta d\phi \right|
\]

Transforming to Fourier space and applying the triangle inequality to prevent unpredictable cancelation effects,

\[
\leq 4\pi^2 \sum_{n=-\infty}^{\infty} \sum_{m=-\infty}^{\infty} \left| \tilde{E}(n,m) \right| \left| \tilde{T}_{\ell,r_0}^s(-n,-m) \right|
\]

This remains an accurate upper bound due to the fast decay of \(\tilde{E}\) for sufficiently large values of \(N_\phi\) and \(N_\theta\). See Fig. 1. Furthermore, this does not use phase information in computing the error, which accounts for more directional choices of \(r_0\) and \(r\) by taking advantage of the shift theorem.

#### 3.2.1 Choosing \(N_\phi\)

The worst case for \(\varepsilon_1\) in terms of \(N_\phi\) occurs when \(r\) and \(r_0\) are aligned with the \(z\)-axis. This causes all spectral information to be contained in the \(\theta\)-direction and makes \(\varepsilon_1\) a function of \(N_\phi\) only. It leads to

\[
|\varepsilon_1| \leq 4\pi^2 \sum_{m=-\infty}^{\infty} \left| \tilde{E}_{\ell,r_0}^{\phi,H}(0,m) \right| \left| \tilde{T}_{\ell,r_0}^s(0,-m) \right|
\]
and applying the Jacobi-Anger identity,

$$e^{i\kappa \cdot r} = \sum_{n=-\infty}^{\infty} i^n J_n(\kappa |r|) e^{in\varphi}$$

where $\cos(\varphi) = \hat{r} \cdot \mathbf{s}$, this simplifies this to

$$|\varepsilon I| = 4\pi^2 \sum_{m=-\infty}^{\infty} \left| J_{M(N_\theta,m)}(\kappa |r|) \right| \left| \tilde{T}_{\kappa r_0}(0, m) \right|$$

which can be used to efficiently search for a value $N_\theta$ via the following algorithm sketched below.

Choose $N_\theta^n$ sufficiently larger than $2\ell + 1$.

1. $T_k \leftarrow T_{\kappa |r_0|}^n(0, \frac{2\pi k}{N_\theta^n}), k = 0, \ldots, N_\theta^n - 1$.
2. $T_m \leftarrow$ absolute value of Fourier series of $T_k$.
3. $E_m \leftarrow |J_{M(N_\theta,m)}(\kappa |r|)|$.
4. For $N_\theta$ from $2\ell$ to $N_\theta^n$ by 2
   a. $E_{m}^* \leftarrow E_{m}^{\ast}$
   b. If $E^* \cdot T < \varepsilon / 4\pi^2$, return $N_\theta$.

Since $N_\theta^n$ is typically only a small constant larger than $2\ell + 1$, the algorithm as presented is dominated by the computation of the $O(\ell)$ modified transfer function values and requires $O(\ell^2)$ operations. Important optimizations include noting the symmetry $E^*_m = E^*_{-m}$ and $T_m = T_{-m}$ and also taking advantage of the very fast decay of $J_n$ to accelerate the inner product by not including terms which cannot contribute. More advanced searching methods also provide improved performance.

### 3.2.2 Choosing $N_\theta$

After determining an appropriate $N_\theta$, a significant optimization is made by allowing $N_\theta$ to be a function of $\phi$. This significantly reduces the number of quadrature points at negligible cost to the error. The worst case for the integration error occurs when $r$ and $r_0$ are in the $xy$-plane. Without loss of generality, suppose $\hat{r} = \hat{x}$. Then considering a constant $\phi = \phi_j$ and using the Jacobi-Anger identity, the plane wave can be expressed as

$$e^{i\kappa \cdot r} = \sum_{n=-\infty}^{\infty} i^n J_n(\kappa |r|) \sin(\phi_j) e^{in\theta}$$

and noting that $J_n(\kappa \sin(\phi_j) |r|)$ is exponentially small when

$$n \sim O(\kappa |r| \sin(\phi_j))$$

implying that we can truncate the series at $N_\theta(\phi_j) \sim O(\kappa |r| \sin(\phi_j))$ without incurring any appreciable error. Estimates of $N_\theta(\phi_j)$ can also be developed by determining when $J_n(\kappa |r| \sin(\phi_j))$ is exponentially small, as in the computation of the excess bandwidth formula (EBF) in (Chew et al., 2001). However, we find the EBF generated quadrature typically has an overestimated sampling rate.

