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<p>Principal Author: 2Lt. Anthony Alvin Illig</p>	<p>Other Author(s): Maj. August Roesener, Maj. Shane Knighton, Maj. Shane Hall</p>	
<p>Principal Author's Organization: Air Force Institute of Technology Complete mailing address: 346 Evergreen Ave. Niceville, FL 32578</p>	<p><input checked="" type="checkbox"/> <i>Anthony Illig</i> Principal Author's Signature: ↑ Date: <i>29 Apr 08</i> Phone: <i>(850) 883-5010</i> FAX: <i>(850) 882-6277</i> Email: <i>anthony.illig@afotec.af.mil</i></p>	
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<p>Complete mailing address: <i>2950 HOBSON WAY WPAFB OH 45433-7765</i></p>	<p>Phone: _____ FAX: _____</p>	



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A NETWORK FLOW APPROACH TO THE INITIAL SKILLS TRAINING SCHEDULING PROBLEM

Lt. Tony Illig

Maj. August Roesener

Maj. Shane Knighton

Maj. Shane Hall

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Overview



- Problem Description
- Motivation
- Network Flow Approach
- Why Network Flow Works
- Results
- Future Research

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Problem Description



- Initial Skills Training Scheduling Problem (ISTSP):
 - After commissioning, new Air Force officers typically require training to be certified to perform their duties.
 - Commissioning sources: USAFA, ROTC, OTS
 - All officers must attend ASBC. (Mandated by CSAF)
 - Following ASBC, many officers continue training to be certified
 - Pilots
 - Combat Systems Officers (CSOs)
 - Air Battle Managers
 - Intel Officers
 - Space and Missile Officers
 - Air Field Operations Officers

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Problem Description



- Initial Skills Training Scheduling Problem (ISTSP):
 - Course sequences are rigid and vary by AFSC
 - Class sizes and intervals vary
 - Course length is fixed
 - Class blends must be met
 - AETC desires that certain blends are met according to commissioning source and AFSC (in ASBC).

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Motivation



- Average 2nd Lt Pay:
 - \$124 / day (2007 dollars)
- Number of Annually Commissioned 2nd Lts:
 - ~ 4000
- Estimated savings of 1 day:
 - \$496,000
- Current average number of down days:
 - ~ 180
- Huge Potential for Savings!!!!

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Network Flow Approach



- General personnel scheduling techniques are computational intensive
- ISTSP has structure which can be exploited
- Developed network flow approach (NFISTSP)
 - How it works
 - Commissioning sources are source nodes.
 - Create nodes to represent each class.
 - Connect the nodes with arcs if it is possible to go from the starting node's class to the ending nodes class.
 - The cost on the arc is the number of days between the classes.
 - Use minimum cost flow to solve.

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NFISTSP Approach



- What it accomplishes
 - Enforces blends:
 - Source of commissioning (SOC) blends enforced by creating 3 separate nodes at each class (one for each SOC) then capacitating arcs accordingly.
 - AFSC blends are done in a similar way at ASBC.
 - Does not allow skipping.

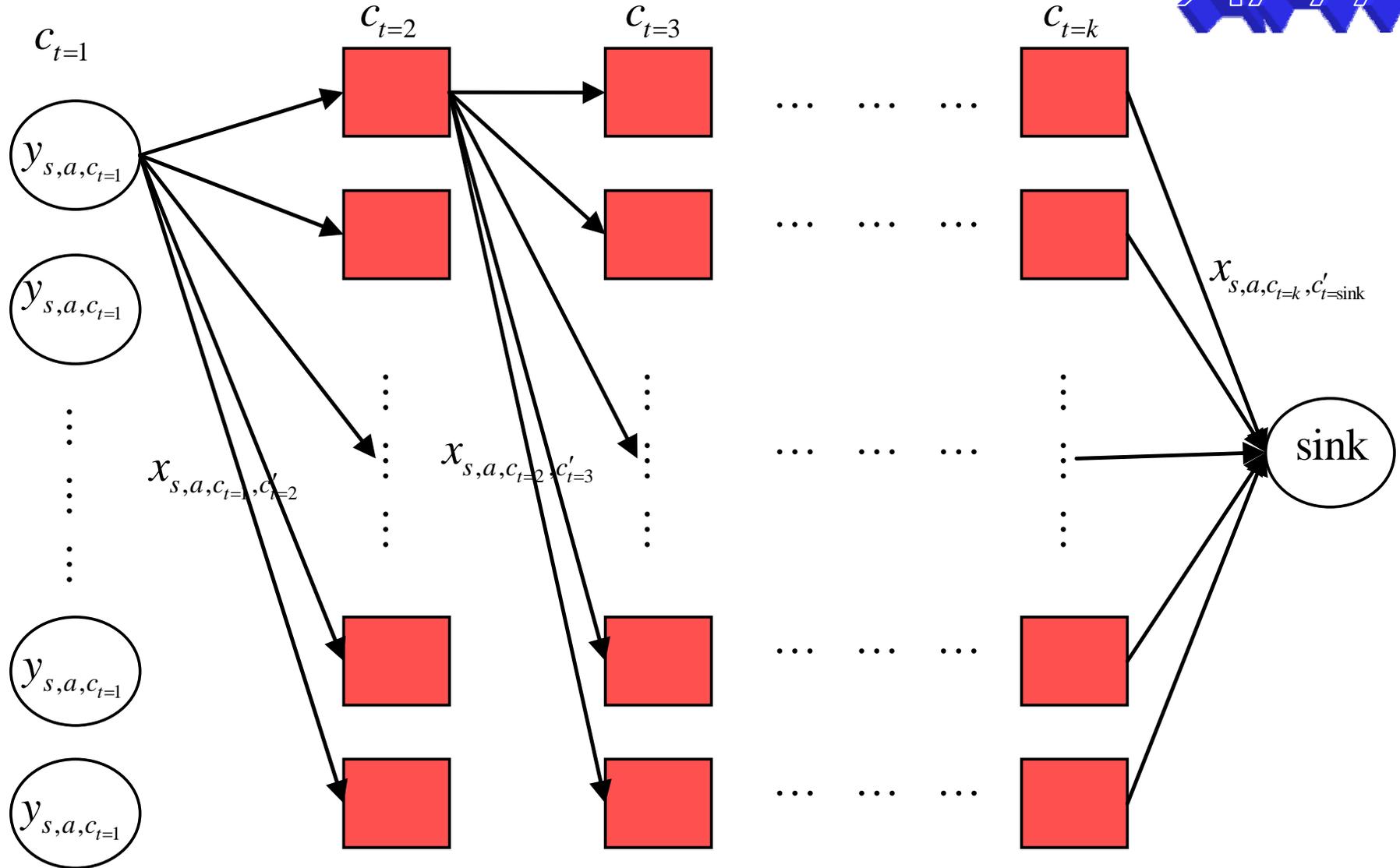
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NFISTSP Graph



