Distributed convergence to Nash equilibria in two-network zero-sum games

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Abstract

This paper considers a class of strategic scenarios in which two networks of agents have opposing objectives with regards to the optimization of a common objective function. In the resulting zero-sum game, individual agents collaborate with neighbors in their respective network and have only partial knowledge of the state of the agents in the other network. For the case when the interaction topology of each network is undirected, we synthesize a distributed saddle-point strategy and establish its convergence to the Nash equilibrium for the class of strictly concave-convex and locally Lipschitz objective functions. We also show that this dynamics does not converge in general if the topologies are directed. This justifies the introduction, in the directed case, of a generalization of this distributed dynamics which we show converges to the Nash equilibrium for the class of strictly concave-convex differentiable functions with globally Lipschitz gradients. The technical approach combines tools from algebraic graph theory, nonsmooth analysis, set-valued dynamical systems, and game theory.

Key words: adversarial networks, distributed algorithms, zero-sum game, saddle-point dynamics, Nash equilibria

1 Introduction

Recent years have seen an increasing interest on networked strategic scenarios where agents may cooperate or compete with each other towards the achievement of some objective, interact across different layers, have access to limited information, and are subject to evolving interaction topologies. This paper is a contribution to this body of work. Specifically, we consider a class of strategic scenarios in which two networks of agents are involved in a zero-sum game. We assume that the objective function can be decomposed as a sum of concave-convex functions and that the networks have opposing objectives regarding its optimization. Agents collaborate with the neighbors in their own network and have partial information about the state of the agents in the other network. Such scenarios are challenging because information is spread across the agents and possibly multiple layers, and networks, by themselves, are not the decision makers. Our aim is to design a distributed coordination algorithm that can be used by the agents to converge to the Nash equilibrium. Note that, for a 2-player zero-sum game of the type considered here, a pure Nash equilibrium corresponds to a saddle point of the objective function.

Literature review. Multiple scenarios involving networked systems and intelligent adversaries in sensor networks, filtering, finance, and wireless communications (Kim and Boyd, 2008; Wan and Lemmon, 2009) can be cast into the strategic framework described above. In such scenarios, the network objective arises as a result of the aggregation of agent-to-agent adversarial interactions regarding a common goal, and information is naturally distributed among the agents. The present work has connections with the literature on distributed optimization and zero-sum games. The distributed optimization of a sum of convex functions has been intensively studied in recent years, see e.g. (Nedic and Ozdaglar, 2009; Wan and Lemmon, 2009; Johansson et al., 2009; Zhu and Martínez, 2012). These works build on consensus-based dynamics (Olfati-Saber et al., 2007; Ren and Beard, 2008; Bullo et al., 2009; Mesbahi and Egerstedt, 2010) to find the solutions of the optimization problem in a variety of scenarios and are...
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designed in discrete time. Exceptions include (Wang and Elia, 2010, 2011) on continuous-time distributed optimization on undirected networks and (Gharesifard and Cortés, 2012c) on directed networks.

Regarding zero-sum games, the works (Arrow et al., 1958; Maistrovskii, 1977; Nedic and Ozdaglar, 2009) study the convergence of discrete-time subgradient dynamics to a saddle point. Continuous-time best-response dynamics for zero-sum games converges to the set of Nash equilibria for both convex-concave (Hofbauer and Sorin, 2006) and quasiconvex-quasiconcave (Barron et al., 2010) functions. Under strict convexity-concavity assumptions, continuous-time subgradient flow dynamics converges to a saddle point (Arrow et al., 1951, 1958). Asymptotic convergence is also guaranteed when the Hessian of the objective function is positive definite in one argument and the function is linear in the other (Arrow et al., 1958; Feijer and Paganini, 2010).

The saddle-point and learning dynamics reviewed above are not directly applicable to the problem considered here because of the information about the game is distributed across the networked scenario and is not centrally available anywhere. The distributed computation of Nash equilibria in noncooperative games, where all players are adversarial, is a challenging problem. The algorithm in (Li and Başar, 1987) relies on all-to-all communication and does not require players to know each other’s payoff functions (which must be strongly convex). In (Frihauf et al., 2012; Stankovic et al., 2012), players are unaware of their own payoff functions but have access to the payoff value of an action once it has been executed. These works design distributed strategies based on extremum seeking techniques to seek the set of Nash equilibria.

**Statement of contributions.** We introduce the problem of distributed convergence to Nash equilibria for two networks engaged in a strategic scenario. The networks aim to either maximize or minimize a common objective function which can be written as a sum of concave-convex functions. Individual agents collaborate with neighbors in their respective network and have partial knowledge of the state of the agents in the other one. Our first contribution is the introduction of an aggregate objective function for each network which depends on the interaction topology through its Laplacian and the characterization of a family of points with a saddle property for the pair of functions. We show the correspondence between these points and the Nash equilibria of the overall game. When the graphs describing the interaction topologies within each network are undirected, the gradients of these aggregate objective functions are distributed. Building on this observation, our second contribution is the synthesis of a consensus-based saddle-point strategy for adversarial networks with undirected topologies. We show that the proposed dynamics is guaranteed to asymptotically converge to the Nash equilibrium for the class of strictly concave-convex and locally Lipschitz objective functions. Our third contribution focuses on the directed case. We show that the transcription of the saddle-point dynamics to directed topologies fails to converge in general. This leads us to propose a generalization of the dynamics, for strongly connected weight-balanced topologies, that incorporates a design parameter. We show that, by appropriately choosing this parameter, the new dynamics asymptotically converges to the Nash equilibrium for the class of strictly concave-convex differentiable objective functions with globally Lipschitz gradients. The technical approach employs notions and results from algebraic graph theory, nonsmooth and convex analysis, set-valued dynamical systems, and game theory. As an intermediate result in our proof strategy for the directed case, we provide a generalization of the known characterization of cocoercivity of concave functions to concave-convex functions.

The results of this paper can be understood as a generalization to competing networks of the results we obtained in (Gharesifard and Cortés, 2012c) for distributed optimization. This generalization is nontrivial because the payoff functions associated to individual agents now also depend on information obtained from the opposing network. This feature gives rise to a hierarchy of saddle-point dynamics whose analysis is technically challenging and requires, among other things, a reformulation of the problem as a constrained zero-sum game, a careful understanding of the coupling between the dynamics of both networks, and the generalization of the notion of cocoercivity to concave-convex functions.

**Organization.** Section 2 contains preliminaries on nonsmooth analysis, set-valued dynamical systems, graph theory, and game theory. In Section 3, we introduce the zero-sum game for two adversarial networks involved in a strategic scenario and introduce two novel aggregate objective functions. Section 4 presents our algorithm design and analysis for distributed convergence to Nash equilibrium when the network topologies are undirected. Section 5 presents our treatment for the directed case. Section 6 gathers our conclusions and ideas for future work. Appendix A contains the generalization to concave-convex functions of the characterization of cocoercivity of concave functions.

### 2 Preliminaries

We start with some notational conventions. Let $\mathbb{R}$, $\mathbb{R}_0$, $\mathbb{Z}$, $\mathbb{Z}_{\geq 1}$ denote the set of real, nonnegative real, integer, and positive integer numbers, respectively. We denote by $|| \cdot ||$ the Euclidean norm on $\mathbb{R}^n$, $d \in \mathbb{Z}_{\geq 1}$ and also use the short-hand notation $1_d = (1, \ldots , 1)^T$ and $0_d = (0, \ldots , 0)^T \in \mathbb{R}^d$. We let $I_d$ denote the identity matrix in $\mathbb{R}^{d \times d}$. For matrices $A \in \mathbb{R}^{d_1 \times d_2}$ and $B \in \mathbb{R}^{e_1 \times e_2}$, $d_1, d_2, e_1, e_2 \in \mathbb{Z}_{\geq 1}$, we let $A \otimes B$ denote their Kronecker product. The function $f : X_1 \times X_2 \to \mathbb{R}$, with $X_1 \subset \mathbb{R}^{d_1}$, $X_2 \subset \mathbb{R}^{d_2}$ closed and convex, is concave-convex if it is concave in its first argument and convex in the second one (Rockafellar, 1997). A point $(x_1^*, x_2^*) \in X_1 \times X_2$ is a saddle point of $f$ if $f(x_1, x_2^*) \leq f(x_1^*, x_2^*) \leq f(x_1^*, x_2)$.
for all \( x_1 \in X_1 \) and \( x_2 \in X_2 \). Finally, a set-valued map \( f : \mathbb{R}^d \rightarrow \mathbb{R}^d \) takes elements of \( \mathbb{R}^d \) to subsets of \( \mathbb{R}^d \).

2.1 Nonsmooth analysis

We recall some notions from nonsmooth analysis (Clarke, 1983). A function \( f : \mathbb{R}^d \rightarrow \mathbb{R} \) is locally Lipschitz at \( x \in \mathbb{R}^d \) if there exists a neighborhood \( U \) of \( x \) and \( C_x \in \mathbb{R}_{\geq 0} \) such that \( |f(y) - f(z)| \leq C_x |y - z| \), for \( y, z \in U \). A solution to this dynamical system is an absolutely continuous curve \( x : [0, T] \rightarrow X \) which satisfies (2) almost everywhere. The set of equilibria of (2) is denoted by \( \psi(\Psi) = \{ x \in X | 0 \in \Psi(x) \} \).

