A METHODOLOGY FOR REENGINEERING RELATIONAL DATABASES TO AN OBJECT-ORIENTED DATABASE
THESIS
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AFIT/GCS/ENG/96J-01

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A METHODOLOGY FOR REENGINEERING RELATIONAL DATABASES TO AN OBJECT-ORIENTED DATABASE

THESIS

Presented to the faculty of the Graduate School of Engineering of the Air Force Institute of Technology

Air University

In Partial Fulfillment of the Requirements for the Degree of Master of Science in Computer Science

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Pedro Arthur Linhares Lima
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Abstract

This research proposes and evaluates a methodology for reengineering a relational database to an object-oriented database. We applied this methodology to reengineering the Air Force Institute of Technology Student Information System (AFITSIS) as our test case. With this test case, we could verify the applicability of the proposed methodology, especially because AFITSIS comes from an old version of Oracle RDBMS. We had the opportunity to implement part of the object model using an object-oriented database, and we present some peculiarities encountered during this implementation. The most important result of this research is that it demonstrated that the proposed methodology can be used for reengineering an arbitrarily selected relational database to an object-oriented database. It appears that this approach can be applied to any relational database.
1 Introduction

The software reengineering process has been used to solve many problems involving legacy systems. It has been helping companies to recover and to update documentation, design, and requirements of important systems. Most of the time thousands of lines of code are the only source of the business rules, and are the starting point in the process of reverse engineering. Software reengineering has been playing an important role and has been proven to be very effective in extending the lifetime of many applications.

All systems have a limited lifetime. Each implemented change erodes the structure which makes any following changes more expensive. As time goes on, the cost to implement a change will be too high, and the system will not be able to support its intended task. The reengineering process plays an important role by not allowing the system to reach this condition.

1.1 Background.

The goal of reengineering is to mechanically reuse past development efforts in order to reduce maintenance expense and improve software flexibility. Reengineering is applicable to diverse software such as programming code, databases, and inference logic [1].

There are many possible motives for the reverse engineering of databases [2]:

1
• **Migration between database paradigms.** One may want to migrate between database paradigms, for example from past hierarchical, network, and relational databases to modern relational and object-oriented databases;

• **Migration within a database paradigm.** A more mundane task would be to migrate between different implementations of a database paradigm, for example from one vendor's relational database to another relational database;

• **Documentation.** Reverse engineering can elucidate poorly documented existing software when the developers are no longer available for advice;

• **Tentative requirements.** Reverse engineering of existing software can yield tentative requirements for the new replacement system. Reverse engineering ensures that the functionality of the existing system is not overlooked or forgotten;

• **Assessment of software.** The quality of the database design is an indicator of the quality of the software as a whole. An understanding of the concepts supported by the underlying database schema allows one to better judge functionality claims;

• **Integration.** Reverse engineering facilitates integration of related legacy applications and purchased applications. A logical model of encompassed software is a prerequisite for integration;
- **Conversion of legacy data.** One must fully understand the logical correspondence between the old database and the new database before attempting to convert data.

1.2 Problem.

The main difficulty to reengineering relational databases is the lack of a robust process that can be applied in all cases. Most of the existing processes for database reverse engineering are inadequate; they assume too high a quality of input information [2].

1.3 Hypothesis.

The maintenance of a relational database application can be improved by:

1. Reverse engineering the system to develop an object-oriented model;
2. Redesigning the system using an Object-Oriented Methodology;
3. Changing the Database Management System to one that supports an object-oriented approach.

1.4 Research Objectives.

In order to solve the problem stated above and establish the validity of the above hypothesis, the following objectives were established:

1. Define an appropriate reverse-engineering methodology;
2. Determine an appropriate database application to be a test case;
3. Analyze and reverse engineer the test case using this methodology;
4. Redesign the test case using object-oriented methods;
5. Implement a portion of the new design in an Object-Oriented Database Management System prototype system;

6. Analyze the methodology based on this experience.

1.5 Test case.

With the intention of conducting directly useable research in the field of software reengineering, the director of the Communication Computer System of the Air Force Institute of Technology (AFIT/SC) was contacted. Discussions led to the discovery that his working group was facing a significant reengineering task which could be used as a basis test case for this thesis research.