AFIT





NFISTSP Graph



Sum all nodes for class UB

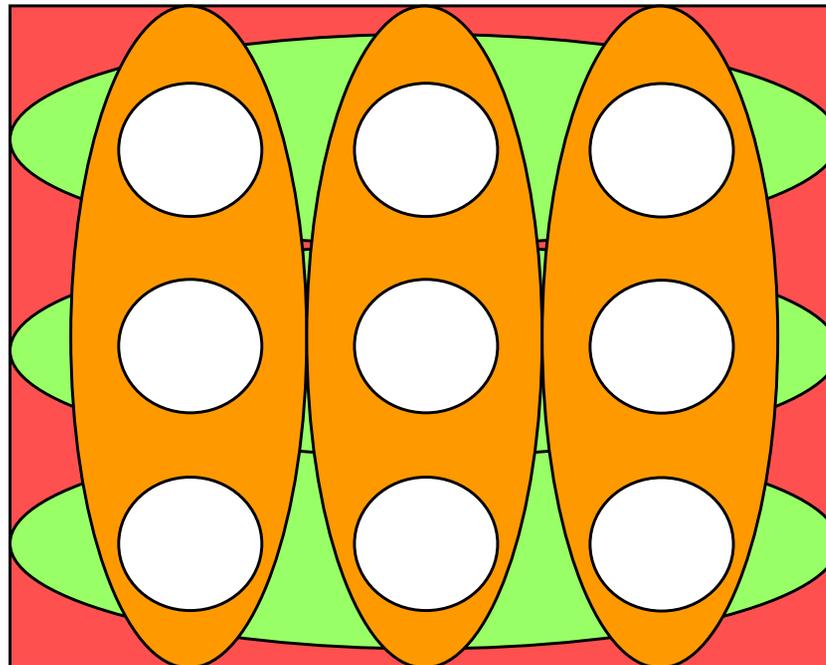
$$\sum_{\forall s \in S} \sum_{a \in A'} \sum_{c \in C_t} x_{s,a,c,c'} \leq UB_{c'}$$

Sum all nodes for AFSC UB

$$\sum_{\forall s \in S} \sum_{c \in C_t} x_{s,a,c,c'} \leq UB_{a,c'}$$

Sum all nodes for SOC UB

$$\sum_{a \in A'} \sum_{c \in C_t} x_{s,a,c,c'} \leq UB_{s,c'}$$





Why does it solve quickly?



- Using this formulation allows large personnel scheduling problems to solve in seconds
 - Why?
- Total Unimodularity (TU)
 - Definition: \mathbf{A} is totally unimodular iff every square sub matrix in \mathbf{A} has a determinant of -1, 0 or 1.
 - Implications: Suppose \mathbf{A} is the constraint matrix for a integer programming (IP) problem. If \mathbf{A} is TU and the RHS is all integer then the LP relaxation of the IP will yield integer results.