2.2 Set-valued dynamical systems

Let \( \Psi : \mathbb{R}^d \rightarrow \mathbb{R}^d \) be upper semicontinuous with nonempty, compact, and convex values, there exists a solution to (2) from any initial condition.

The LaSalle Invariance Principle for set-valued continuous-time systems is helpful to establish the asymptotic stability properties of systems of the form (2). A set \( W \subset X \) is weakly positively invariant with respect to \( \Psi \) if for any \( x \in W \), there exists \( \tilde{x} \in X \) such that \( \tilde{x} \in \Psi(x) \). The set \( W \) is strongly positively invariant with respect to \( \Psi \) if \( \Psi(x) \subset W \), for all \( x \in W \). Finally, the set-valued Lie derivative of a differentiable function \( V : \mathbb{R}^d \rightarrow \mathbb{R} \) with respect to \( \Psi \) at \( x \in \mathbb{R}^d \) is defined by \( \mathcal{L}_x V(x) = \{ \nu \cdot \nabla V(x) | \nu \in \Psi(x) \} \).

Theorem 2.4 (Set-valued LaSalle Invariance Principle): Let \( W \subset X \) be a strongly positively invariant under (2) and \( V : X \rightarrow \mathbb{R} \) a continuously differentiable function. Suppose the evolutions of (2) are bounded and \( \max \mathcal{L}_x V(x) \leq 0 \) or \( \mathcal{L}_x V(x) = \emptyset \), for all \( x \in W \). If \( S_{\Psi, V} = \{ x \in X | 0 \in \mathcal{L}_x V(x) \} \), then any solution \( x(t) \), \( t \in \mathbb{R}_{\geq 0} \), starting in \( W \) converges to the largest weakly positively invariant set \( M \) contained in \( S_{\Psi, V} \cap W \). When \( M \) is a finite collection of points, then the limit of each solution equals one of them.

2.3 Graph theory

We present some basic notions from algebraic graph theory following the exposition in (Bullo et al., 2009). A directed graph, or simply digraph, is a pair \( G = (\mathcal{V}, E) \), where \( \mathcal{V} \) is a finite set called the vertex set and \( E \subset \mathcal{V} \times \mathcal{V} \) is the edge set. A digraph is undirected if \( (u, v) \in E \) anytime \( (u, v) \in E \). We refer to an undirected digraph as a graph. A path is an ordered sequence of vertices such that any ordered pair of vertices appearing consecutively is an edge of the digraph. A digraph is strongly connected if there is a path between any pair of distinct vertices. For a graph, we refer to this notion simply as connected. A weighted digraph is a triplet \( G = (\mathcal{V}, E, A) \), where \( (\mathcal{V}, E) \) is a digraph and \( A \in \mathbb{R}^{n \times n} \) is the adjacency matrix of \( G \), with the property that \( a_{ij} > 0 \) if \( v_i, v_j \in E \) and \( a_{ij} = 0 \), otherwise. The weighted out-degree matrix \( D_{out} \) is the diagonal matrix defined by \( D_{out} = d_{out} \), for all \( i \in \{1, \ldots, n\} \). The Laplacian matrix is \( L = D_{out} - A \). Note that \( L1_n = 0 \). If \( G \) is strongly connected, then zero is a simple eigenvalue of \( L \). \( G \) is undirected if \( L = L^T \) and weight-balanced if \( d_{out} = d_{out} \), for all \( v \in V \). Equivalently, \( G \) is weight-balanced if and only if \( L1_n = 0 \) if and only if \( L = L^T \) is positive semidefinite. Furthermore, if \( G \) is weight-balanced and strongly connected, then zero is a simple eigenvalue of \( L + L^T \). Note that any undirected graph is weight-balanced.
2.4 Zero-sum games

We recall basic game-theoretic notions following Başar and Olsder (1999). An n-player game is a triplet \( G = (P, X, U) \), where \( P \) is the set of players with \( |P| = n \in \mathbb{Z}_{\geq 2} \), \( X = X_1 \times \ldots \times X_n \), \( X_i \subset \mathbb{R}^{d_i} \) is the set of (pure) strategies of player \( v_i \in P \), \( d_i \in \mathbb{Z}_{\geq 1} \), and \( U = (u_1, \ldots, u_n) \), where \( u_i : X \to \mathbb{R} \) is the payoff function of player \( v_i \), \( i \in \{1, \ldots, n\} \). The game \( G \) is called a zero-sum game if \( \sum_{i=1}^n u_i = 0 \). An outcome \( x^* \in X \) is a (pure) Nash equilibrium of \( G \) if for all \( i \in \{1, \ldots, n\} \) and all \( x_i \in X_i \),

\[
u_i(x^*_i, x^-_{i}) \geq u_i(x_i, x^-_{i}),
\]

where \( x^-_i \) denotes the actions of all players other than \( v_i \).

In this paper, we focus on a class of two-player zero-sum games which have at least one pure Nash equilibrium as the next result states.

**Theorem 2.5 (Minmax theorem):** Let \( X_1 \subset \mathbb{R}^{d_1} \) and \( X_2 \subset \mathbb{R}^{d_2} \), \( d_1, d_2 \in \mathbb{Z}_{\geq 1} \), be nonempty, compact, and convex. If \( u : X_1 \times X_2 \to \mathbb{R} \) is continuous and the sets \( \{x' \in X_1 \mid u(x', y) \geq \alpha \} \) and \( \{y' \in X_2 \mid u(x, y') \leq \alpha \} \) are convex for all \( x \in X_1, y \in X_2 \), and \( \alpha \in \mathbb{R} \), then

\[
\max_{x} \min_{y} u(x, y) = \min_{y} \max_{x} u(x, y).
\]

Theorem 2.5 implies that the game \( G = (\{v_1, v_2\}, X_1 \times X_2, (u, -u)) \) has a pure Nash equilibrium.

3 Problem statement

Consider two networks \( \Sigma_1 \) and \( \Sigma_2 \) composed of agents \( \{v_1, \ldots, v_m\} \) and agents \( \{w_1, \ldots, w_n\} \), respectively. Throughout this paper, \( \Sigma_1 \) and \( \Sigma_2 \) are either connected undirected graphs, c.f. Section 4, or strongly connected weight-balanced digraphs, c.f. Section 5. Since the latter case includes the first one, throughout this section, we assume the latter. The state of \( \Sigma_1 \), denoted \( x_1 \), belongs to \( X_1 \subset \mathbb{R}^{d_1}, d_1 \in \mathbb{Z}_{\geq 1} \). Likewise, the state of \( \Sigma_2 \), denoted \( x_2 \), belongs to \( X_2 \subset \mathbb{R}^{d_2}, d_2 \in \mathbb{Z}_{\geq 1} \). In this paper, we do not get into the details of what these states represent (as a particular case, the network state could correspond to the collection of the states of agents in it). In addition, each agent \( v_i \) in \( \Sigma_1 \) has an estimate \( x_1^i \in \mathbb{R}^{d_1} \) of what the network state is, which may differ from the actual value \( x_1 \). Similarly, each agent \( w_j \) in \( \Sigma_2 \) has an estimate \( x_2^j \in \mathbb{R}^{d_2} \) of what the network state is. Within each network, neighboring agents can share their estimates. Networks can also obtain information about each other. This is modeled by means of a bipartite directed graph \( \Sigma_{eng} \), called engagement graph, with disjoint vertex sets \( \{v_1, \ldots, v_m\} \) and \( \{w_1, \ldots, w_n\} \), where every agent has at least one out-neighbor. According to this model, an agent in \( \Sigma_1 \) obtains information from its out-neighbors in \( \Sigma_{eng} \) about their estimates of the state of \( \Sigma_2 \), and vice versa.

Figure 1 illustrates this concept.

