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The research has focused on a technology for "mobile information management". Underlying this technology is a mathematical foundation enabling the use of formal methods in developing and reasoning about the construction of mobile information management components and their use in database integration and transformation. The salient features of our approach are: use of unmaterialized views; dynamic integration of data consumers and data sources, using mobile query processes; data interface specifications, based on XML schemas; specifications for transformations, based on XML query languages; the development of formal methods, focusing on query and constraint reformulation.

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(CVAL TANNEN)

Enclosure 3
1 Personnel

In addition to the PIs, this contract has partially supported PhD students Alin Deutsch, Wenfei Fan, Carmem Hara, Hartmut Liefke, Lucian Popa, Arnaud Sahuguet, and Wang-Chiew Tan.

2 Goals

The research has focused on a technology for “mobile information management”. Underlying this technology is a mathematical foundation enabling the use of formal methods in developing and reasoning about the construction of mobile information management components and their use in database integration and transformation. The salient features of our approach are:

1. Use of unmaterialized views.
2. Dynamic integration of data consumers and data sources, using mobile query processes.
3. Data interface specifications, based on XML schemas.
4. Specifications for transformations, based on XML query languages.
5. The development of formal methods, focusing on query and constraint reformulation.

3 Project Summaries

3.1 Mobile queries and distributed query languages

Sahuguet and Tannen have worked on ubQL a new distributed query language for programming large-scale distributed query systems such as resource sharing systems. The language is obtained by adding a small set of mobile process primitives (communication channels, migration operators, etc.) on top of any traditional query language. Queries are encapsulated into processes and can migrate between sites thus enabling cooperation. An important methodological device is the separation of the installation (including migration) of query processes from the distributed execution of the queries. ubQL allows the encoding of widely used distributed query patterns such as chaining, referral, subscription, leasing, recruiting, query/data/hybrid shipping, etc., and evaluate some
language-based rewrite strategies for the installation of ubQL queries that use only partial and distributed knowledge of execution costs.

Sahuguet and Tannen (with Pierce) have worked on new mechanisms in distributed query optimization. This work outlines a flexible framework for optimizing and deploying distributed queries in wide area networks. The database field has developed very powerful techniques for finding efficient execution plans for declaratively specified queries. However, applying these optimization techniques in the setting of distributed information management requires centralized knowledge of the entire network and assumes passive behavior from the data sources. The reality of the Web is different. Future distributed query optimizers must handle (in fact, exploit!) a rich variety of information flow mechanisms like chaining, referral, proxying, brokering, publish-subscribe, leasing, etc. We look to mobile agent technologies for the combination of flexibility and precision needed for handling these mechanisms. Our language-based approach uses a mobile process calculus based on the pi-calculus in combination with a powerful query-plan language. The salient characteristic of the language is that messaging, migration, and database operations all live in the same semantic space and interact, creating new opportunities for optimization.

3.2 Query reformulation and optimization

Popa and Tannen have studied a class of path-conjunctive queries and constraints (dependencies) defined over complex values with dictionaries. This class includes the relational conjunctive queries and embedded dependencies, as well as many interesting examples of complex value and oodb queries and integrity constraints. We show that some important classical results on containment, dependency implication, and chasing extend and generalize to this class.

Deutsch, Popa and Tannen have continued the work on an optimization method and algorithm designed for several objectives: physical data independence, using materialized views/cached queries, semantic optimization, and generalized tableau minimization. The method relies on generalized forms of chase and “backchase” with constraints (dependencies). By using dictionaries (finite functions) in physical schemas we can capture with constraints useful access structures such as indexes, materialized views, source capabilities, access support relations, gmaps, etc. In this reporting period, we have shown that the method is usable in realistic optimizers by extending it to bag and mixed (i.e. bag-set) semantics as well as to grouping views and by showing how to integrate it with standard cost-based optimization. We understand materialized views broadly, including user-defined views, cached queries and physical access structures (such as join indexes, access support relations, and gmaps). Moreover, our internal query representation supports object features hence the method applies to OQL and (extended) SQL:1999 queries. Chase and backchase supports a very general class of integrity constraints, thus being able to find execution plans using views that do not fall in the scope of other methods. In fact, we prove completeness theorems that show that our method will find the best plan in the presence of common and practically important classes of constraints and views, even when bag and set semantics are mixed.