In 1987 the Air Force Institute of Technology (AFIT) contracted the development of an automated system called Student Tracking and Registration System (STARS). This system is used for scheduling courses, registering students in courses, tracking academic histories of students, and generating related reports. The STARS application uses the Structured Query Language (SQL) to access an Oracle Relational Database Management System (RDBMS) Version 6. This system also uses the following tools: the SQL-Forms, SQL-ReportWriter, SQL-Menu, VMS, and Batch files [3]. From the time the system was designed until this thesis effort, requirements have been changing. Some of these changes were implemented, while others were not.

Even though this system is only eight years old, it is already considered old or a legacy system. This quick obsolescence was caused mainly by the following [4]:

1. Changes were made to incorporate some new requirements; however, documentation was not updated;
2. Past leaders who lacked software knowledge;

3. New technology;

4. Poor training;

5. Lack of focus on changing needs.

The Air Force Institute of Technology Student Information System (AFITSIS) was chosen as the test case in implementing a new method for reengineering relational databases to an Object-Oriented database.

AFITSIS is currently designed and implemented using relational technology and unfriendly user interface mechanisms. This old design and technology cause the maintenance to be difficult, because there are no maintainability features. This lack of maintainability demands a lot of time and effort every time new requirements are implemented on the system. Additionally, the system is inflexible and complex, requiring for each change up to five hundred forms and reports to be updated and checked for consistency.

1.6 Assumptions.

The following assumptions were made for the thesis research:

1. The decision to reengineer AFITSIS instead of starting the analysis and design of a new system is the best decision;

2. Access to AFITSIS and query information from the database are available;

3. Access to an Object-Oriented Database Management System (OODBMS) is available for use.
1.7 Sequence of Presentation.

The thesis is divided into six chapters. Chapter I, Introduction, has provided an overview of the work. Chapter II, Summary of Current Knowledge, discusses the background information which provides the foundation for this research. Chapter III, The Methodology, presents a proposed methodology for reengineering a relational database to an object-oriented database. Chapter IV, Application of the Methodology, presents the application of the proposed methodology using AFITSIS as a test case. Chapter V, Implementation Issues, discusses how the selected part of AFITSIS was implemented using an OODBMS. Lastly, Chapter VI, Analysis, Conclusions, and Recommendations, analyzes the results obtained from the application and implementation of the methodology, draws conclusions from this analysis, and makes recommendations for future applications of this methodology.
2 Summary of Current Knowledge

2.1 Treatment and Organization.

This literature review provides the foundation to create a methodology for reengineering relational database applications to an object-oriented database. This chapter is divided into three sections: software reengineering, reengineering of relational databases, and object-oriented methodology. The software reengineering section gives an overview of the software reengineering process. The reengineering of relational databases section presents the basic steps when reverse engineering relational databases. The object-oriented methodology section describes the stages used by developers to analyze a problem, design a system, and implement the system into a usable product.

2.2 Software Reengineering.

Reengineering, also viewed as both renovation and reclamation, is the examination and alteration of a system to reconstitute it in a new form. Reengineering usually includes some form of reverse engineering (to achieve a more abstract description) followed by some form of forward engineering or re-structuring [5].

Reverse engineering is a process of examination and analysis of the subject to identify its components and create a higher level form of abstraction [5]. It can start at any stage of the life-cycle and it does not involve changes to the subject. Its sub-products include the design recovery and the redocumentation of the subject. Forward engineering can be easily understood as a process of moving from a high-level of
abstraction to low-level or physical details. It is the same as the traditional method of
developing a new system. This term is used only to distinguish this process from
reverse engineering. Figure 1 illustrates the basic ideas of software reengineering
using, for simplicity, only three life-cycle stages of software.

Figure 1: Relationship between terms [5]
The objectives of Software reengineering can be grouped into four main areas [6]:

1. **Improve maintainability.** The maintenance efforts can be reduced by reengineering smaller modules with more explicit interfaces. However, it is not easy to measure progress toward this goal.