![Fig. 1. Networks Σ1 and Σ2 engaged in a strategic scenario. Both networks are strongly connected and weight-balanced, with weights of 1 on each edge. Edges which correspond to Σeng are dashed.](image)

For each \( i \in \{1, \ldots, n\} \), let \( f_i^1 : X_1 \times X_2 \to \mathbb{R} \) be a locally Lipschitz concave-convex function only available to agent \( v_i \in \Sigma_1 \). Similarly, let \( f_j^2 : X_1 \times X_2 \to \mathbb{R} \) be a locally Lipschitz concave-convex function only available to agent \( w_j \in \Sigma_2 \), \( j \in \{1, \ldots, n\} \). The networks \( \Sigma_1 \) and \( \Sigma_2 \) are engaged in a zero-sum game with payoff function \( U : X_1 \times X_2 \to \mathbb{R} \)

\[
U(x_1, x_2) = \sum_{i=1}^{n_1} f_i^1(x_1, x_2) - \sum_{j=1}^{n_2} f_j^2(x_1, x_2),
\]

where \( \Sigma_1 \) wishes to maximize \( U \), while \( \Sigma_2 \) wishes to minimize it. The objective of the networks is therefore to settle upon a Nash equilibrium, i.e., to solve the following maxmin problem

\[
\max_{x_1 \in X_1, x_2 \in X_2} \min_{x_1, x_2} U(x_1, x_2).
\]
Likewise, the network $\Sigma_2$ does something similar with $m_2$ channels with noise power $y_1$, $n - 1 - m_2$ channels with noise power $y_2$, and one last channel with noise power $C - m_2 y_1 - (n - 1 - m_2) y_2$. Each network is aware of the information received from their neighbors in the information topology. For $i = 1, \ldots, n$, this takes the form

$$f^i(x, y) = \log \left(1 + \frac{\beta x_n}{\sigma_i + y_b}\right),$$

for some $a, b \in \{1, 2\}$. Here $x = (x_1, x_2)$ and $y = (y_1, y_2)$. For $i = n$, it takes instead the form

$$f^n(x, y) = \log \left(1 + \frac{\beta (P - m_1 x_1 - (n - m_1) x_2)}{\sigma_n + C - m_2 y_1 - (n - 1 - m_2) y_2}\right).$$

Note that $\sum_{i=1}^n f^i(x, y)$ is the total capacity of the $n$ communication channels.

### 3.1 Reformulation of the 2-network zero-sum game

In this section, we describe how agents in each network use the information obtained from their neighbors to compute the value of their own objective functions. Based on these estimates, we introduce a reformulation of the $G_{\text{adv-net}} = (\Sigma_1, \Sigma_2, \Sigma_{\text{eng}}, U)$ which is instrumental for establishing some of our results.

Each agent in $\Sigma_1$ has a locally Lipschitz, concave-convex function $\tilde{f}_1^i : \mathbb{R}^{d_1} \times \mathbb{R}^{d_2 n_2} \to \mathbb{R}$ with the properties:

(Extension of own payoff function): for any $x_1 \in \mathbb{R}^{d_1}, x_2 \in \mathbb{R}^{d_2}$,

$$\tilde{f}_1^i(x_1, 1_{n_2} \otimes x_2) = f_1^i(x_1, x_2). \quad (5a)$$

(Distributed over $\Sigma_{\text{eng}}$): there exists $f_1 : \mathbb{R}^{d_1} \times \mathbb{R}^{d_2 |\Sigma_{\text{eng}}(v_i)|} \to \mathbb{R}$ such that, for any $x_1 \in \mathbb{R}^{d_1}, x_2 \in \mathbb{R}^{d_2 n_2}$,

$$\tilde{f}_1^i(x_1, x_2) = f_1^i \left(x_1, \pi_1^i(x_2)\right), \quad (5b)$$

with $\pi_1^i : \mathbb{R}^{d_2 n_2} \to \mathbb{R}^{d_2 |\Sigma_{\text{eng}}(v_i)|}$ the projection of $x_2$ to the values received by $v_i$ from its out-neighbors in $\Sigma_{\text{eng}}$.

Equation (5a) states the fact that, when the estimates of all neighbors of an agent in the opponent’s network agree, its evaluation should coincide with this estimate. Equation (5b) states the fact that agents can only use the information received from their neighbors in the interaction topology to compute their new estimates.

Each agent in $\Sigma_2$ has a function $\tilde{f}_2^j : \mathbb{R}^{d_1 n_1} \times \mathbb{R}^{d_2} \to \mathbb{R}$ with similar properties. The collective payoff functions of the two networks are

$$\tilde{U}_1(x_1, x_2) = \sum_{i=1}^{n_1} \tilde{f}_1^i(x_1, x_2), \quad (6a)$$

$$\tilde{U}_2(x_1, x_2) = \sum_{j=1}^{n_2} \tilde{f}_2^j(x_1, x_2). \quad (6b)$$

In general, the functions $\tilde{U}_1$ and $\tilde{U}_2$ need not be the same. However, $\tilde{U}_1(1_{n_1} \otimes x_1, 1_{n_2} \otimes x_2) = \tilde{U}_2(1_{n_1} \otimes x_1, 1_{n_2} \otimes x_2)$, for any $x_1 \in \mathbb{R}^{d_1}, x_2 \in \mathbb{R}^{d_2}$. When both functions coincide, the next result shows that the original game can be lifted to a (constrained) zero-sum game.

**Lemma 3.2 (Reformulation of the 2-network zero-sum game):** Assume that the individual payoff functions $\{\tilde{f}_1^i\}_{i=1}^{n_1}, \{\tilde{f}_2^j\}_{j=1}^{n_2}$ satisfying (5) are such that the network payoff functions defined in (6) satisfy $\tilde{U}_1 = \tilde{U}_2$, and let $\tilde{U}$ denote this common function. Then, the problem (4) on $\mathbb{R}^{d_1} \times \mathbb{R}^{d_2}$ is equivalent to the following problem on $\mathbb{R}^{n_1 d_1} \times \mathbb{R}^{n_2 d_2}$,

$$\max_{x_1 \in \mathbb{R}^{n_1 d_1}, x_2 \in \mathbb{R}^{n_2 d_2}} \tilde{U}(x_1, x_2),$$

subject to $L_1 x_1 = 0_{n_1 d_1}, \quad L_2 x_2 = 0_{n_2 d_2}, \quad (7)$

with $L_\ell = L_\ell \otimes I_{d_1}$ and $L_\ell$ the Laplacian of $\Sigma_\ell, \ell \in \{1, 2\}$.

**Proof.** The proof follows by noting that (i) $\tilde{U}(1_{n_1} \otimes x_1, 1_{n_2} \otimes x_2) = U(x_1, x_2)$ for all $x_1 \in \mathbb{R}^{d_1}$ and $x_2 \in \mathbb{R}^{d_2}$ and (ii) since $G_1$ and $G_2$ are strongly connected, $L_1 x_1 = 0_{n_1 d_1}$ and $L_2 x_2 = 0_{n_2 d_2}$ if $x_1 = 1_{n_1} \otimes x_1$ and $x_2 = 1_{n_2} \otimes x_2$ for some $x_1 \in \mathbb{R}^{d_1}$ and $x_2 \in \mathbb{R}^{d_2}$.

**Remark 3.3 (Restrictions on extensions):** The assumption of Lemma 3.2 does not hold in general for all sets of extensions satisfying (5a) and (5b). If the interaction topology is one-to-one (i.e., both networks have the same number of agents, the interaction topology is undirected, and each agent in the first network obtains information only from one agent in the opposing network), the natural extensions satisfy the assumption. Example 5.5 later provides yet another instance of a different nature. In general, determining if it is always possible to choose the extensions in such a way that the assumption holds is an open problem.

We denote by $G_{\text{adv-net}} = (\Sigma_1, \Sigma_2, \Sigma_{\text{eng}}, \tilde{U})$ the constrained zero-sum game defined by (7) and refer to this situation by saying that $G_{\text{adv-net}}$ can be lifted to $G_{\text{adv-net}}$. Our objective is to design a coordination algorithm that is implementable with the information that agents in $\Sigma_1$ and $\Sigma_2$ possess and leads them to find a Nash equilibrium of $G_{\text{adv-net}}$, which corresponds to a Nash equilibrium of $G_{\text{adv-net}}$ by Lemma 3.2. Achieving this goal, however, is nontrivial because individual agents, not networks themselves, are the decision makers. From the point of view of agents in each network,
the objective is to agree on the states of both their own network and the other network, and that the resulting states correspond to a Nash equilibrium of $G_{adv-net}$.

The function $\tilde{U}$ is locally Lipschitz and concave-convex. Moreover, from Lemma 2.2, the elements of $\partial_{x_1} \tilde{U}(x_1, x_2)$ are of the form

$$\tilde{g}(x_1, x_2) = (\tilde{g}_1(x_1, x_2), \ldots, \tilde{g}_n(x_1, x_2)) \in \partial_{x_1} \tilde{U}(x_1, x_2),$$

where $\tilde{g}_i(x_1, x_2) \in \partial_2 f_1(x_1, x_2)$, for $i \in \{1, \ldots, n\}$. Note that, because of (5b), we have $\partial_{x_1} f_1(x_1, 1_n \otimes x_2) = \partial_{x_2} f_1(x_1, x_2)$. A similar reasoning can be followed to describe the elements of $\partial_{x_2} \tilde{U}(x_1, x_2)$.

Next, we present a characterization of the Nash equilibria of $G_{adv-net}$, instrumental for proving some of our upcoming results.