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Test Cases/Scenarios



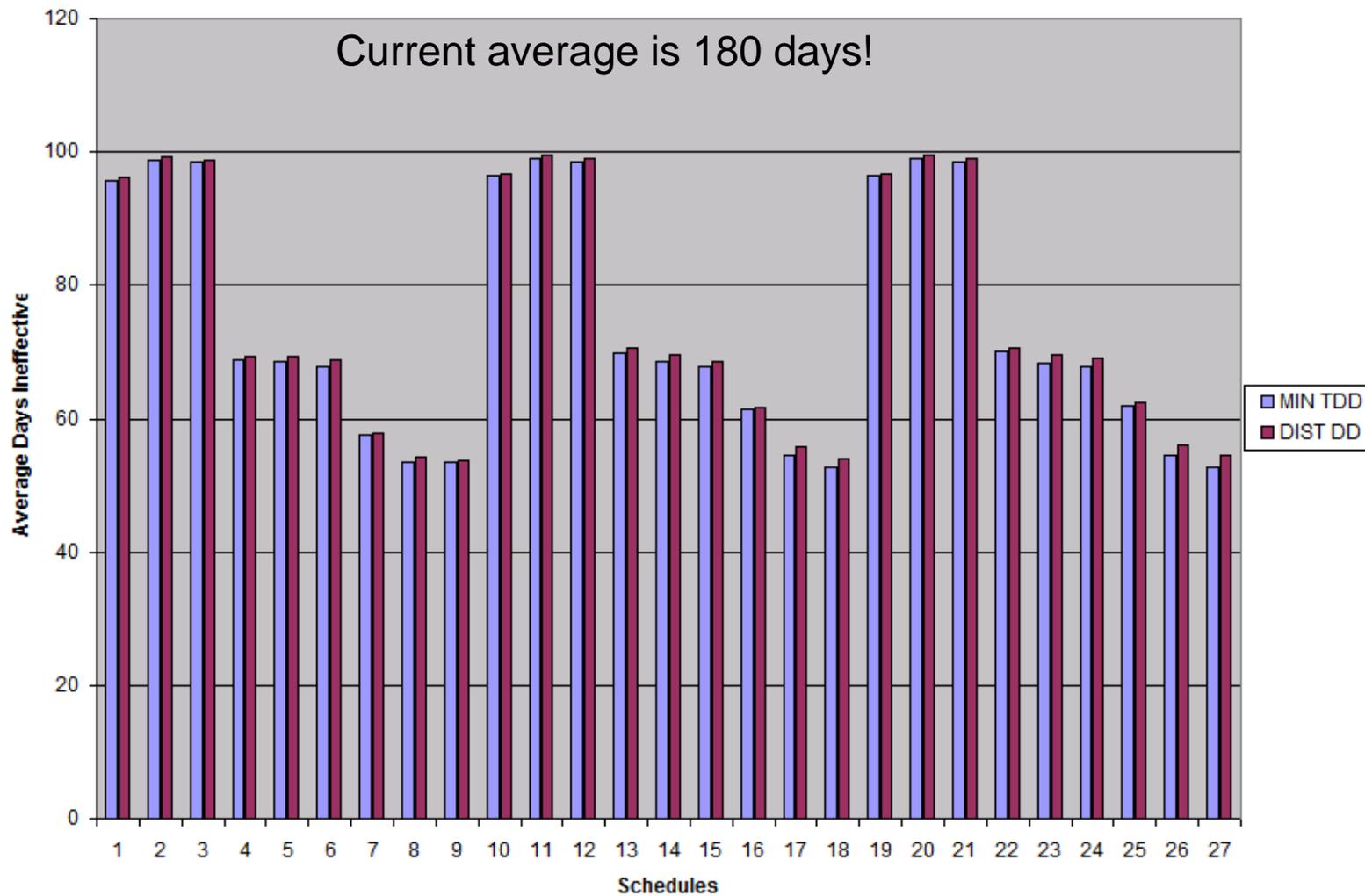
- Min USAFA Leave
 - 0, 30, 60 Days 3 Levels
- Max ROTC Delay
 - 0, 180, 365 Days 3 Levels
- Class Blend Levels
 - 50%, 75%, No Forced Blend 3 Levels
- Arc Weighting Schemes (OF)
 - Minimize TDD, Distribute DD 2 Levels