First, it must be noted that letting $N_\theta$ be a function of $\phi_j$ requires an additional step in the computation of the modified transfer function. Sec. 3.1 computed the transfer function on a $N_\theta/2 + 1 \times N_\theta$ grid. With $N_\theta = N_\theta(\phi_j)$, the data computed for each $\phi_j$ must be Fourier interpolated from length $N_\theta$ to length $N_\theta(\phi_j)$.

To accurately and reliably compute $N_\theta(\phi_j)$ the same procedure as in Sec. 3.2.1 is applied but with $r$ and $r_0$ in the $xy$-plane. This represents the worst case for the integration error as a function of $N_\theta$. For a given $\phi_j$, we search for a $N_\theta(\phi_j)$ such that

$$|\varepsilon I| \leq 2\pi \sum_{n=-\ell}^{\ell} \sum_{m=-\infty}^{\infty} \left| J_{M(N_\theta(\phi_j),n)}(\kappa |r| \sin(\phi_j)) \right| \left| \tilde{T}_{\kappa r_0}^n(n; \phi_j) \right|$$

is bounded by a prescribed error. This is accomplished via the following sketched algorithm.

Set the poles $N_\theta(\phi_0) = N_\theta(\phi_{N_\theta/2}) = 1$.

1. Choose $N_\theta^n$ sufficiently larger than $2\ell + 1$.
2. For $\phi_j$, $j = 1, \ldots, N_\theta/2 - 1$
   a. $T_k \leftarrow T_{\kappa |r_0|}^n(0, \frac{2\pi k}{N_\theta^n}, \phi_j), k = 0, \ldots, 2\ell$.
   b. $T_n \leftarrow$ absolute value of Fourier series of $T_k$.
   c. $E_n \leftarrow |J_n(\kappa |r| \sin(\phi_j))|$.
   d. For $N_\theta(\phi_j)$ from 2 to $N_\theta^n$ by 2
      i. $E_{m}^* \leftarrow E_{M(N_\theta(\phi_j),m)}$
      ii. If $E^* \cdot T < \varepsilon / 4\pi^2$, save $N_\theta(\phi_j)$ and $j \leftarrow j + 1$.

Since $N_\theta^n$ is only a small constant larger than $2\ell + 1$, the algorithm as presented is dominated by the computation of the modified transfer function and requires
\( \mathcal{O}(\ell^3) \) operations. Again, using more advanced searching methods accelerates this algorithm and symmetry and decay properties of \( E^* \) and \( T \) should be used to improve the inner product computation. Using the EBF as an initial guess in the search for \( N_\theta(\phi_j) \) further improves the searching speed. Additionally, only half of the \( N_\theta(\phi_j) \)'s should be computed due to symmetry.

### 3.2.3 Choosing \( |x| \) and \( |r_0| \)

The previous algorithms require representative values of \( |x| \) and \( |r_0| \) for each level of the tree. The worst-case transfer vectors, \( r_0 \), are well known to be those with minimum magnitude. If \( a_l \) is the box size at level \( l \), then \( |r_0| = 2a_l \) is the smallest transfer vector length in the common one buffer box case.

The worst case values of \( |x| \) is well known to be the largest. For a box of size \( a_i \), \( |x| \leq a_i \sqrt{3} \). However, using \( |x| = a_i \sqrt{3} \) in the previous methods is a very conservative choice. This case only occurs when two points are located in the exact corners of the boxes - rare indeed. See Fig. 2. Instead, we let \( |x| = a_i \sqrt{3} \) for some \( \alpha \in [0,1] \). A high \( \alpha \) will strongly guarantee an upper bound on the error generated by the quadrature, but the points which actually generate this error will be more and more rare. A lower value of \( \alpha \) will yield a smaller quadrature, but more points may fall outside the radius \( |x| \) where the upper bound on the error is guaranteed. A typically reliable choice appears to be \( \alpha = 0.8 \). A study of this parameter will appear in a subsequent paper.