**Proposition 3.4 (Characterization of the Nash equilibria of $G_{adv-net}$):** For $\Sigma_1, \Sigma_2$ strongly connected and weight-balanced, define $F_1$ and $F_2$ by

$$F_1(x_1, z_1, x_2) = -\tilde{U}(x_1, x_2) + x_1^T L_1 z_1 + \frac{1}{2} x_1^T L_1 x_1,$$

$$F_2(x_2, z_2, x_1) = \tilde{U}(x_1, x_2) + x_2^T L_2 z_2 + \frac{1}{2} x_2^T L_2 x_2.$$

Then, $F_1$ and $F_2$ are convex in their first argument, linear in their second one, and concave in their third one. Moreover, assume $(x_1^*, z_1^*, x_2^*, z_2^*)$ satisfies the following saddle property for $(F_1, F_2)$: $(x_1^*, z_1^*)$ is a saddle point of $(x_1, z_1) \mapsto F_1(x_1, z_1, x_2^*)$ and $(x_2^*, z_2^*)$ is a saddle point of $(x_2, z_2) \mapsto F_2(x_2, z_2, x_1^*)$.

(i) $(x_1^*, z_1^*, x_2^*, z_2^*)$ satisfies the saddle property for $(F_1, F_2)$ for any $a_1 \in \mathbb{R}^{d_1}$, $a_2 \in \mathbb{R}^{d_2}$, and $x_1 \in \mathbb{R}^{d_1}$, $x_2 \in \mathbb{R}^{d_2}$.

(ii) $(x_1^*, x_2^*)$ is a Nash equilibrium of $G_{adv-net}$.

Furthermore,

(iii) if $(x_1^*, x_2^*)$ is a Nash equilibrium of $G_{adv-net}$ then there exists $z_1^*, z_2^*$ such that $(x_1^*, z_1^*, x_2^*, z_2^*)$ satisfies the saddle property for $(F_1, F_2)$.

**Proof.** The statement (i) is immediate. For (ii), since $(x_1^*, z_1^*, x_2^*, z_2^*)$ satisfies the saddle property, and the networks are strongly connected and weight-balanced, we have $x_1^* = 1_n \otimes x_1^*$, $x_2^* = 1_n \otimes x_2^*$, $x_2^* = 1_n \otimes x_2^*$, $x_2^* = 1_n \otimes x_2^*$, $L_1 z_1^* = -\partial_{x_1} \tilde{U}(x_1^*, x_2^*)$, and $L_2 z_2^* = -\partial_{x_2} \tilde{U}(x_1^*, x_2^*)$. Thus there exist $g_1(x_1, x_2) \in \partial_{x_1} f_1(x_1, x_2)$, $i \in \{1, \ldots, n\}$, and $g_2(x_1, x_2) \in \partial_{x_2} f_2(x_1, x_2)$, $j \in \{1, \ldots, n\}$, such that

$$L_1 z_1^* = (g_1(x_1, x_2), \ldots, g_n(x_1, x_2))^T,$$

$$L_2 z_2^* = (g_1(x_1, x_2), \ldots, g_n(x_1, x_2))^T.$$

Noting that, for $\ell \in \{1, 2\}$, $(1_{n_\ell} \otimes l_{d_\ell}) L_{\ell} = (1_{n_\ell} \otimes l_{d_\ell}) (L_{\ell} \otimes l_{d_\ell}) = 1_{n_\ell} L_{\ell} \otimes l_{d_\ell} = 0_{d_\ell \times d_\ell}$, we deduce that

$$\sum_{i=1}^{n_1} g_i(x_1, x_2) = 0_{d_1},$$

and

$$\sum_{j=1}^{n_2} g_j(x_1, x_2) = 0_{d_2},$$

i.e., $(x_1^*, x_2^*)$ is a Nash equilibrium. Finally for proving (iii), note that $x_1^* = 1_n \otimes x_1^*$ and $x_2^* = 1_n \otimes x_2^*$. The result follows then from the fact that $0 \in \partial_{x_1} \tilde{U}(x_1^*, x_2^*)$ implies that there exists $x_1^* \in \mathbb{R}^{n_1 d_1}$ and $z_2^* \in \mathbb{R}^{n_2 d_2}$ with $L_1 z_1^* \in \partial_{x_1} \tilde{U}(x_1^*, x_2^*)$ and $L_2 z_2^* \in -\partial_{x_2} \tilde{U}(x_1^*, x_2^*)$.

**4 Distributed convergence to Nash equilibria for undirected topologies**

In this section, we introduce a distributed dynamics which solves (7) when $\Sigma_1$ and $\Sigma_2$ are undirected. In particular, we design gradient dynamics to find points with the saddle property for $(F_1, F_2)$ prescribed by Proposition 3.4. Consider the set-valued dynamics $\Psi_{\text{Nash-undir}}: (\mathbb{R}^{d_1 n_1})^2 \times (\mathbb{R}^{d_2 n_2})^2 \mapsto ((\mathbb{R}^{d_1 n_1})^2 \times (\mathbb{R}^{d_2 n_2})^2$,

$$\dot{x}_1 + L_1 x_1 + L_1 z_1 \in \partial_{x_1} \tilde{U}(x_1, x_2),$$

$$\dot{z}_1 = L_1 x_1,$$

$$\dot{x}_2 + L_2 x_2 + L_2 z_2 \in -\partial_{x_2} \tilde{U}(x_1, x_2),$$

$$\dot{z}_2 = L_2 x_2,$$

where $x_1, z_1 \in \mathbb{R}^{n_1 d_1}, \ell \in \{1, 2\}$. Note that (8a)-(8b) and (8c)-(8d) correspond to saddle-point dynamics of $F_1$ in $(x_1, z_1)$ and $F_2$ in $(x_2, z_2)$, respectively. Local solutions to this dynamics exist by virtue of Lemmas 2.1 and 2.3. We characterize next its asymptotic convergence properties.

**Theorem 4.1 (Distributed convergence to Nash equilibrium for undirected networks):** Consider the zero-sum game $G_{\text{adv-net}} = (\Sigma_1, \Sigma_2, \Sigma_{\text{eng}}, U)$, with $\Sigma_1$ and $\Sigma_2$ connected undirected graphs, $X_1 \subset \mathbb{R}^{d_1}$, $X_2 \subset \mathbb{R}^{d_2}$, compact and convex, and $U : X_1 \times X_2 \to \mathbb{R}$ strictly concave-convex and locally Lipschitz. Assume $G_{\text{adv-net}}$ can be lifted to $G_{\text{adv-net}}$. Then, the projection onto the first and third components of the solutions of (8) asymptotically converge to agreement on the Nash equilibrium of $G_{\text{adv-net}}$.

**Proof.** Throughout this proof, since property (5b) holds, without loss of generality and for simplicity of notation, we assume that agents in $\Sigma_1$ have access to $x_2$ and, similarly, agents in $\Sigma_2$ have access to $x_1$. By Theorem 2.5, a solution to (6) exists. By the strict concavity-convexity properties, this solution is, in fact, unique. Let us denote this solution by $x_1^* = 1_n \otimes x_1^*$ and $x_2^* = 1_n \otimes x_2^*$. By Proposition 3.4(iii), there exists $z_1^*$ and $z_2^*$ such that $(x_1^*, z_1^*, x_2^*, z_2^*) \in \operatorname{Eq}(\Psi_{\text{Nash-undir}})$. First, note that given any initial condition $(x_1^0, z_1^0, x_2^0, z_2^0) \in (\mathbb{R}^{n_1 d_1})^2 \times (\mathbb{R}^{n_2 d_2})^2$, the set

$$W_{x_1^0, x_2^0} = \{(x_1, z_1, x_2, z_2) | (1_{n_\ell} \otimes l_{d_\ell}) x_1 = (1_{n_\ell} \otimes l_{d_\ell}) x_1^0, \ell \in \{1, 2\}\}$$

is a Nash equilibrium.
is strongly positively invariant under (8). Consider the function $V : (\mathbb{R}^{d_1n_1})^2 \times (\mathbb{R}^{d_2n_2})^2 \to \mathbb{R}_{\geq 0}$ defined by

$$V(x_1, z_1, x_2, z_2) = \frac{1}{2} (x_1 - x_1^*)^T (x_1 - x_1^*) + \frac{1}{2} (z_1 - z_1^*)^T (z_1 - z_1^*) + \frac{1}{2} (x_2 - x_2^*)^T (x_2 - x_2^*) + \frac{1}{2} (z_2 - z_2^*)^T (z_2 - z_2^*).$$

The function $V$ is smooth. Next, we examine its set-valued Lie derivative along $\Psi_{\text{Nash-undir}}$. Let $\xi \in \mathcal{L}^{\Psi_{\text{Nash-undir}}} V(x_1, z_1, x_2, z_2)$. By definition, there exists $v \in \Psi_{\text{Nash-undir}}(x_1, z_1, x_2, z_2)$, given by

$$v = ( - L_1 x_1 - L_1 z_1 + g_1(x_1, x_2), - L_2 x_2 - L_2 z_2 - g_2(x_2, x_1, z_2), L_1 x_1, L_2 x_2 ),$$

where $g_1(x_1, x_2) \in \partial_{x_1} U(x_1, x_2)$ and $g_2(x_2, x_1) \in \partial_{x_1} U(x_1, x_2)$, such that

$$\xi = v \cdot \nabla V(x_1, z_1, x_2, z_2) = (x_1 - x_1^*)^T (-L_1 x_1 - L_1 z_1 + g_1(x_1, x_2)) + (x_2 - x_2^*)^T (-L_2 x_2 - L_2 z_2 - g_2(x_2, x_1, z_2)) + (z_1 - z_1^*)^T L_1 x_1 + (z_2 - z_2^*)^T L_2 x_2.$$