2. **Migration.** This task usually deals with altering and converting program structure. This goal can be easily measured, since the system will perform the same operation in the new environment.

3. **Achieve greater reliability.** This goal can be easily reached because the restructuring process usually causes most of the potential defects to appear. The other factor that contributes to better software reliability is the extensive testing required to prove the functional equivalence between the old and the new system. This goal can be readily measured by fault analysis.

4. **Preparation for functional enhancement.** Once the programs are decomposed into smaller modules, it is easier to isolate them from one another. This makes it simpler to change or add new functions without affecting other modules.

2.3 Reengineering of Relational Databases.

The goal of reengineering is to mechanically reuse past development efforts in order to reduce maintenance expense and improve software flexibility. According to Hainaut [7] the most tractable approach for database applications is to first reverse
engineering the database and then deal with the programming code. Object-oriented models provide a natural language for facilitating the reengineering process. An object-oriented model can describe the existing software, the reverse-engineering semantic intent, and the forward-engineered new system.

In general, the mapping between object models and a database schema is many-to-many. Various optimizations and design decisions can be used to forward engineer an object model into a database schema. Similarly, when reverse engineering a database, alternate interpretations of the structure and data can yield different object models. Usually there is no obvious, single correct answer for reverse engineering. Multiple interpretations can all yield plausible results [2].

A good way to begin reverse engineering is by entering the existing schema into a CASE tool. Associations will often be found in a degraded form such as relational database foreign keys. Inheritance must be implemented in a degraded manner for current relational database managers. The schema may then be gradually transformed to a logical model as underlying relationships are inferred.

Jacobson [8] presents a good approach for reengineering old systems to an object-oriented architecture, but he does not give much information when dealing with relational databases. The same problem exists when considering other approaches for reengineering like those of Bennett [9] and Sneed [6]; they are not focusing on relational databases.

A good approach is suggested by Blaha [1], [2]. His papers present some typical implementation strategies that are used for forward engineering. He explains in
detail each step to be taken for reverse engineering of relational databases. The basic steps he suggests are:

**Step 1.** Prepare an initial object model.
- Represent each table as a tentative class. All columns of tables become attributes of classes.

**Step 2.** Determine candidate keys.
- Look for unique indexes. Automated scanning of data can yield potential candidate keys.

**Step 3.** Determine foreign-key groups.
- Try to resolve homonyms, attributes with the same name that refer to different things, and synonyms, attributes with different names that refer to the same thing.
- Matching attribute names, data types, and/or domains may suggest foreign keys.
- During this step do not attempt to determine specific reference-referent attribute pairs – but merely groups of attributes within which foreign keys may be found.

**Step 4.** Refine tentative classes.
- Agglomerate horizontally partitioned classes into a single class. (horizontally partitioned classes must also have the same semantic intent.)
- Detect functions and constraints that are represented as tables.
Step 5. Discover generalizations.

- Analyze large foreign-key groups, particularly those with 5, 10, or more cross-related attributes.
- Look for patterns of many replicated attributes.
- Look for patterns of data where a class has mutually exclusive subsets of attributes.
- When discovering generalizations do not forget there may be a forest of generalizations with multiple superclass roots and intermediate levels.

Step 6. Discover associations.

- Convert a tentative class to an association when a candidate key is a concatenation of two or more foreign keys.
- Introduce a qualified association when a candidate key combines a foreign key with non-foreign key attributes.
- The remaining associations are buried and manifest as foreign keys.
- Note minimum multiplicity for associations. Optional multiplicity is the permissive case; a lower limit of one (or another number) is more restrictive.
- Note maximum multiplicity for associations. Many multiplicity is the permissive case; an upper limit of one (or another number) is more restrictive.
Apply semantic understanding and restate some associations as aggregations. Aggregation is the “a-part-of” relationship.