![Diagram](image)

Fig. 2: The worst case \( x \) and \( r_0 \), projected from the 3D box. Here, \( |r_0| = 2a_l \) and \( i \) and \( j \) on at opposite corners of the box so that \( |x| = |r_{ic}| + |r_{cj}| = a_l \sqrt{3} \).

### 3.2.4 Number of Quadrature Points

Recall from Sec. 2.1 that the typical approach in the standard MLFMM is to use \( N + 1 \) uniform points in the \( \theta \) direction and \( N_{\phi}^{\ell+1} \) Gauss-Legendre points in the \( \phi \) direction so that all \( Y_n^m \), \(-n \leq m \leq n, 0 \leq n \leq N \) are integrated exactly. In (Koc et al., 1999), Chew et. al. takes \( N_{\phi}^{\ell+1} = \ell + 1 \), which accounts for the rapid decay of the spherical harmonics in the plane wave expansion. This results in

\[
M_\theta = 2(\ell + 1)^2 \approx 2\ell^2
\]

quadrature points.

For a given Gegenbauer series truncation \( \ell \), the total number of active quadrature points required in the Fourier based MLFMM is approximately

\[
M_f \approx \frac{N_\theta}{2} \sqrt{\int_0^\pi N_\theta(\phi) \, d\phi} \approx (\ell + C) / (2(\ell + 1))
\]

where \( C \geq 1 \) is a small integer dependent on \( \ell \) numerically computed from the method in Sec. 3.2.1. Keeping only the dominant term,

\[
M_f \approx \frac{4}{\pi} \ell^2
\]

Thus, the method presented in this paper uses approximately 0.6 times the number of quadrature points in the standard MLFMM. However, it may be possible that the same \( N_\theta \) optimization can be applied to the standard MLFMM for the same reasons it was applied in Sec. 3.2.2 to reduce the standard quadrature to a comparable size.

### 3.3 Interpolation and Anterpolation

Most importantly, the Fourier based MLFMM directly uses FFTs in the interpolation and anterpolation steps. This makes the time critical upward pass and downward pass especially fast and easy to implement.

Characterize a quadrature by an array of length \( N_\phi/2 + 1 \),

\[
Q = [1, N_\theta(\phi_1), \ldots, N_\theta(\phi_{N_\phi/2-1}), 1]
\]

noting that \( N_\theta(\phi_j) = N_\theta(\phi_{N_\phi/2-j}) \) and \( N_\theta(\phi_j) = N_\theta(\phi_{N_\phi/2-j}) \). A transform of the data \( F(\theta_1, \phi_j) \) sampled on a quadrature \( Q \) to a quadrature \( Q' \) is performed by a sequence of Fourier interpolations and anterpolations. Let

\[
N_\theta = \max( \max_{0 \leq j \leq N_\phi/2} N_\theta(\phi_j), \max_{0 \leq j \leq N_\phi/2} N_\theta(\phi_j) )
\]

Then, the following steps, as illustrated in Fig. 3, perform an exact interpolation/anterpolation using only FFTs.

1. For each \( \phi_j, 0 \leq j \leq N_\phi/2 \), Fourier interpo- late the data \([F(\theta_i = 0, \ldots, N_\theta(\phi_{j-1}), \phi_j)]\) from length \( N_\theta(\phi_j) \) to \( N_\theta \).
2. For each \( \theta_i, 0 \leq i < N_\theta/2 \), wrap the data to construct the periodic sequence from the rest of the line \([F(\theta_i, \phi_{j=0, \ldots, N_\phi/2}), F(\theta_{i+N_\theta/2}, \phi_{j=N_\theta/2-1, \ldots, 1})]\).
3. For each $\theta_i$, $0 \leq i < N_\theta/2$, Fourier interpolate the data $[F(\theta_i, \phi_j=0, ..., N_\phi-1)]$ from length $N_\theta$ to $N_\theta'$.

4. For each $\theta_i$, $0 \leq i < N_\theta/2$, unwrap the data $[F(\theta_i, \phi_j=0, ..., N_\phi'-1)]$ to construct the sequences $[F(\theta_i, \phi_j=0, ..., N_\phi'/2)]$ and $[F(\theta_i+N_\theta/2, \phi_j=0, ..., N_\phi'/2)]$.