Note that $-L_1 x_1 - L_1 z_1 + g_1(x_1, x_2) \in -\partial_{x_1} F_1(x_1, z_1, x_2)$, $L_1 x_1 \in \partial_{x_1} F_1(x_1, z_1, x_2) - L_2 x_2 - L_2 z_2 - g_2(x_2, x_1, z_2) \in -\partial_{x_2} F_2(x_2, z_2, x_1)$, and $L_2 x_2 \in \partial_{x_1} F_2(x_2, z_2, x_1)$. Using the first-order convexity property of $F_1$ and $F_2$ in their first two arguments, one gets

$$\xi \leq -\bar{U}(x_1^*, x_2) + \bar{U}(x_1, x_2^*) - z_1^* L_1 x_1 - \frac{1}{2} x_1^* L_1 x_1 - z_2^* L_2 x_2 - \frac{1}{2} x_2^* L_2 x_2.$$

By rearranging, we thus have

$$\xi \leq -F_2(x_2, z_2^*, x_1^*) - F_1(x_1, z_1^*, x_2^*).$$

Next, since $F_2(x_1^*, z_1^*, x_2^*) + F_1(x_2^*, z_2^*, x_1^*) = 0$, we have

$$\xi \leq -F_1(x_1^*, z_1^*, x_1^*) + F_1(x_1, z_1^*, x_2^*) + F_2(x_2, z_2^*, x_1^*) - F_2(x_2, z_2^*, x_1^*),$$

yielding that $\xi \leq 0$. As a result,

$$\max \mathcal{L}^{\Psi_{\text{Nash-undir}}} V(x_1, z_1, x_2, z_2) \leq 0.$$

As a by-product, we conclude that the trajectories of (8) are bounded. By virtue of the set-valued version of the LaSalle Invariance Principle, cf. Theorem 2.4, any trajectory of (8) starting from an initial condition $(x_1^0, z_1^0, x_2^0, z_2^0)$ converges to the largest positively invariant set $M$ in $\mathcal{S}_{\Psi_{\text{Nash-undir}}} V \cap V^{-1}(\leq V(x_1^0, z_1^0, x_2^0, z_2^0))$. Let $(x_1, z_1, x_2, z_2) \in M$. Because $M \subset \mathcal{S}_{\Psi_{\text{Nash-undir}}} V$, then $F_1(x_1^*, z_1^*, x_2^*) - F_1(x_1, z_1^*, x_2^*) = 0$, i.e.,

$$-\bar{U}(x_1^*, x_2^*) + \bar{U}(x_1, x_2^*) - x_1^* L_1 x_1 - \frac{1}{2} x_1^* L_1 x_1 = 0.$$

Define now $G_1 : \mathbb{R}^{n_1d_1} \times \mathbb{R}^{n_1d_1} \times \mathbb{R}^{n_2d_2} \to \mathbb{R}$ by

$$G_1(x_1, z_1, x_2) = F_1(x_1, z_1, x_2) - \frac{1}{2} x_1^* L_1 x_1.$$ As a result, $G_1(x_1^0, z_1^0, x_2^0) \leq 0$, or equivalently, $-\bar{U}(x_1^*, x_2^*) + \bar{U}(x_1, x_2^*) - x_1^* L_1 x_1 \leq 0$. Combining this with (10), we have that $L_1 z_1 = 0$ and $-\bar{U}(x_1^*, x_2^*) + \bar{U}(x_1, x_2^*) = 0$. Since $\bar{U}$ is strictly concave in its first argument $x_1 = x_1^*$. A similar argument establishes that $x_2 = x_2^*$. Using now the fact that $M$ is weakly positively invariant, one can deduce that $L_2 z_2 = 0$, for $\ell \in \{1, 2\}$, and thus $(x_1, x_2, z_2, x_2) \in \mathcal{E}(\Psi_{\text{Nash-undir}})$.

5 Distributed convergence to Nash equilibria for directed topologies

Interestingly, the saddle-point dynamics (8) fails to converge when transcribed to the directed network setting. This observation is a consequence of the following result, which studies the stability of the linearization of the dynamics (8), when the payoff functions have no contribution to the linear part.

Lemma 5.1 (Necessary condition for the convergence of (8) on digraphs): Let $\Sigma_d$ be strongly connected and $f^i_\ell = 0$, $i \in \{1, \ldots, n_\ell\}$, for $\ell \in \{1, 2\}$. Then, the set of network agreement configurations $\mathcal{S}_{\text{agree}} = \{(1_1 \otimes x_1, 1_1 \otimes z_1, 1_2 \otimes x_2, 1_2 \otimes z_2) \in (\mathbb{R}^{n_1d_1})^2 \times (\mathbb{R}^{n_2d_2})^2 \mid x_\ell, z_\ell \in \mathbb{R}^{d_\ell}, \ell \in \{1, 2\}\}$, is stable under (8) iff, for any nonzero eigenvalue $\lambda$ of the Laplacian $L_\ell$, $\ell \in \{1, 2\}$, one has $\sqrt{3}|\text{Im}(\lambda)| \leq \text{Re}(\lambda)$.

Proof. In this case, (8) is linear with matrix

$$\begin{pmatrix} \begin{pmatrix} -1 & -1 \\ 1 & 0 \end{pmatrix} & \otimes L_1 \\ 0 & \begin{pmatrix} -1 & -1 \\ 1 & 0 \end{pmatrix} & \otimes L_2 \end{pmatrix}$$

and has $\mathcal{S}_{\text{agree}}$ as equilibria. The eigenvalues of (11) are of the form $\lambda_\ell \left( \frac{-1}{2} \pm \frac{\sqrt{3}i}{2} \right)$, with $\lambda_\ell$ an eigenvalue of $L_\ell$, for $\ell \in \{1, 2\}$ (since the eigenvalues of a Kronecker product are the product of the eigenvalues of the corresponding matrices). Since $L_\ell = L_\ell \otimes I_{d_\ell}$, each eigenvalue of $L_\ell$ is an eigenvalue of $L_\ell$. The result follows by noting that $\text{Re}(\lambda_\ell \left( \frac{-1}{2} \pm \frac{\sqrt{3}i}{2} \right)) = \frac{1}{2}(\pm \sqrt{3}\text{Im}(\lambda_\ell) - \text{Re}(\lambda_\ell)).$
It is not difficult to construct examples of strictly concave-convex functions that have zero contribution to the linearization of (8) around the solution. Therefore, such systems cannot be convergent if they fail the necessary condition identified in Lemma 5.1. The counterexample provided in our recent paper (Gharesifard and Cortés, 2012c) of strongly connected, weight-balanced digraphs that do not meet the stability criterion of Lemma 5.1 is therefore valid in this context too.

From here on, we assume that the payoff functions are differentiable. We elaborate on the reasons for this assumption in Remark 5.3 later. Motivated by the observation made in Lemma 5.1, we introduce a parameter $\alpha \in \mathbb{R}_{>0}$ in the dynamics of (8) as

\begin{align}
\dot{x}_1 + \alpha L_1 x_1 + L_1 z_1 &= \nabla \tilde{U}(x_1, x_2), \\
\dot{z}_1 &= L_1 x_1, \\
\dot{x}_2 + \alpha L_2 x_2 + L_2 z_2 &= -\nabla \tilde{U}(x_1, x_2), \\
\dot{z}_2 &= L_2 x_2.
\end{align}

We next show that a suitable choice of $\alpha$ makes the dynamics convergent to the Nash equilibrium.

**Theorem 5.2 (Distributed convergence to Nash equilibria for directed networks):** Consider the zero-sum game $G_{\text{adv-net}} = (\Sigma_1, \Sigma_2, \Psi_{\text{Nash-dir}}, U)$, with $\Sigma_1$ and $\Sigma_2$ strongly connected and weight-balanced digraphs, $X_1 \subseteq \mathbb{R}^{d_1}$, $X_2 \subseteq \mathbb{R}^{d_2}$ compact and convex, and $U : X_1 \times X_2 \rightarrow \mathbb{R}$ strictly concave-convex and differentiable with globally Lipschitz gradient. Assume $G_{\text{adv-net}}$ can be lifted to $G_{\text{adv-net}}^\ast$ such that $\tilde{U}$ is differentiable and has a globally Lipschitz gradient. Define $h : \mathbb{R}_{>0} \rightarrow \mathbb{R}$ by

\begin{equation}
h(r) = \frac{1}{2} \Lambda_s^\min \left( \sqrt{\frac{r^4 + 3r^2 + 2}{r}} - 4 - \frac{r^4 + 3r^2 + 2}{r} \right) + \frac{Kr^2}{(1 + r^2)},
\end{equation}

where $\Lambda_s^\min = \min_{t=1,2} \{ \Lambda_s(L_t^+ + \frac{1}{L_t^-}) \}$, $\Lambda_s(\cdot)$ denotes the smallest non-zero eigenvalue and $K \in \mathbb{R}_{>0}$ is the Lipschitz constant of the gradient of $U$. Then there exists $\beta^* \in \mathbb{R}_{>0}$ with $h(\beta^*) = 0$ such that for all $0 < \beta < \beta^*$, the projection onto the first and third components of the solutions of (12) with $\alpha = \frac{\beta^2 + 2}{\beta}$ asymptotically converge to agreement on the Nash equilibrium of $G_{\text{adv-net}}$.

