MICROCOPY RESOLUTION TEST CHART
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Experience with Berkeley UNIX* Interprocess Communications

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Abstract

At many universities and institutions throughout the world it is now very common to have a network of computers, each running the Berkeley 4.3BSD version of UNIX or an equivalent version such as ULTRIX. This paper is to help users of these versions of UNIX to explore and experiment with the interprocess communications and networking facilities. We present a series of client/server programs that can be used as a model for writing distributed applications. We describe how users can test and experiment with these programs. Readers are assumed to be familiar with the C programming language and some version of UNIX.

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1 Introduction

Currently, many universities and research organizations have a network of computers each running 4.3BSD version of UNIX or a compatible version such as ULTRIX or DG/UX. A major addition to these versions of UNIX is the Interprocess Communications (IPC) facilities that allow unrelated processes on the same or different machines to communicate and exchange messages. This makes it possible to write distributed applications and provide a better environment for users to collaborate and access remote resources and utilities (See for example, Abdel-Wahab et. al. [1,2]).

In this paper we introduce the important basic features of the IPC facilities and show how to write distributed applications based on these facilities. The users are assumed to be familiar with some version of UNIX (see for example, Kernighan and Pike [4]), and the C programming language (Kernighan and Ritchie[5]). A tutorial introduction to IPC is given in Coffield [3] and Sechrest [8]. For a comprehensive coverage see Leffler et. el. [6].

There are several domains of communication, the most frequently used ones are the internet domain (INET) and the UNIX domain. The INET domain employs the DARPA standard protocols TCP/IP and UDP/IP (Leiner et. el. [7]). The UNIX domain is basically used for efficient communication between unrelated processes residing on the same machine. The INET domain allows processes running on different machines (as well as on the same machine) to communicate. In most cases the INET domain is preferable to the UNIX domain, since applications written in the INET domain may run on any configurations of machines.

In each domain, there are two methods of communication: stream and datagram. In the stream method the two communicating processes have to establish a connection or a virtual circuit before they can exchange messages. This style of communication provides bidirectional, reliable, sequenced and unduplicated flow of data. On the other hand, datagram communication does not require any connection between the two processes as each message is addressed individually but
there is no guarantee regarding delivery, sequencing, or duplication of messages. The choice between the two methods is usually based on the application semantics and the required performance.

This paper is organized as follows. Section 2 is concerned with the INET domain, the first part deals with stream communication and the second with datagram communication. Each part is described by giving a pair of server and client programs. The programs are tested using two terminals (or two windows of a single terminal) connected to two different hosts or to a single host. Section 3 is devoted to the UNIX domain. Section 4 deals with the cases where a server communicates with two or more clients and input multiplexing is required. Section 5 is our conclusion.

2 Internet Domain

In this section we describe how two processes in the INET domain communicate with each other. The first step for a process to communicate with another is to create a socket. A socket is an end point of communication and is created with socket system call:

\[ s = \text{socket}(\text{domain, type, protocol}); \]

For example, to create a stream socket in the internet domain we use:

\[ s = \text{socket}(\text{AF_INET, SOCK_STREAM, 0}); \]

Here, we specify 0 in place of the protocol arguments so the system will use the default standard protocol.

We distinguish between the two communicating processes by calling one the server and the other the client. The server process has to "bind" its socket to a known "name" so that the client process can use this name to establish a connection with the server. The name format in the INET domain consists of two parts host_address and port_number. The host_address is a 32-bit network wide address assigned to the machine (for example, 128.109.135.1 is the address of a machine at
North Carolina State University called “ncsu” and 128.109.136.82 is for a machine at UNC-Chapel Hill called “unc”). The port number is an integer in the range 0-50,000 (the subrange 0-1024 is reserved for privileged use, for example, the remote login program `rlogin` uses port number 513 and the file transfer program `ftp` uses port number 21). To bind a socket `s` to a name we use:

```c
bind(s, name, sizeof(name));
```

`name` is a structure where its fields are to be filled out appropriately as in the following example.

```c
struct sockaddr_in name;

do

name.sin_family = AF_INET;
name.sin_addr.s_addr = INADDR_ANY;
name.sin_port = 0;
bind(s, &name, sizeof(name));
```

Here, the constant `INADDR_ANY`, means any valid address of the host. If we know the host address, e.g., 128.109.136.82, we use it as:

```c
name.sin_addr.s_addr = inet_addr("128.109.136.82");
```

and if the host name is known, e.g., "ncsu" it is used as:

```c
struct hostent *hp;

do

hp = gethostbyname("ncsu");
bcopy(hp->h_addr, &name.sin_addr.s_addr, hp->h_length);
```