5. For each $\phi_j$, $0 \leq j < N_\phi/2$, Fourier anterpolate the data $[F(\theta=0, ..., N_\theta(\phi_j)-1, \phi_j)]$ from length $N_\theta$ to $N_\theta'(\phi_j)$.

Fig. 3: The data profile at each step in an anterpolation from a large quadrature $Q$ with $N_\theta = 30$ to a smaller quadrature $Q'$ with $N_\theta' = 24$. The data corresponding to a pole has been darkened for clarity.

4. NUMERICAL RESULTS

A direct computation was used to compute the optimal Gegenbauer truncation $\ell$ and the methods described Sec. 3.2 were used to construct a quadrature. For a given box size $s$, the quadrature and truncation are constructed with $|r| = 0.8s\sqrt{3}$, $|r_0| = 2s$, and target error $\epsilon$. The total measured error, $\epsilon$, is

$$\epsilon = \epsilon_n |r+\epsilon_0| [r + r_0] - \sum_{m=1}^{N_\theta} \sum_{n=1}^{N_\phi} \omega_n,m e^{i\xi n, \phi m} r F^{n, r}_{n, \theta} (s, n, m)$$

The total Gegenbauer truncation error, $\epsilon_G$, is

$$\epsilon_G = \frac{\epsilon_n |r+\epsilon_0| [r + r_0]}{[r + r_0]} - \sum_{k=0}^{\ell} (-1)^n (2n+1) G_n(r, r_0)$$

The total integration error, $\epsilon_I$, is

$$\epsilon_I = \epsilon - \epsilon_G$$

The plotted errors in Fig. 4 represent the maximum error found over many directions $\hat{r}$ and magnitudes $|r| \leq 0.8s\sqrt{3}$, where $s$ is the box size.

Figure 5 shows the recorded running times of the Fourier based MFMM and the direct matrix-vector product on an Intel Core(2) Quad CPU Q9450 2.66GHz with 4GB of RAM. Note that the intersection point is less than $N = 10,000$.

5. CONCLUSIONS

We have proposed using a Fourier basis in the spherical variables $\theta$ and $\phi$: $e^{i\theta} e^{i\phi}$. By modifying the Helmholtz kernel approximation and using a uniform quadrature we can take advantage of very fast, exact, and well-known FFT interpolation and anterpolation methods. By exploiting symmetries, the number of uniform quadrature points required is almost identical to the number of Gauss-Legendre quadrature points typically used for spherical quadratures and can be improved to use even fewer points. However, the Fourier based MFMM requires careful changes to the precomputation stage where the modified transfer functions, $T_{\ell, r, \theta}(s)$, are computed. Since $|\sin(\phi)|$ is not smooth, we must accurately precompute a bandlimited version of the modified transfer function.

The Fourier based MFMM approach has a number of advantages. Since the interpolation and anterpolation algorithms are exact, significant errors in the algorithm are a function of only the truncation parameter $\ell$ and the quadrature, specifically the bandwidth in the $\phi$-direction. The truncation error $\epsilon_I$ has been extensively studied and is well understood. The integration error $\epsilon_I$ can be accounted for during the precomputation stage and precise bounds on the final error can be made. The error analysis relies on well-known properties of the Fourier basis. We have used the results in constructive algorithms to determine the appropriate quadrature and constrain the MFMM error a-priori. We find that bounding integration error is accurate and efficiently exploitable to search for the optimal quadrature. Although not required in the Fourier based MFMM, computing the optimal quadrature can aid in error analysis and im-
prove performance. Since the Fourier Based MLFMM uses a uniform quadrature and well-known FFT algorithms, the time-critical stages of the algorithm are much easier to implement. Modern FFT packages are very common, extremely fast, numerically exact, and easy to use. FFTs scale much better than other fast interpolation algorithms of constant accuracy. Although the asymptotic complexity is not improved, the smaller constant in the FFT’s $O(N \log N)$ complexity yields improved running times in relation to algorithms of comparable accuracy.

References