**Proof.** Similarly to the proof of Theorem 4.1, we assume, without loss of generality, that agents in $\Sigma_1$ have access to $x_2$ and agents in $\Sigma_2$ to $x_1$. For convenience, we denote the dynamics described in (12) by $\Psi_{\text{Nash-dir}} : (\mathbb{R}^{d_1 n_1})^2 \times (\mathbb{R}^{d_2 n_2})^2 \rightarrow (\mathbb{R}^{d_1 n_1})^2 \times (\mathbb{R}^{d_2 n_2})^2$. Let $(x_1^0, x_2^0, z_1^0, z_2^0)$ be an arbitrary initial condition. Note that the set $W_{x_1^0, x_2^0}$ defined by (9) is invariant under the evolutions of (12). By an argument similar to the one in the proof of Theorem 4.1, there exists a unique solution to (7), which we denote by $x_1^* = 1_{n_1} \otimes x_1^*$ and $x_2^* = 1_{n_2} \otimes x_2^*$. By Proposition 3.4(i), there exists $(x_1^1, z_1^1, x_2^1, z_2^1) \in \Theta(\Psi_{\text{Nash-dir}}) \cap W_{x_1^0, x_2^0}$. Consider the function $V : (\mathbb{R}^{d_1 n_1})^2 \times (\mathbb{R}^{d_2 n_2})^2 \rightarrow \mathbb{R}_{\geq 0}$,

\begin{equation}
V(x_1, z_1, x_2, z_2) = \frac{1}{2} (x_1 - x_1^*)^T (x_1 - x_1^*) + \frac{1}{2} (x_2 - x_2^*)^T (x_2 - x_2^*) + \frac{1}{2} (y(x_1, z_1) - y(x_1^*, z_1^*))^T (y(x_1, z_1) - y(x_1^*, z_1^*)) + \frac{1}{2} (y(x_2, z_2) - y(x_2^*, z_2^*))^T (y(x_2, z_2) - y(x_2^*, z_2^*))
\end{equation}

where $y(x_1, z_1) = \beta x_1 + z_1$, $\ell \in \{1, 2\}$, and $\beta \in \mathbb{R}_{>0}$ satisfies $\beta^2 - \alpha \beta + 2 = 0$. This function is quadratic, hence smooth. Next, we consider $\xi = \lim_{\rightarrow \Psi_{\text{Nash-dir}}} V(x_1, z_1, x_2, z_2)$ given by

\begin{equation}
\xi = (-\alpha L_1 x_1 - L_1 z_1 + \nabla \tilde{U}(x_1, x_2), L_1 x_1, -\alpha L_2 x_2 - L_2 z_2 - \nabla \tilde{U}(x_1, x_2), L_2 x_2) \cdot \nabla V(x_1, z_1, x_2, z_2).
\end{equation}

After some manipulation, one can show that

\begin{equation}
\xi = \sum_{\ell=1}^2 \frac{1}{2} (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T A_\ell (x_\ell, y(x_\ell, z_\ell)) + \sum_{\ell=1}^2 \frac{1}{2} (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T L_{\ell\ell} - \beta L_{\ell\ell}.
\end{equation}

This equation can be written as

\begin{equation}
\xi = \sum_{\ell=1}^2 \frac{1}{2} (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T Q_\ell (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))
\end{equation}

\begin{equation}
+ \sum_{\ell=1}^2 (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T A_\ell (x_\ell^*, y(x_\ell^*, z_\ell^*))
\end{equation}

\begin{equation}
+ \sum_{\ell=1}^2 (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T L_{\ell\ell} - \beta L_{\ell\ell},
\end{equation}

\begin{equation}
+ \sum_{\ell=1}^2 (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T L_{\ell\ell} - \beta L_{\ell\ell},
\end{equation}

\begin{equation}
+ \sum_{\ell=1}^2 (x_\ell - x_\ell^*, y(x_\ell, z_\ell) - y(x_\ell^*, z_\ell^*))^T L_{\ell\ell} - \beta L_{\ell\ell}.
\end{equation}
where $Q_{\ell}$, $\ell \in \{1, 2\}$, is given by

$$Q_{\ell} = (L_{\ell} + L_{\ell}^T) \otimes \begin{pmatrix} -\left(\frac{\beta^2 + 2}{\beta} - \beta \right) & -1 \\ -1 & -\beta \end{pmatrix}.$$ (14)

Note that, we have

$$A_1(x_1^*, y(x_1^*; z_1)) = -(L_1 y(x_1^*; z_1), \beta L_1 y(x_1^*; z_1))$$
$$= (\nabla x_1 \tilde{U}(x_1, x_2), \beta \nabla x_1 \tilde{U}(x_1^*, x_2))$$
$$= (\nabla x_2 \tilde{U}(x_1, x_2), \beta \nabla x_2 \tilde{U}(x_1, x_2)).$$

Thus, after substituting for $y(x_1^*, z_1)$, we have

$$\xi = \sum_{\ell=1}^{2} \frac{1}{2} (\xi_{\ell} - x_1^*, z_1 - z_1^*)^T \tilde{Q}_\ell (x_1 - x_1^*, z_1 - z_1^*) + (1 + \beta^2) (x_1 - x_1^*)^T \nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2)$$
$$- (1 + \beta^2) (x_2 - x_2^*)^T \nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2)$$
$$+ \beta (z_1 - z_1^*)^T (\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2))$$
$$- \beta (z_2 - z_2^*)^T (\nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2)),$$ (15)

where

$$\tilde{Q}_\ell = \begin{pmatrix} -\beta^3 - \left(\frac{\beta^2 + 2}{\beta} - \beta \right) - (1 + \beta^2) \\ -(1 + \beta^2) & -\beta \end{pmatrix} \otimes (L_{\ell} + L_{\ell}^T),$$

for $\ell \in \{1, 2\}$. Each eigenvalue of $\tilde{Q}_\ell$ is of the form

$$\eta_\ell = \lambda_\ell - (\beta^4 + 3 \beta^2 + 2) \pm \sqrt{(\beta^4 + 3 \beta^2 + 2)^2 - 4 \beta^2},$$ (16)

where $\lambda_\ell$ is an eigenvalue of $L_{\ell} + L_{\ell}^T$, $\ell \in \{1, 2\}$. Using now Theorem A.1 twice, one for $(x_1, x_2)$, $(x_1^*, x_2)$, and another one for $(x_1^*, x_2)$, $(x_1, x_2)$, we have

$$(x_1 - x_1^*)^T (\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2)) \leq$$
$$- \frac{1}{K} \left( ||\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2)||^2 \right),$$

$$(x_2 - x_2^*)^T (\nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2)) \leq$$
$$- \frac{1}{K} \left( ||\nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2)||^2 \right),$$

where $K \in \mathbb{R}_{>0}$ is the Lipschitz constant of $\nabla \tilde{U}$. We thus conclude that

$$\xi \leq \sum_{\ell=1}^{2} \frac{1}{2} (x_1 - x_1^*, z_1 - z_1^*)^T \tilde{Q}_\ell (x_1 - x_1^*, z_1 - z_1^*)$$
$$- \frac{1}{K} \left( ||\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2)||^2 \right)$$
$$+ ||\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2)||^2$$
$$+ \beta (z_1 - z_1^*)^T (\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2))$$
$$- \beta (z_2 - z_2^*)^T (\nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2)).$$

One can write this inequality as displayed in (17), where

$$X = (x_1 - x_1^*, z_1 - z_1^*, x_2 - x_2^*, z_2 - z_2^*)$$
$$\nabla x_1 \tilde{U}(x_1, x_2) - \nabla x_1 \tilde{U}(x_1^*, x_2),$$
$$\nabla x_2 \tilde{U}(x_1, x_2) - \nabla x_2 \tilde{U}(x_1^*, x_2).$$