To choose a port number, we either use:

```c
name.sin_port = 0;
```

in which case, the system selects the next free port, or if you know the number of a particular free port, e.g., 1234, it is used as:

```c
name.sin_port = 1234;
```

If the system selects a free port, the selected port number may be displayed as follows:
length = sizeof(name);
getsockbyname(s, &name, &length);
printf("%d", ntohs(name.sin_port));

After the server creates a socket and binds it to an INET name, it executes:
listen(s, n);
to specify, as n, the maximum number of outstanding connections to be queued awaiting acceptance by the server. This system call is non-blocking, it just sets up the socket and makes it ready for accepting connections.

The client process creates its own socket, sc, but it does not need to bind it to a name like the server. Then it initiates a connection to the server process using the connect system call:
connect(sc, &server_name, sizeof(server_name));

When the connection request arrives at the server, it is accepted using:
s1 = accept(s, 0, 0);

Where s1 is an “auxiliary” socket descriptor to be used in communicating with the client process. This system call causes the server process to be blocked waiting for the client process to be connected. The server may accept other connections as shown in Fig. 1.

Once a connection is established between a client and a server, data may flow in both directions. The client may send data using:
send(sc, data, sizeof(data), 0);
and the server may receive the arrived data using:
recv(s1, buffer, sizeof(buffer)), 0);

When the server and the client are done with their conversation, each may close its end of the connection using the close call:
close(socket);

Example Programs

In this paper we choose the names of our example programs to reflect their functions using the following abbreviations:
Fig. 1: Server and two clients
Each program includes a "def" file that contains:

```
#include <sys/types.h>
#include <sys/socket.h>
#include <sys/un.h>
#include <stdio.h>
#include <netinet/in.h>
#include <netdb.h>
#include <signal.h>
#include <sys/time.h>
#define TRUE
```

### 2.1 INET Virtual-circuit programs

In order to test and practice with the concept of communication between two processes in the INET domain using a virtual circuit connection, we wrote two programs in Fig. 2: a server program IVS and a client program IVC.

The server process IVS starts by creating a socket ss (lines 10-14), then binds it to any free port in the host machine (lines 16-23), advertises the selected port number by displaying the message: socket has port # .... and finally listens for a connection attempt by the client process IVC (line 33). When the client process attempts to connect to the server process, the connection is accepted in socket s1 (line 34). Following the connection establishment, the server enters a
0 #include "def"
1 main()
2 {
3     int ss, s1;
4     struct sockaddr_in name;
5     char buf[1024];
6     int sc;
7     int length;
8     struct sockaddr from;
9     /* Create socket from which to read. */
10    ss = socket(AF_INET,SOCK_STREAM,0);
11    if ( ss<0 ) {
12        perror("opening virtual circuit socket");
13        exit(-1);
14    };
15    /* form INET socket name and bind it to ss */
16    name.sin_family = AF_INET;
17    name.sin_addr.s_addr = INADDR_ANY;
18    name.sin_port = 0;
19    if ( bind( ss, &name, sizeof(name) ) ) {
20        close(ss);
21        perror("binding name to virtual circuit socket");
22        exit(-1);
23    };
24    /* Find out assigned port number and print it out */
25    length = sizeof(name);
26    if ( gethostname (ss,&name, &length) ) {
27        perror("getting socket name");
28        exit(0);
29    }
30    printf("Socket has port #\%d", ntohs(name.sin_port));
31    /* listen for a connection from the client process via */
32    printf("... waiting for connection ...
");
33    listen(ss);
34    s1 = accept(ss, 0, 0);
35    printf("Connected to client\n");
36    printf("Waiting for messages ...
");
37    do { //
38        if ( (recv(s1,buf,sizeof(buf),0) ) < 0 )
39            perror("receiving virtual circuit packet");
40        if ( ce > 0 ) {
41            buf[ce] = NULL;
42            printf("message received: %s
", buf);
43            } else {
44                printf("message received: EOF \n");
45                close(s1);
46                close(sc);
47                printf("... disconnect\n");
48                exit(0) ;
49            }
50        
51    while (TRUE);
52
53    }
54
55    printf("EOF... disconnect\n");
56    close(ss);
57    exit(0);
58
59    printf("... waiting for connection ...
");
60    /* Create socket on which to send. */
61    sc = socket(AF_INET,SOCK_STREAM,0);
62    if ( sc<0 ) {
63        perror("opening virtual circuit socket");
64        exit(-1);
65    };
66    /* Construct name of socket to send to and connect to it */
67    server.sin_family = AF_INET;
68    if ( (hp = gethostname(argv[1])) == NULL ) {
69        close(sc);
70        fprintf(stderr, "Can't find host \%s, argv[1]\n");
71        exit(-1);
72    }
73    bcopy (hp->h_addr, &server.sin_addr, hp->h_length);
74    server.sin_port = htons(atoi(argv[2]));
75    printf("server\.sin_family = AF_INET\n");
76    if ( connect(sc, &server, sizeof(server)) < 0 ) {
77        close(sc);
78        perror("connecting stream socket");
79        exit(0);
80    }
81    bcopy (hp->h_addr, &server.sin_addr.sin_addr, hp->h_length);
82    server.sin_port = htons(atoi(argv[2]));
83    printf("server\.sin_family = AF_INET\n");
84    if ( connect(sc, &server, sizeof(server)) < 0 ) {
85        close(sc);
86        perror("connecting virtual circuit socket");
87        exit(-1);
88    }
89    /* Send message. */
90    do {
91        read(0,buf,sizeof(buf));
92        if (send(sc, buf, rc, 0 ) <0 )
93            perror("sending virtual circuit message\n");
94        } while (rc > 0);
95    printf("EOF... disconnect\n");
96    close(sc);
97    exit (0);
98
99    Fig. 2a: /IC (client)
100   Fig. 2b: INET Virtual-circuit programs