Total Scenarios: 54

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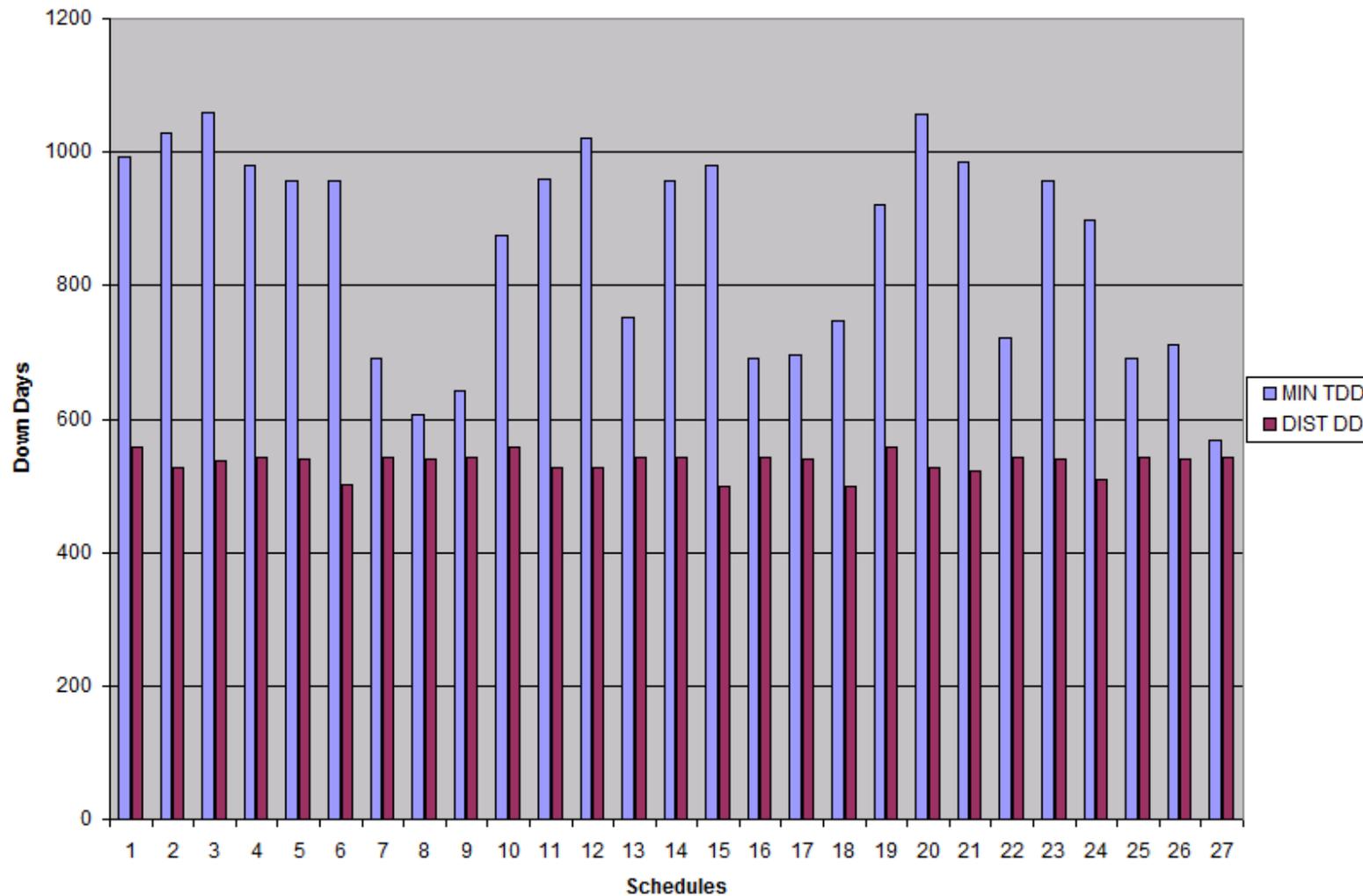