Since $(x_1^*, z_1^*, x_2^*, z_2^*) \in W_{\tilde{U}}$, we have $(1 \otimes I_{d_1})(z_1^* - z_1^*) = 0_{d_1}$, $\ell \in \{1, 2\}$, and hence it is enough to establish that $Q$ is negative semidefinite on the subspace $W = \{(v_1, v_2, v_3, v_4, v_5, v_6) \in (\mathbb{R}^{n_1 d_1})^2 \times (\mathbb{R}^{n_2 d_2})^2 \times \mathbb{R}^{n_1 d_1} \times \mathbb{R}^{n_2 d_2} | (1 \otimes I_{d_1})v_2 = 0_{d_1}, (1 \otimes I_{d_2})v_4 = 0_{n_2}\}$. Using the fact that $-\frac{1}{K}(1 + \beta^2)I_{n_1 d_1}$ is invertible, for $\ell \in \{1, 2\}$, we can express $Q$ as

$$Q =$$

$$\begin{pmatrix} Q_1 & 0 & 0 \\ 0 & Q_2 & 0 \\ 0 & 0 & -\frac{1}{K}(1 + \beta^2)I_{n_1 d_1} \end{pmatrix}$$
$$= N D$$

where

$$Q_{\ell} = \tilde{Q}_{\ell} + \frac{K \beta^2}{1 + \beta^2} \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \ell \in \{1, 2\},$$

and

$$N = \begin{pmatrix} I_{n_1 d_1} & 0 & 0 & 0 & 0 & 0 \\ 0 & I_{n_1 d_1} & 0 & -\frac{K}{1 + \beta^2} I_{n_1 d_1} & 0 \\ 0 & 0 & I_{n_2 d_2} & 0 & 0 \\ 0 & 0 & 0 & I_{n_1 d_1} & 0 \\ 0 & 0 & 0 & 0 & I_{n_2 d_2} \end{pmatrix}.$$

Noting that $W$ is invariant under $N^T$ (i.e., $N^T W = W$), all we need to check is that the matrix $D$ is negative semidefinite on $W$. Clearly,

$$\begin{pmatrix} -\frac{1}{K}(1 + \beta^2)I_{n_1 d_1} & 0 \\ 0 & -\frac{1}{K}(1 + \beta^2)I_{n_2 d_2} \end{pmatrix}$$
is negative definite. On the other hand, for \( \ell \in \{1, 2\} \), on \((\mathbb{R}^{n_d l})^2\), 0 is an eigenvalue of \( \bar{Q}_\ell \) with multiplicity 2\(d\ell \) and eigenspace generated by vectors of the form \((1_{n_l} \otimes a, 0)\) and \((0, 1_{n_l} \otimes b)\), with \(a, b \in \mathbb{R}^{d\ell} \). However, on \(((v_1, v_2) \in (\mathbb{R}^{n_d l})^2 \mid (1_{n_l} \otimes d\ell) v_2 = 0_{d\ell})\), 0 is an eigenvalue of \( \bar{Q}_\ell \) with multiplicity \(d\ell \) and eigenspace generated by the vectors of the form \((1_{n_l} \otimes a, 0)\). Moreover, on \(((v_1, v_2) \in (\mathbb{R}^{n_d l})^2 \mid (1_{n_l} \otimes d\ell) v_2 = 0_{d\ell})\), the eigenvalues of \( \frac{K\beta^2}{(1+\beta^2)} \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} \) are \( \frac{K\beta^2}{(1+\beta^2)} \) with multiplicity \(n_d \ell - d\ell \) and 0 with multiplicity \(n_d d\ell \). Therefore, using Weyl's theorem (Horn and Johnson, 1985, Theorem 4.3.7), we deduce that the nonzero eigenvalues of the sum \( \bar{Q}_\ell \) are upper bounded by \( \Lambda_* (\bar{Q}_\ell) + \frac{K\beta^2}{(1+\beta^2)} \). Thus, the eigenvalues of \( \bar{Q} = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} \) are upper bounded by \( \min_{\ell=1,2} \Lambda_* (\bar{Q}_\ell) + \frac{K\beta^2}{(1+\beta^2)} \). From (16) and the definition of \( h \) in (13), we conclude that the nonzero eigenvalues of \( \bar{Q} \) are upper bounded by \( h(\beta) \). It remains to show that there exists \( \beta^* \in \mathbb{R}_{>0} \) with \( h(\beta^*) = 0 \) such that for all \( 0 < \beta < \beta^* \) we have \( h(\beta) < 0 \). For \( r > 0 \) small enough, \( h(r) < 0 \) holds. Furthermore, \( \lim_{r \to \infty} h(r) = K > 0 \). Hence, using the Mean Value Theorem, we deduce the existence of \( \beta^* \). Therefore we conclude that \( \mathcal{L}_{\Psi_{\text{Nash}}: \text{dir}} (x_1, z_1, x_2, z_2) \leq 0 \). As a by-product, the trajectories of (8) are bounded. Consequently, all assumptions of the LaSalle Invariance Principle, cf. Theorem 2.4, are satisfied. This result then implies that any trajectory of (8) starting from an initial condition \( (x_1^0, z_1^0, x_2^0, z_2^0) \) converges to the largest invariant set \( M \) in \( S_{\Psi_{\text{Nash}}: \text{dir}} \cap \mathcal{W}_{z^0, x^0} \). Note that if \( (x_1, z_1, x_2, z_2) \in S_{\Psi_{\text{Nash}}: \text{dir}} \cap \mathcal{W}_{z^0, x^0} \), then \( N^T X \in \ker(Q) \times \{0\} \). From the discussion above, we know \( \ker(Q) \) is generated by vectors of the form \((1_{n_l} \otimes a_1, 0, 0, 0), (0, 0, 1_{n_l} \otimes a_2, 0), a_1 \in \mathbb{R}^{d\ell}, j \in \{1, 2\} \), and hence \( x_1 = x_1^* + 1_{n_l} \otimes a_1, z_1 = x_1^* \). Using the strict concavity-convexity, this then implies that \( x_1 = x_1^* \). Finally, for \( (x_1^*, z_1, x_2, z_2) \in M \), using the positive invariance of \( M \), one deduces that \( (x_1^*, z_1, x_2, z_2) \in \text{Eq}(\Psi_{\text{Nash}}: \text{dir}) \).

Remark 5.3 (Assumptions on payoff function): Two observations are in order regarding the assumptions in Theorem 5.2 on the payoff function. First, the assumption that the payoff function has a globally Lipschitz gradient is too restrictive given that, since the state spaces are compact, standard boundedness conditions on the gradient imply the globally Lipschitz condition. Second, we restrict our attention to differentiable payoff functions because locally Lipschitz functions with globally Lipschitz generalized gradients are in fact differentiable, see (Gharesifard and Cortés, 2012c, Proposition A.1).

Remark 5.4 (Comparison with best-response dynamics): Using the gradient flow has the advantage of avoiding the cumbersome computation of the best-response map. This, however, does not come for free. There are concave-convex functions for which the (distributed) gradient flow dynamics, unlike the best-response dynamics, fails to converge to the saddle point, see (Feijer and Paganini, 2010) for an example.

We finish this section with an example.

Example 5.5 (Distributed adversarial selection of signal and noise power via (12)): Recall the communication scenario described in Remark 3.1. Consider 5 channels, \( \{c_{1h}, c_{2h}, c_{3h}, c_{4h}, c_{5h}\} \), for which the network \( \Sigma_1 \) has decided that \( \{c_{1h}, c_{3h}\} \) have signal power \( x_1 \) and \( \{c_{2h}, c_{4h}\} \) have signal power \( x_2 \). Channel \( c_{5h} \) has its signal power determined to satisfy the budget constraint \( P \in \mathbb{R}_{>0} \), i.e., \( P - 2x_1 - 2x_2 \). Similarly, the network \( \Sigma_2 \) has decided that \( c_{1h} \) has noise power \( y_1 \), \( \{c_{2h}, c_{3h}, c_{4h}\} \) have noise power \( y_2 \), and \( c_{5h} \) has noise power \( C - y_1 - 3y_2 \) to meet the budget constraint \( C \in \mathbb{R}_{>0} \). We let \( x = (x^1, x^2, x^3, x^4, x^5) \) and \( y = (y^1, y^2, y^3, y^4, y^5) \), where \( x^i \in [0, P]^2 \) and \( y^i \in [0, C]^2 \) for each \( i \in \{1, \ldots, 5\} \).

The networks \( \Sigma_1 \) and \( \Sigma_2 \), which are weight-balanced and strongly connected, and the engagement topology \( \Sigma_{\text{eng}} \) are shown in Figure 2. Note that, according to this topology, each agent can observe the power employed by its adversary in its channel and, additionally, the agents in channel 2 can obtain information about the estimates of the opponent in channel 4 and vice versa. The payoff functions of the agents are given in Remark 3.1, where for simplicity we take \( \sigma_1 = \sigma_1 \), for \( i \in \{1, 3, 5\} \), and \( \sigma_1 = \sigma_2 \), for \( i \in \{2, 4\} \), with \( \sigma_1, \sigma_2 \in \mathbb{R}_{>0} \).