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Fig. 3a: Testing using two terminals

Fig. 3b: Testing using two windows

Fig. 3: Testing the server and the client programs
loop (lines 37-52), receiving (lines 38-39) and displaying (lines 40-43) any message that arrives on the connection until the client sends an EOF signaling the end of conversation. Upon receiving the EOF message, the IVS process terminates (lines 44-50).

In Fig. 2b, the client process IVC is initiated by typing:

```
% IVC host_name port_number
```

where host_name is the name of the machine running the server process IVS, and port_number is the number displayed by the server process. The IVC process starts by creating a socket sc (lines 11-15), forming the server_name (lines 17-24) from the host name (argv[1]) and the port number (argv[2]). Then it connects the socket sc to the server_name (lines 25-29) and enters a loop (lines 33-38) reading any message typed on the keyboard (line 34) and sending it to the server process (lines 35-36) until CTRL-D is typed to end of conversation with the server.

To test how the two programs interact with each other, we may use two terminals T, and T, beside each other as shown in Fig. 3a. We use terminal T, to run the server program. The server process displays the selected socket number (e.g., Socket has port # 1245) and waits for a connection by the client process. We use the second terminal T, to run the client process IVC. To run IVC we provide two arguments: the remote machine name and the port number where the server is waiting. After the client process is connected to the server process it reads any character typed on the keyboard and sends it out to the server process IVS. Whatever typed on T, will appear instantly at T,. For example, if we type the message “Hello World!” on T, it will appear on T, as “Message received: Hello World!”.

To end the session we type CTRL-D on T, and both processes will terminate after displaying the message “EOF ... disconnect”.

If we have access to only one machine, we can still conduct the above test, since in the INET domain processes in the same machine or in different machines can communicate with each other. We use terminal T, to run the server IVS and

---

1In UNIX an End-Of-File is signaled by typing CTRL-D
after it displays the port number, use terminal $T_s$ to run the client, where the
machine name is the one running both the server and the client processes.

If we like to use a single terminal for this test, we need to create two “windows”
as shown in Fig. 3b. Windows can be created using any available window man-
agement system such as the Berkeley UNIX 4.3BSD window program that runs on
any ASCII terminal, or the MIT X-windows that runs on a variety of workstations
including SUNs and DEC VAXs. In our test, the top window is used as terminal
$T_s$ while the bottom window is used as terminal $T_c$.

\section{INET Datagram Programs}

To create a datagram socket in the INET domain we use:

\begin{verbatim}
 s = socket (AF_INET, SOCK_DGRAM, 0):
\end{verbatim}

Fig. 4a shows the server program $IDS$ and Fig. 4b shows the client program
$IDC$. The server process $IDS$ starts by creating socket $ss$ (lines 11-15), binds it
to an INET port (lines 17-24) and advertises the chosen port (lines 26-31). Then
enters a loop (lines 34-46) where it receive messages sent by the client process. In
contrast to the virtual circuit program ($IVS$), this program does not execute listen
to queue connections and does not accept connections. Also, instead of using recv,
it uses recvfrom to read the incoming messages (line 35).