Minimize TDD vs Distribute DD





Minimize TDD vs Distribute DD

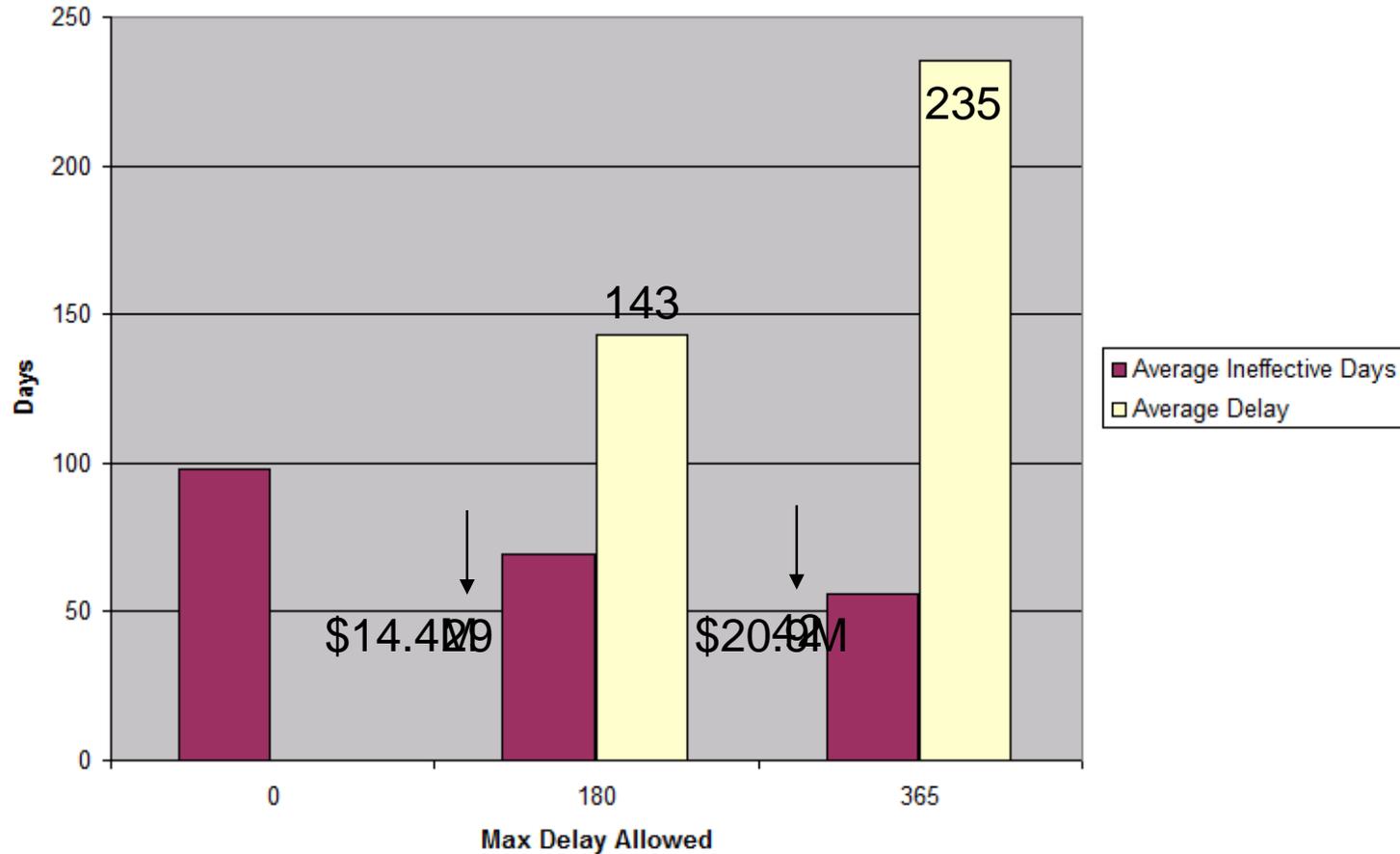




Impact of ROTC Delay

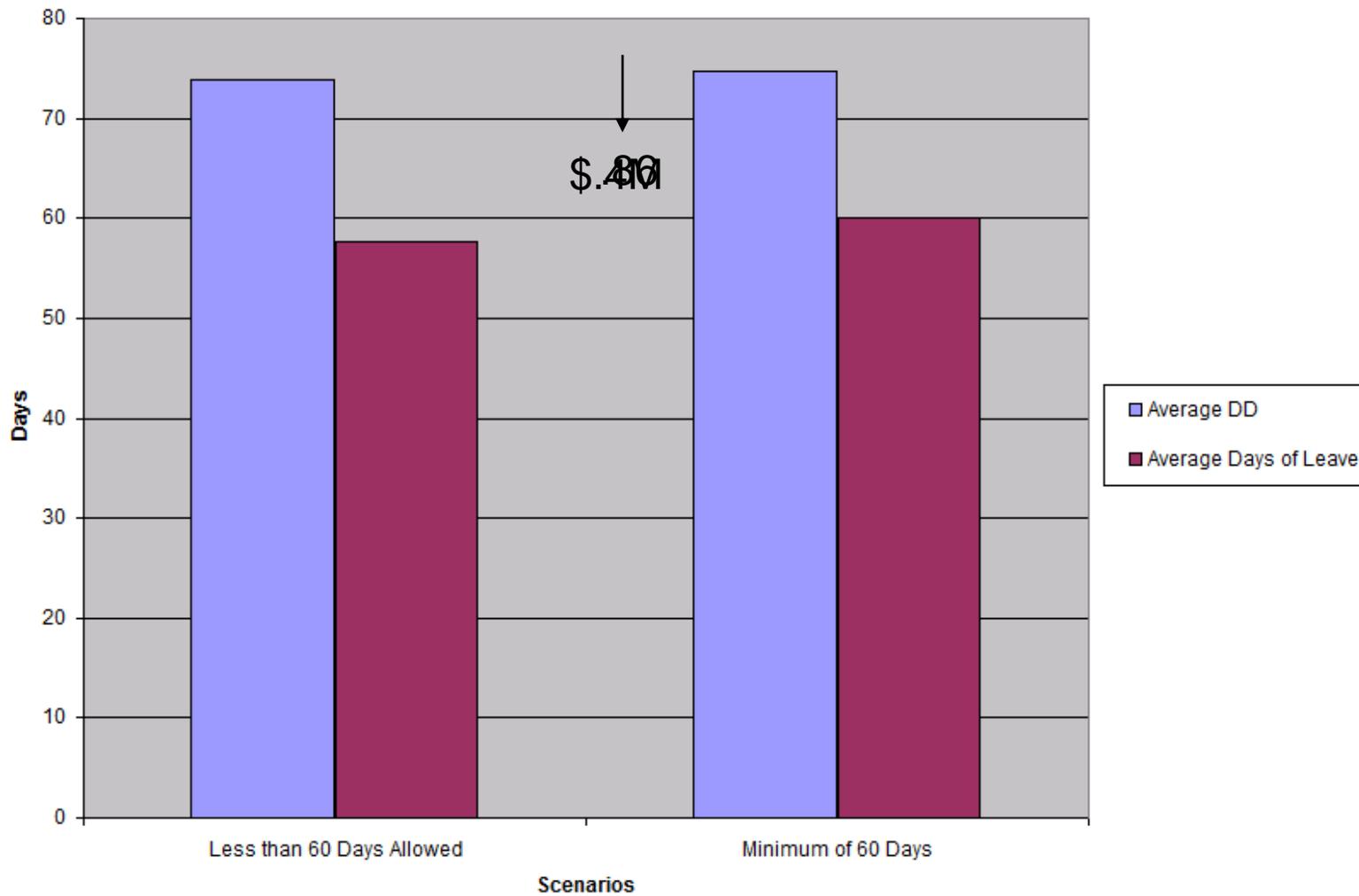


Average Days Ineffective vs ROTC Delay





Impact of USAFA Leave

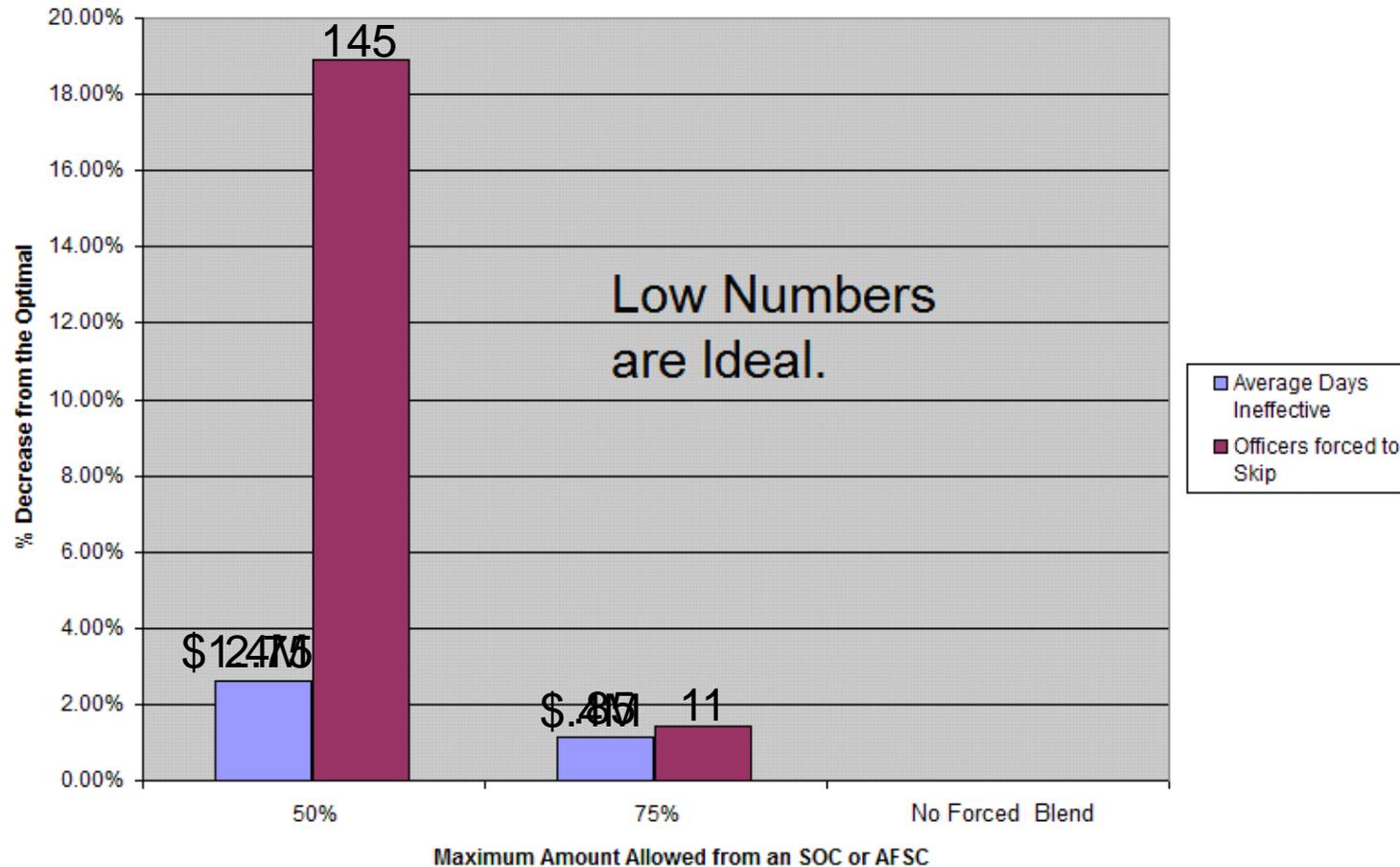




Impact of Blend Levels



Effect of Blend Levels





Computational Conclusions



- The savings (dollars or ineffective days) from reducing the minimum USAFA Leave is not significant.
- ROTC delay extension results in a sizable effect.
- Restricting blends significantly impacts total number of down days; partial blends occur naturally.