The search space for query plans is defined and enumerated in a novel manner: the chase phase rewrites the original query into a “universal” plan that integrates all the access structures and alternative pathways that are allowed by applicable constraints. Then, the backchase phase produces optimal plans by eliminating various combinations of redundancies, again according to constraints.

This method is applicable (sound) to a large class of queries, physical access structures, and semantic constraints. We prove that it is in fact complete for path-conjunctive queries and views with complex
objects, classes and dictionaries, going beyond previous theoretical work on processing queries using materialized views.

Popa, Deutsch, Sahuguet and Tannen have studied the practicality this novel method for generating alternative query plans that uses chasing (and back-chasing) with logical constraints. The method brings together use of indexes, use of materialized views, semantic optimization and join elimination (minimization). Each of these techniques is known separately to be beneficial to query optimization. The novelty of our approach is in allowing these techniques to interact systematically, e.g. non-trivial use of indexes and materialized views may be enabled only by semantic constraints.

We have implemented our method for a variety of schemas and queries. We examine how far we can push the method in term of complexity of both schemas and queries. We propose a technique for reducing the size of the search space by "stratifying" the sets of constraints used in the (back)chase. The experimental results demonstrate that our method is practical (i.e., feasible and worthwhile).

Hara and Davidson have studied functional dependencies for nested data. Functional dependencies add semantics to a database schema, and are useful for studying various problems, such as database design, query optimization and how dependencies are carried into a view. In the context of a nested relational model, these dependencies can be extended by using path expressions instead of attribute names, resulting in a class of dependencies that we call nested functional dependencies (NFDs). NFDs define a natural class of dependencies in complex data structures; in particular they allow the specification of many useful intra- and inter-set dependencies (i.e., dependencies that are local to a set and dependencies that require consistency between sets).

### 3.3 XML and semistructured data

XML has become an increasingly popular data-format embraced by a lot of different communities. XML is extremely attractive because it offers a simple, intuitive and uniform text-based syntax and is extensible. One can find today XML proposals for messages, text content delivery and presentation, data content, documents, software components, scientific data, real-estate ads, financial products, cooking recipes, etc. Unfortunately this also means that XML is far too general and if people plan to use it in serious applications (mainly for Electronic Document Interchange, in a broad sense), they will need to provide a specification (i.e. structure, constraints, etc.) for their XML, which XML itself cannot offer. In order to specify and enforce this structure, people have been using Document Type Definitions (DTDs), inherited from SGML and more recently, XML Schema.

Buneman, Davidson, Fan, Hara, and Tan. have investigated integrity constraints for XML data. Both DTDs and the XML Schema proposal lack a clean and general treatment of key dependencies. We discuss the definition of keys for XML documents, paying particular attention to the concept of a relative key, which is commonly used in hierarchically structured documents and scientific databases. We also investigate the (finite) implication problems associated with these dependencies. In contrast to other proposals of keys for XML, these two classes of keys can be reasoned about efficiently. In particular, we show that their (finite) implication problems are finitely axiomatizable and are decidable in polynomial time.

Buneman, Deutsch and Tan have worked on a deterministic model for semistructured data and Buneman and Pierce have worked on union types for semistructured data. Semistructured databases are treated as dynamically typed: they come equipped with no independent schema or type system to constrain the data. Query languages that are designed for semistructured data, even when used
with structured data, typically ignore any type information that may be present. The consequences of this are what one would expect from using a dynamic type system with complex data: fewer guarantees on the correctness of applications. For example, a query that would cause a type error in a statically typed query language will silently return the empty set when applied to a semistructured representation of the same data. We describe a system of untagged UNION TYPES that can accommodate variations in structure while still allowing a degree of static type checking.

Sahuguet has obtained some preliminary results that explore how DTDs are being used for specifying the structure of XML documents. By looking at some publicly available DTDs, we look at how people are actually (mis)using DTDs, show some shortcomings, list some requirements and discuss possible replacements.

Liefke has worked on horizontal query optimization on ordered semistructured data. The exchange and storage of XML data is becoming increasingly important. In contrast to conventional semistructured data, the labels in a document-oriented representation such as XML are ordered. Furthermore, regular expressions (DTDs) describe the horizontal (and vertical) structure. Conventional query languages for semi-structured data ignore the horizontal order and are therefore limited in their expressiveness and optimizability. We describe a query language for querying ordered semistructured data. This query language provides primitives for specifying more powerful queries on ordered semistructured data. Furthermore, we describe how horizontal type information in DTDs is used to optimize queries based on finite automata.