In Fig. 4b, the client process $IDC$ creates a socket $sc$ (lines 11-15), forms the
Then sends every line as an individual datagram (line 26-32) using:

\begin{verbatim}
 sendto(sc, message, message_length, 0, &server, sizeof(server)):
\end{verbatim}

Note that the server name $&server$ is attached with every message going from
process $IDC$ so that the server process $IDS$ can identify the source of the message.
The two processes terminate when CTRL-D is typed.

To test these two program we follow the same procedure we described earlier
for testing the virtual circuit programs $IVS$ and $IVC$. 

11
```c
#include "def"
main()
{
  let
Ss;
  let
name;
  argc;
  argv;
  length;
  struct
sockaddr.in

  /* Create socket from which to read. */
  ss = socket (AF_INET,SOCK_DGRAM,0);
  if ( ss<0 ) { 
    perror("opening datagram socket");
    exit(-1);
  }
  /* form INET socket name and bind it to ss */
  name.sin_family = AF_INET;
  name.sin_addr.s_addr = INADDR_ANY;
  name.sin_port = 0;
  if ( bind ( ss, &name, sizeof(name)) ) {
    close(ss);
    perror("binding name to datagram socket");
    exit(-1);
  }
  /* Find out assigned port number and print it out */
  length = sizeof(name);
  if ( getsockname ( ss,&name,&length ) ) {
    perror("getting socket name");
    exit(0);
  }
  printf("Socket has port %d", ntohs(name.sin_port));
  /* Read from the socket */
  printf("... waiting for messages ...
");
  for ( ; ) {
    if ( (cc=recvfrom( ss,buf,sizeof(buf), 0,0,0)) < 0 ) 
      perror("receiving datagram packet");
    if ( cc > 0 ) {
      buf[cc] = NULL;
      printf("message received: %s", buf);
    } else{
      printf("message received: EOF ... exit");
      close(ss);
      exit(0);
    }
  }
}

Fig. 4a: IDS (server)

Fig. 4b: IDC (client)

Fig. 4: INET datagram programs
```
3 Unix Domain

One difference between the UNIX and the INET domains is the address format. In UNIX domain, addresses are file path names, e.g., /tmp/file1. To create a stream socket in the UNIX domain we use:

\[ s = \text{socket} (AF\_UNIX, \text{SOCK\_STREAM}, 0); \]

and to create a datagram socket we use:

\[ s = \text{socket} (AF\_UNIX, \text{SOCK\_DGRAM}, 0); \]

3.1 UNIX Virtual-circuit Programs

Fig. 5 shows the code for the server program UVS and the client program UVC.

The UVS process creates a socket ss (lines 13-16), binds it to a UNIX file name (lines 18-24). Then it listens to socket ss (line 27), and when a client process initiates a connection, it is accepted into an auxiliary socket s1 (line 28). After that it enters a loop (lines 31-46) to receive (lines 32-33) and display messages sent by the client process (lines 34-37) until an EOF is received. Upon receiving EOF, it closes the socket and deletes the associated file (lines 41-42).

The client process UVC in Fig. 5b creates socket sc (lines 8-12), and forms the server name (lines 14-15). Contrast this with the INET program IVC, where the server name is not known until execution time and is formed from the command line arguments: argv[1] and argv[2]. After connecting to the server process, it repeatedly accepts the user input and sends it to the server process until CTRL-D is typed.

To test the interaction between UVS and UVC, we may use two terminals (connected to the same host), or create two windows on one terminal as described before. In invoking the client process UVC, we provide no arguments, since the host name is the same for both the client and the server and the socket is bound to a path name known to the client process in priori.
```c
#include "def"
#include "derf"

main() {
    struct sockaddr_un name;
    struct sockaddr from;
    char buf[1024];

    char *pathname = "tmp/unsock";
    /* Create socket from which to read. */
    ss = socket(AF_UNIX.SOCK_STREAM,0);
    if (ss<0) {
        perror("opening virtual circuit socket");
        exit(-1);
    }
    /* Form name and bind it to ss */
    name.sun_family = AF_UNIX;
    strcpy(name.sun_path,pathname);
    if (bind(ss, &name, sizeof(struct sockaddr_un)-1) ) {
        close(ss);
        perror("binding name to ss");
        exit(-1);
    }
    /* listen for a connection from the client process uvc */
    printf("... waiting for connection ...
");
    si = accept(ss, 0, 0);
    printf("connected to client\n");
    printf("... waiting for messages ...
");
    do {
        if ( (cc=recv(si,buf,sizeof(buf),0)) < 0 )
            perror("receiving virtual circuit message");
        else {
            printf("message received: %sn", buf);
            close(ss);
            unlink(pathname);
            exit(0);
        }
    } while (TRUE);
}

Fig. 5a: UVS (server)