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Robustness of Formulation



- Theoretically, any timeline can be scheduled.
- This formulation handles heterogeneous workforces well.
 - In this example we had 2 heterogeneous levels (AFSC and SOC).
- Can handle any number of types and levels efficiently.

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Future Research



- Use of Decision Analysis
- Use new optimization package in SAS/OR
- Use model to schedule optimal class sizes and class start dates
- Use probabilities and stochastic analysis to include drop out rates

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Questions?

Thanks for your time ...

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NFISTSP Mathematical Formulation



- Sets

\mathbf{S} = set of all commissioning sources = $\{s_1, s_2, \dots, s_n\}$

\mathbf{A} = set of all AFSCs = $\{a_1, a_2, \dots, a_m\}$

\mathbf{T} = set of all training courses = $\{t_1, t_2, \dots, t_k\}$

\mathbf{C}_t = set of all classes of training course t = $\{c_{t_1}, c_{t_2}, \dots, c_{t_i}\}$

- Variables

x_{s,a,c_t,c'_t} = number of officers from commissioning source s
in AFSC a going from class c_t to class c'_t



NFISTSP Mathematical Formulation



- Parameters

d_{c_t, c'_t} = number of days between class c_t and class c'_t

y_{s, a, c_t} = number of officers from commissioning source s
in AFSC a graduating from class c_t

- Objective Function

Minimize

$$Z = \sum_{\forall t \in T} \sum_{\forall c'_t \in C_{t+1}} \sum_{\forall c_t \in C_t} d_{c_t, c'_t} \sum_{\forall a \in A} \sum_{\forall s \in S} x_{s, a, c_t, c'_t}$$



NFISTSP Mathematical Formulation



- Subject to

$$y_{s,a,c_1} = \sum_{\forall c' \in C_2} x_{s,a,c_1,c'} \quad \forall s \in S, a \in A, c_1 \in C_1$$

$$\sum_{\forall c_t \in C_t} x_{s,a,c_t,c'_{t+1}} = \sum_{\forall c'_{t+2} \in C_{t+2}} x_{s,a,c_{t+1},c'_{t+2}} \quad \forall s \in S, a \in A, c'_{t+1} = c_{t+1} \in C_{t+1}, t = 1, 2, \dots, n-1$$

$$\sum_{\forall s \in S} \sum_{\forall a \in A} \sum_{\forall c_k \in C_k} x_{s,a,c_k,c'} = \sum_{\forall s \in S} \sum_{\forall a \in A} \sum_{\forall c_t \in C_1} y_{s,a,c_t} \quad c' = \text{"sink"}$$

$$\sum_{\forall s \in S} \sum_{a \in A'} \sum_{\forall c_t \in C_t} x_{s,a,c_t,c'_{t+1}} \leq UB_{c'_{t+1}} \quad \forall c'_{t+1} \in C_{t+1}, t = 1, 2, \dots, n-1$$

$$\sum_{\forall s \in S} \sum_{\forall c_t \in C_t} x_{s,a,c_t,c'} \leq UB_{a,c'_2} \quad \forall a \in A, c'_2 \in C_2$$

$$\sum_{\forall a \in A'} \sum_{\forall c_t \in C_t} x_{s,a,c_t,c'_{t+1}} \leq UB_{s,c'_{t+1}} \quad \forall s \in S, c' \in C_{t+1}, t = 1, 2, \dots, n-1$$



Properties of Total Unimodularity



- Properties of TU that will be useful:
 - If A is TU then A^T is TU
 - If A is TU then (A, I) is TU
 - Row and Column swaps in A do not affect TU.
 - The Node-Arc incidence matrix of a Network Flow problem is always TU
 - If A is TU and B is obtained by removing a row or column of A , then B is TU
- The first four are from Bazaraa, Jarvis and Serali



Proving Properties of Total Unimodularity



- If A is TU and B is obtained by removing a row or column of A, then B is TU

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \cdots & a_{1j} \\ a_{21} & a_{22} & a_{23} & a_{24} & \cdots & a_{2j} \\ a_{31} & a_{32} & a_{33} & a_{34} & \cdots & a_{3j} \\ a_{41} & a_{42} & a_{43} & a_{44} & \cdots & a_{4j} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{i1} & a_{i2} & a_{i3} & a_{i4} & \cdots & a_{ij} \end{bmatrix}$$



Total Unimodularity



- A look at the constraint matrix for a network (the node incidence matrix)

Arcs
↓

Nodes →

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \cdots & a_{1j} \\ 1 & -1 & 0 & 0 & \cdots & 0 \\ -1 & 1 & 1 & 0 & \cdots & 0 \\ 0 & 0 & -1 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{i1} & a_{i2} & a_{i3} & a_{i4} & \cdots & a_{ij} \end{bmatrix}$$