**Step 7.** Perform transformation.

- Convert a class to a link class as needed.
- Lightweight one-to-one associations should be more simply represented as an attribute.
- Nonatomic n-ary associations should be decomposed into their constituent associations of lesser order.
- Consider shifting associations via transitive closure.
- Double-buried associations should be merged into a single association.
- You may need to insert an intermediate class in a generalization hierarchy to recognize common semantics, attributes, and associations.
- Transitive closure also arises through the combination of generalization and association. Where possible, eliminate an imprecise association to a superclass in favor of a more restrictive association to a subclass.
- Similarly, eliminate associations to subclasses by recognizing patterns of commonality.
2.4 Object-Oriented Methodology.

One of the primary reasons for adopting object technology is the promise of faster development and reduced maintenance costs. In traditional systems, ongoing maintenance costs amount to more than 80% of the overall cost of the system [10]. Object-oriented systems promise to reduce maintenance costs through reusable objects that can dramatically reduce maintenance. In many cases, developers only need to identify an object class that functions like the object that they desire to create, and specify the differences between the object and their new object. This type of code reusability can dramatically reduce development and maintenance costs.

Object-oriented methodology allows developers to analyze problems and divide them into entities residing in specific states and exhibiting certain dynamic behaviors. The entities become objects in the system. The designer defines the relationships between the objects to determine how the system functions as a whole. The four specific stages of object-oriented methodology are [11:4-6]:

1. **Analysis.** During the analysis stage, the developer defines the system requirements. Objects are identified and their relationships to other objects are recorded. There are no implementation decisions in this stage. Three models are defined in this stage: an object relationship model, a dynamic model, and a functional model;

2. **System Design.** In this stage the system’s architecture is determined. The application is broken into subsystems. Control mechanisms are defined for
each subsystem. The focus is on what needs to be done, and not how it is to be done;

3. **Object Design.** During this phase, the object relationship model, dynamic model, and functional model are evaluated to determine what operations must be implemented for each object. Structures for representing the relationships between objects are defined.

4. **Implementation.** The final stage involves transforming the design into an executable system. This is dependent on whether the software language selected supports object-oriented programming.

### 2.5 Conclusion

This literature review has provided an overview of the basic concepts of software reengineering, the reengineering of relational databases, and object-oriented methodology. All three of these areas are required for the successful analysis and implementation of the new methodology.
3 Methodology

3.1 Introduction

This chapter presents the methodology for reengineering a relational database to an object-oriented database. It shows the methodology step by step explaining each step in detail, including some discussion of typical implementation techniques that one can find during the process of reverse engineering.

This methodology is based on Blaha [1] [2] with some changes. His papers were selected because they are focused specifically on reverse engineering of a relational database to an object-oriented database and they are the only ones that give detailed information on this subject.

Some changes were introduced on his approach just to facilitate the transition from relational to an object-oriented view. The most important changes are:

1. Construct an entity-relationship model instead of going directly from the tables to an object model;
2. Besides the object model, prepare a functional model to facilitate the implementation of the system.

3.2 The Methodology

This methodology is presented in a linear fashion for ease of understanding, but, except for the first and last step, the others steps are weakly ordered since during the process of reverse engineering there is much iteration and backtracking. The steps are as follows:
Step 1. Prepare an entity-relationship (ER) model.

This step can be easily accomplished by using an automated tool. Otherwise proceed as listed below:

- Represent each table as an entity.

- Determine candidate keys. Look for unique indexes, but some candidate keys may not be enforced by unique indexes. Automated scanning of data can yield potential candidate keys.

- Determine primary keys. Ordinarily every table should have a primary key. But exceptions can be encountered as follow:

  1. Tables with temporary data or tables which the performance overhead can not be tolerated.

  2. Missing primary key without cause. Some applications enforce primary keys with custom code and do not rely upon the database manager.

  3. Null primary key attributes. Some relational database managers require that one define a unique index to enforce a primary key. Indexed attributes are permitted to be null, unless “not null” is specified for each of the attributes. This violates the definition of primary key; attributes in a primary key may not be null.
4. Extraneous primary key attributes. By definition a primary key must also be minimal; no attribute can be discarded from the primary key without destroying uniqueness. The reverse engineer must regard all primary key declarations with suspicion, and look for attributes that do not seem semantically justified.