```
3.2 **UNIX Datagram Programs**

To experiment with datagrams in the **UNIX** domain, see Fig. 6 for the server program **UDS** and the client program **UDC**. The server process **UDS** creates a socket (lines 11-15), binds it to name (lines 17-23), and without waiting for connections, enters a loop (lines 26-41) where it receives the messages sent by the client process and displays it until an **EOF** is received.

The client process in Fig. 6b creates a socket (lines 8-12) and sends messages with the server name attached to each message (line 20). To test the interaction between these two programs we follow the same procedure for testing the virtual circuit programs in the **UNIX** domain.

4 **Input Multiplexing**

If a server process is connected concurrently to more than one client process, it needs to multiplex the input received from all of the clients. In this section we present an example in the **INET** domain.

Fig. 7 shows the code for a server program called **IV_SEL** which is essentially the **IVS** program except that it serves two clients instead of just one. The program uses the select system call to monitor the arrival of data from any of the clients. As usual, the program starts by creating a socket (lines 9-13), binds it to a port (lines 15-20), announces the selected port (lines 22-27) and listens for connections (line 30). It accepts the first connection on socket **s1** (line 31) and the second connection on socket **s2** (line 34). Fig. 1 illustrates the relationship between the server and its two clients.

To monitor the two sockets **s1** and **s2** for the arrival of data, we use a bit mask called **read_template** (line 3). The mask is cleared using the **FD_ZERO** macro (line 40) and we add **s1** and **s2** to the mask using the **FD_SET** macro (line 41-42). The select call:

\[
\text{nb} = \text{select} (\text{FD_SETSIZE, \&read\_template, 0, 0, \&wait});
\]
#include "def"

struct sockaddr_un sockname;
char buf[1024];
int cc;
struct sockaddr from;
int fromlen;
int ss;

#include "def"

char *pathname = "/tmp/unxsock";
/* Create socket from which to read. */
ss = socket(AF_UNIX, SOCK_DGRAM, 0);
if (ss < 0) { perror("opening datagram socket"); exit(-1);
};
/* form socketname and bind it to ss */
sockname.sun_family = AF_UNIX;
strcpy(sockname.sun_path, pathname);
if (bind(ss, &sockname, sizeof(struct sockaddr_un) - 1) ) { close(ss);
    perror("binding socketname to ss");
    exit(-1);
};
/* Read from the socket */
printf("... waiting for messages ... ");
do { if ( (cc = recv(ss, buf, sizeof(buf), 0) ) < 0 ) { perror("receiving virtual circuit message");
    if (cc > 0) {buf[cc] = NULL;
        printf("message received: \n", buf);
    } else { printf("message motived: EOF ");
        printf("... exit()");
        close (ss);
        unlink(pathname);
        exit(0);
    }
} while (TRUE);

0 #include "def"

int sc;
struct sockaddr_un name;
char buf[512];
int rc;
/* Create socket on which to send. */
sc = socket(AF_UNIX, SOCK_DGRAM, 0);
if (sc < 0) { perror("opening datagram socket");
    exit(-1);
};
/* Construct name of socket to send to. */
name.sun_family = AF_UNIX;
strcpy(name.sun_path, "/tmp/unxsock");
printf("... type any message, to finish type CTRL-D \n");
/* Send message. */
do { if (sendto(sc, buf, sizeof(buf), 0, &name, sizeof(struct sockaddr_un)-1) < 0 ) { close(sc);
    unlink(pathname);
    printf("sending datagram message");
} while (rc > 0);
printf("EOF... exit()");
exit(0);

Fig. 6b: UDC (client)

Fig. 6a: UDS (server)

Fig. 6: UNIX Datagram programs
is blocked for an amount of time equal to the value of the variable wait (in Fig. 7, lines 37-38, the value is set to 1 second). When the select call returns, then either data has arrived on one of the sockets, the wait time-out has been exceeded, or an error has occurred. If nb, the value returned by select, is greater than zero then new data has arrived on at least one of the multiplexed sockets. We use the FD_ISSET macro to test whether there is data available at a particular socket (line 49 and 62). For example, if socket s1 has data, it is received (lines 50-51) and displayed (lines 52-55).