Total Unimodularity



- Another sufficient condition for TU is:
 - For every $J \subseteq N = \{1, \dots, n\}$, there exists a partition J_1, J_2 of J such that

$$\left| \sum_{j \in J_1} a_{ij} - \sum_{j \in J_2} a_{ij} \right| \leq 1 \quad \forall i = 1, \dots, m$$



Total Unimodularity



- The side constraints for this problem are of this form:

$$C = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \cdots & a_{1j} \\ a_{21} & a_{22} & a_{23} & a_{24} & \cdots & a_{2j} \\ a_{31} & a_{32} & a_{33} & a_{34} & \cdots & a_{3j} \\ a_{41} & a_{42} & a_{43} & a_{44} & \cdots & a_{4j} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{i1} & a_{i2} & a_{i3} & a_{i4} & \cdots & a_{ij} \end{bmatrix}$$

$$C \in \{0,1\}^{m \times n}$$



Total Unimodularity



- The side constraints for this problem are of this form:

$$C' = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \cdots & a_{1j} \\ a_{21} & a_{22} & a_{23} & a_{24} & \cdots & a_{2j} \\ a_{31} & a_{32} & a_{33} & a_{34} & \cdots & a_{3j} \\ a_{41} & a_{42} & a_{43} & a_{44} & \cdots & a_{4j} \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{i1} & a_{i2} & a_{i3} & a_{i4} & \cdots & a_{ij} \end{bmatrix}$$

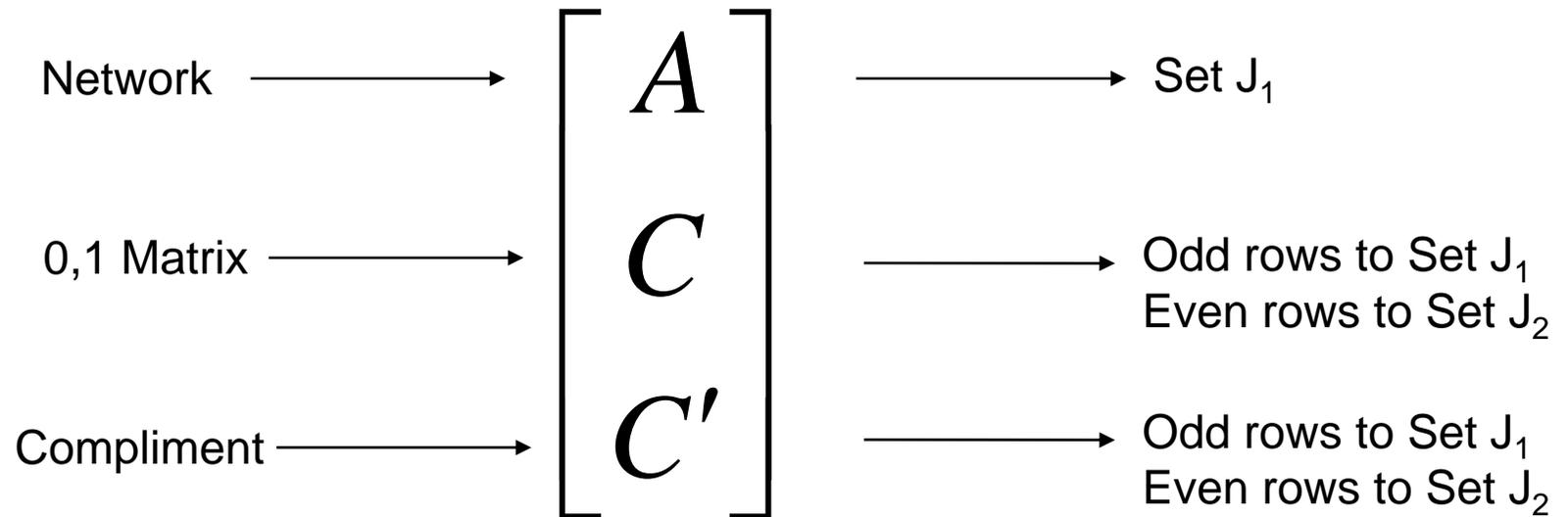
$$C' \in \{1, 0\}^{m \times n}$$



Total Unimodularity



- Consider the Matrix:





Total Unimodularity



- Another sufficient condition for TU is:
 - For every $J \subseteq N = \{1, \dots, n\}$, there exists a partition J_1, J_2 of J such that

$$\left| \sum_{j \in J_1} a_{ij} - \sum_{j \in J_2} a_{ij} \right| \leq 1 \quad \forall i = 1, \dots, m$$

\uparrow \uparrow

$n/2$ or $(n+1)/2$ $n/2$ or $(n-1)/2$



Total Unimodularity



- Therefore, this matrix is TU

$$\begin{bmatrix} A \\ C \\ \del{C'} \end{bmatrix}$$