This example fits into the approach described in Section 3.1 by considering the following extended payoff
functions:

\[
\begin{align*}
\tilde{f}_1^1(x^1, y) &= \log(1 + \frac{\beta x_1^1}{\sigma_1 + y_1}), \\
\tilde{f}_1^2(x^2, y) &= \frac{1}{3} \log(1 + \frac{\beta x_2^2}{\sigma_2 + y_2}) + \frac{2}{3} \log(1 + \frac{\beta x_2^2}{\sigma_2 + y_2}), \\
\tilde{f}_1^3(x^3, y) &= \log(1 + \frac{\beta x_3^3}{\sigma_1 + y_3}), \\
\tilde{f}_1^4(x^4, y) &= \frac{1}{3} \log(1 + \frac{\beta x_4^4}{\sigma_2 + y_2}) + \frac{2}{3} \log(1 + \frac{\beta x_4^4}{\sigma_2 + y_2}), \\
\tilde{f}_1^5(x^5, y) &= \log\left(1 + \frac{\beta(P - 2x_1^5 - 2x_2^5)}{\sigma_1 + C - y_1^5 - 3y_2^5}\right), \\
\tilde{f}_1^1(x, y^1) &= \tilde{f}_1^1(x^1, y), \\
\tilde{f}_1^2(x, y^2) &= \frac{2}{3} \log(1 + \frac{\beta x_2^2}{\sigma_2 + y_2}) + \frac{1}{3} \log(1 + \frac{\beta x_2^2}{\sigma_2 + y_2}), \\
\tilde{f}_1^3(x, y^3) &= \frac{1}{3} \log(1 + \frac{\beta x_3^3}{\sigma_2 + y_2}) + \frac{2}{3} \log(1 + \frac{\beta x_3^3}{\sigma_2 + y_2}), \\
\tilde{f}_1^4(x, y^4) &= \frac{1}{3} \log(1 + \frac{\beta x_4^4}{\sigma_2 + y_2}), \\
\tilde{f}_1^5(x, y^5) &= \tilde{f}_1^5(x^5, y).
\end{align*}
\]

Note that these functions are strictly concave and thus the zero-sum game defined has a unique saddle point on the set \([0, P]^2 \times [0, C]^2\). These functions satisfy (5) and \(U_1 = U_2\). Figure 3 shows the convergence of the dynamics (12) to the Nash equilibrium of the resulting 2-network zero-sum game.

6 Conclusions and future work

We have considered a class of strategic scenarios in which two networks of agents are involved in a zero-sum game. The networks aim to either maximize or minimize a common objective function. Individual agents collaborate with neighbors in their respective network and have partial knowledge of the state of the agents in the other one. We have introduced two aggregate objective functions, one per network, identified a family of points with a special saddle property for this pair of functions, and established their correspondence between the Nash equilibrium of the overall game. When the individual networks are undirected, we have proposed a distributed saddlepoint dynamics that is implementable by each network via local interactions. We have shown that, for a class of strictly concave-convex and locally Lipschitz objective functions, the proposed dynamics is guaranteed to converge to the Nash equilibrium. We have also shown that this saddle-point dynamics fails to converge for directed networks, even when they are strongly connected and weight-balanced. Motivated by this fact, we have introduced a generalization that incorporates a design parameter. We have shown that this dynamics converges to the Nash equilibrium for strictly concave-convex and differentiable objective functions with globally Lipschitz gradients for appropriate parameter choices. An interesting venue of research is determining whether it is always possible to choose the extensions of the individual payoff functions in such a way that the lifted objective functions coincide. Future work will also include relaxing the assumptions of strict concavity-convexity and differentiability of the payoff functions, and the globally Lipschitz condition on their gradients, extending our results to dynamic interaction topologies and non-zero sum games, and exploring the application to various areas, including collaborative resource allocation in the presence of adversaries, strategic social networks, collective bargaining, and collaborative pursuit-evasion.

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References


Fig. 3. Execution of (12) over the networked strategic scenario described in Example 5.5, with $\beta = 8$, $\sigma_1 = 1$, $\sigma_2 = 4$, $P = 6$, and $C = 4$. (a) and (b) show the evolution of the agent’s estimates of the state of networks $\Sigma_1$ and $\Sigma_2$, respectively, and (c) shows the value of the Lyapunov function. Here, $\alpha = 3$ in (12) and initially, $x^0 = ((1, 0.5), (0.5, 1), (0.5, 0.5), (0.5, 1), (0.5, 1), (0.5, 1))^T$, $z^0_1 = 0_{10}$, $y^0 = ((1, 0.5), (0.5, 1), (0.5, 1), (0.5, 0.5), (1, 0.5))^T$ and $z^0_2 = 0_{10}$. The equilibrium $(x^*, z^*_1, y^*, z^*_2)$ is $x^* = (1.3371, 1.0315)^T \otimes 1_5$, $y^* = (1.5027, 0.3366)^T \otimes 1_5$, $z^*_1 = (0.7508, 0.5084, 0.1447, 0.5084, 0.1447, -0.1271, -0.5201, -0.1271, -0.5201, -0.7626)^T$ and $z^*_2 = (0.1079, -0.0987, -0.0002, 0.2237, 0.0358, 0.2875, -0.0360, 0.0087, -0.1076, -0.4213)$.


A Appendix

The following result can be understood as a generalization of the characterization of cocoercivity of concave functions (Golshtein and Tretyakov, 1996, Lemma 6.7).

**Theorem A.1 (Concave-convex differentiable functions with globally Lipschitz gradients):** Let $f : \mathbb{R}^d \times \mathbb{R}^d \to \mathbb{R}$ be a concave-convex differentiable function with globally Lipschitz gradient (with Lipschitz constant $K \in \mathbb{R}_{>0}$). For $(x, y), (x', y') \in \mathbb{R}^d \times \mathbb{R}^d$,

$$(x - x')^T(\nabla_x f(x, y) - \nabla_x f(x', y')) + (y - y')^T(\nabla_y f(x, y') - \nabla_y f(x, y)) \leq -\frac{1}{2K}(\|\nabla_x f(x, y') - \nabla_x f(x', y')\|^2 + \|\nabla_y f(x', y') - \nabla_y f(x, y')\|^2),$$

where we have used the fact that $\sup_{x \in \mathbb{R}^d} \tilde{f}(x, y') = \inf_{y \in \mathbb{R}^d} \tilde{f}(x', y) = \tilde{f}(x', y') = 0$. Next, by definition of $\tilde{f}$,

$$\tilde{f}(x, y') = f(x, y') - f(x', y') - (x - x')^T \nabla_x f(x', y'),$$

$$\tilde{f}(x', y) = f(x', y) - f(x', y') - (y - y')^T \nabla_y f(x', y').$$

Using (A.2), we deduce that

$$f(x, y') - f(x', y) - (x - x')^T \nabla_x f(x', y') + (y - y')^T \nabla_y f(x', y') \leq -\frac{1}{2K}(\|\nabla_x f(x, y') - \nabla_x f(x', y')\|^2 + \|\nabla_y f(x', y') - \nabla_y f(x, y')\|^2). \quad (A.3)$$

The claim now follows by adding together (A.3) and the inequality that results by interchanging $(x, y)$ and $(x', y')$ in (A.3). □

**Proof.** We start by noting that, for a concave function $j : \mathbb{R}^d \to \mathbb{R}$ with globally Lipschitz gradient, the following inequality holds, see (Golshtein and Tretyakov, 1996, Equation 6.64),

$$j(x) \leq j^* - \frac{1}{2M} \|\nabla j(x)\|^2, \quad (A.1)$$

where $j^* = \sup_{x \in \mathbb{R}^d} j(x)$ and $M$ is the Lipschitz constant of $\nabla j$. Given $(x', y') \in \mathbb{R}^d \times \mathbb{R}^d$, define the map $\tilde{f} : \mathbb{R}^d \times \mathbb{R}^d \to \mathbb{R}$ by

$$\tilde{f}(x, y) = f(x, y) - f(x', y) - (x - x')^T \nabla_x f(x', y) + f(x, y) - f(x, y') + (y - y')^T \nabla_y f(x', y').$$

Since the gradient of $f$ is Lipschitz, the function $\tilde{f}$ is differentiable almost everywhere. Thus, almost everywhere, we have

$$\nabla_x \tilde{f}(x, y) = \nabla_x f(x, y) - \nabla_x f(x', y) + \nabla_x f(x, y) - \nabla_x f(x, y') - (y - y')^T \nabla_x \nabla_y f(x, y'),$$

$$\nabla_y \tilde{f}(x, y) = \nabla_y f(x, y) - \nabla_y f(x', y) + \nabla_y f(x, y) - \nabla_y f(x, y') - (x - x')^T \nabla_y \nabla_x f(x', y).$$

In particular, note that $\nabla_x \tilde{f}(x', y') = \nabla_y \tilde{f}(x', y') = 0$. Since $x \mapsto \tilde{f}(x, y')$ and $y \mapsto \tilde{f}(x', y)$ are concave and convex functions, respectively, we can use (A.1) to deduce

$$\tilde{f}(x, y') \leq -\frac{1}{2K}\|\nabla_x f(x, y') - \nabla_x f(x', y')\|^2, \quad (A.2a)$$

$$-\tilde{f}(x', y) \leq -\frac{1}{2K}\|\nabla_y f(x', y) - \nabla_y f(x', y')\|^2, \quad (A.2b)$$

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