When a client terminates the conversation with the server by sending an EOF, the corresponding socket should not be added to the bit mask. This is achieved by setting a flag (e.g., eof1 in line 59) when EOF is received. When a flag is set, the corresponding socket is not monitored any more (e.g., line 41). When all clients have sent an EOF, the server terminates (lines 77).

It should be noted that if the value of wait is 0, the select statement will return immediately, i.e., it behaves as if it is "polling" the sockets. On the other hand if &wait is replaced with the constant NULL, then the call is blocked indefinitely until data arrives on one of the monitored sockets.

To test the IV_SEL program, we may use three different terminals (alternatively, we may use three windows of one terminal) connected to either one, two or three machines. Assume that we have access to three machines named: "ncsu", "mcnc" and "unc" and we are using three windows of one terminal as shown in Fig. 8. We use the top window to run IV_SEL process on "ncsu", the middle window to run a client process IVC on the "mcnc" machine and the bottom window to run another client process IVC on "unc".

We should notice that any message typed on the middle or the bottom window will instantly appear at the top window. To terminate a client type CTRL-D in the corresponding window. The IV_SEL program terminates when all the clients have sent EOF messages (lines 78-81). The test may be conducted using other combinations of terminals, windows and machines.
```c
#include "def"
main()
{ 
    fd_set read_template;
    struct timeval wait;
    struct sockaddr_in server;
    char buf[1024];
    int sock, length, s1, s2, nb, cc, ready1, ready2;

    /* Create socket */
    sock = socket(AF_INET, SOCK_STREAM, 0);
    if (sock < 0) ("perror("opening stream socket")
           exit(0);
    }

    /* Name socket using wildcards */
    server.sin_family = AF_INET;
    server.sin_addr.s_addr = INADDR_ANY;
    server.sin_port = 0;
    if (bind(sock, &server髭zeof(server)))
         perror("binding stream socket");
    }

    /* Find out assigned port number and print it out */
    length = sizeof(server);
    if (getsockname(sock, &server, &length))
         perror("getting socket name");
         exit(0);
    
    printf("Socket has port #%d", ntohs(server.sin_port));
    printf("...waiting for connection ...");
    
    /* Start accepting connections */
    listen(sock, 5);
    s1 = accept(sock, 0);
    
    /* Print connected to first client\n");
    s2 = accept(sock, 0);
    
    printf("connected to second client\n");
    wait.tv_sec = 1; wait.tv_usec = 0;
    do
    { FD_ZERO(&read_template);
      if (FD_ISSET(s1, &read_template))
          
    if (cc < 0) /* handling virtual circuit packet */
        (cc > 0)
            
    else
        printf("message from s2: EOP\n");
        close(s2);
        
    } while (ready1 & ready2);
        
    printf("...exit\n");
    close(s1);
    close(s2);
    exit(0);
} 
```

Fig. 7: IV_SEL program
% IVS_SEL

Socket has port# 1234
... waiting for connection ...
connected to first client
connected to second client
message from s1: Hello, first client
message from s2: Hello, second client
message from s1: EOF
message from s2: EOF
... exit...

% IVC ncsu 1234
Connected to server
type message(CTRL-D to finish)
Hello, first client
EOF... disconnect

% IVC ncsu 1234
Connected to server
type message(CTRL-D to finish)
Hello, second client
EOF... disconnect

Fig. 8: Testing input multiplexing from two clients
5 Conclusion

We have surveyed the basic features of the Interprocess Communications facilities available in the popular Berkeley version of UNIX. Then we presented a series of simple server/client programs that can be used as a model of writing practical distributed applications over a network of computers each running 4.3BSD or compatible version of UNIX. We also show how to test and experiment with the programs in a variety of configurations. After the user masters the concepts and techniques presented in our programs, we recommend looking at the source code of some widely used utilities such as the remote login program (rlogin) and the file transfer program (ftp). While the paper focuses on a specific system, the experience and knowledge gained are valuable in dealing with other distributed systems and environments.

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