Liefke and Davidson have investigated view maintenance for hierarchical semistructured data. While several important aspects of XML have been investigated, such as query languages, type systems, and storage models, the issue of incrementally maintaining XML views is largely unstudied. XML views differ from relational views in two essential ways: 1) There is no rigid type system, and 2) The query definition often performs complex restructuring far beyond the typical select-project-join query definition in relational views. We address the problem of incrementally maintaining views over XML data with key constraints. We describe a system called WHAX (Warehouse Architecture for XML) that allows the definition and incremental maintenance of views over existing relational and XML data sources with keys. Our query language supports important operations, such as joins, aggregations, regrouping, and restructuring operations such as flattening. We generalize several well-known results about view maintenance in the relational model based on the notion of "multi-linearity". Furthermore, we demonstrate how incremental view maintenance improves the efficiency for XML views defined on real XML data.

3.4 Data Provenance and Annotation

Buneman and Tan (with Khanna) have investigated definitions and properties of the data provenance concept. With the proliferation of database views and curated databases, the issue of data provenance – where a piece of data came from and the process by which it arrived in the database – is becoming increasingly important, especially in scientific databases where understanding provenance is crucial to the accuracy and currency of data. We describe an approach to computing provenance when the data of interest has been created by a database query. We adopt a syntactic approach and present results for a general data model that applies to relational databases as well as to hierarchical data such as XML. A novel aspect of our work is a distinction between "why" provenance (refers to the source data that had some influence on the existence of the data) and "where" provenance (refers to the location(s) in the source databases from which the data was extracted).
Buneman and Tan (with Bird) have investigated the design of a query language for annotation graphs. The multidimensional, heterogeneous, and temporal nature of speech databases raises interesting challenges for representation and query. Recently, annotation graphs have been proposed as a general-purpose representational framework for speech databases. Typical queries on annotation graphs require path expressions similar to those used in semistructured query languages. However, the underlying model is rather different from the customary graph models for semistructured data: the graph is acyclic and unrooted, and both temporal and inclusion relationships are important. We develop a query language and describe optimization techniques for an underlying relational representation.

3.5 Updates

Davidson and Liefke have worked on the problem of maintaining derived data in the context of database changes. “View maintenance” describes the problem of maintaining a materialized view while updating the source database(s). Updates to the source database are either immediately propagated to the view or are accumulated over time and the view is updated in frequent intervals (for instance, during night). “View update” is the problem of propagating updates to the view to the source database.

They have developed a generic update language, CPL+, for updating complex value databases – databases containing values composed of base values, sets, tuples, and variants. The complex value model is a generalization of the relational model. We propose various simplification and optimizations so that an update on a given database is transformed into a more efficient update expression. Further, they extended this work to the object-oriented data model. A new language, OQL+, has been developed to specify updates for such databases in the flavor of OQL and the update primitives known from SQL. Interesting issues such as efficient execution, non-determinism of updates, and cost-based optimizations are investigated in this project.

3.6 Integrated Access to Genomic Data Sources

Davidson, Tannen, et al, have performed and reported on experiments in applications of databases to bioinformatics. The integration of heterogeneous data sources and software systems is a major issue in the biomedical community and several approaches have been explored: linking databases, “on-the-fly” integration through views, and integration through warehousing. We report on our experiences with two systems that were developed at the University of Pennsylvania: an integration system called K2, which has primarily been used to provide views over multiple external data sources and software systems; and a data warehouse called GUS which downloads, cleans, integrates and annotates data from multiple external data sources. Although the view and warehouse approaches each have their advantages, there is no clear “winner”. Therefore, users must consider how the data is to be used, what the performance guarantees must be, and how much programmer time and expertise is available to choose the best strategy for a particular application. Our experiences also point to some practical tips on how updates should be published by the community, and how XML can be used to facilitate the processing of updates in a warehousing environment.

Davidson and Liefke (with Limsoon Wong) have investigated creating and maintaining curated view databases. The process of building a new database relevant to some field of study in biology involves transforming, integrating, and cleansing multiple external data sources, as well as adding new material and annotations. Creating and maintaining these “view” databases raise a number
of problems: 1) How can we specify and implement the transformation and integration from the underlying source databases to the view database? 2) How can we automate the refresh process? 3) How can we track the origins or "provenance" of data? The work discusses these phases of creating and maintaining curated view databases and contrast solutions where appropriate.

References


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Dear Professor Val Tannen:

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