Even when tables do have a primary key, different realizations may still be chosen. Figure 2 shows relational tables for three different approaches to identify the primary key. All three schemas can be reverse engineered to the same logical model.

- **Artificial identity.** Each object table (shown in Figure 2) has an object identifier as primary key. Association tables (not shown in Figure 2) have a primary key consisting of the identifiers of the related objects.

- **Value-based identity.** The primary key of each object consists of some combination of application attributes. Some primary keys may become lengthy, as attributes are incorporated from foreign key of related tables.

- **Hybrid identity.** One may use artificial identity and value-based identity in the same schema. In the third segment of Figure 2 Bank has artificial identity and Account has identity derived from a reference to a bank combined with an account number.
Reverse engineering input: Artificial identity

Bank

<table>
<thead>
<tr>
<th>bank ID</th>
<th>bank name</th>
</tr>
</thead>
</table>

Account

| account ID | bank ID | account number |

{Candidate key of bank is: bank name.}
{Candidate key of Account is: bank ID + account number.}

Reverse engineering input: Value-based identity

Bank

<table>
<thead>
<tr>
<th>bank name</th>
</tr>
</thead>
</table>

Account

| bank name | account number |

Reverse engineering input: Hybrid identity

Bank

<table>
<thead>
<tr>
<th>bank ID</th>
<th>bank name</th>
</tr>
</thead>
</table>

Account

| bank ID | account number |

{Candidate key of bank is: bank name.}

Reverse engineering output: Logical intent

Bank

| bank name | account number |

Account

Figure 2: Various approaches to identify the primary key [2]

- Determine foreign-keys. Most of the modern RDBs have a foreign-key clause as part of the schema. If you do not have this do the following:
  - Try to resolve homonyms, attributes with the same name that refer to different things, and synonyms, attributes with different names that refer to the same thing.
• Matching attribute names, data types, and/or domains suggest foreign keys.

• Generate the relationships by checking every possible foreign key against every candidate key.

• Finish the ER model by querying the data and determining the multiplicity of each relationship.

**Step 2. Prepare an initial object model.**

Based on the ER diagram, represent each entity as a tentative class and each relationship as a tentative association. All columns of the related tables become attributes of classes.

**Step 3. Refine tentative classes.**

Agglomerate horizontally partitioned classes into a single class. Horizontally partitioned classes have the same schema. Distributed databases often use horizontal partitioning to disperse records. (Horizontally partitioned classes must also have the same semantic intent. Identical schema is a good indicator of same semantic intent.)

Detect functions and constraints that are represented as tables and take these classes out of the tentative object model. Look for classes that do not participate in any foreign key.
Step 4. Discover generalizations.

Analyze large foreign-key groups, particularly those with 5, 10, or more cross-related attributes. Look for a primary key that is entirely composed of a foreign key of another table. Derived identity is symptomatic of an implementation of generalization with distinct superclass and subclass tables or propagation of identity via one-to-one association. Data analysis can increase confidence in the discovery of generalization by revealing subsets of records.

Look for patterns of many replicated attributes. A generalization may have been implemented by pushing superclass attributes down to each subclass.

Look for patterns of data where a class has mutually exclusive subsets of attributes. This may indicate an implementation of generalization where subclass attributes were pushed up to the superclass.

When discovering generalizations one must not forget there may be a forest of generalizations with multiple superclass roots and intermediate levels. Data analysis can help distinguish multiple, disjoint, and overlapping inheritance. (Keep in mind that data analysis only yields hypotheses, and semantic understanding is required to reach firm conclusions.)

Step 5. Discover associations.

Convert a tentative class to an association when a candidate key is a concatenation of two or more foreign keys. Where possible, try to restate ternary and
n-ary associations (confluence of primary keys from three or more classes) as binary associations[2].

Introduce a qualified association when a candidate key combines a foreign key with non-foreign key attributes. This will find some, but not all, qualifiers.

The remaining associations are buried and manifest as foreign keys.

Note minimum multiplicity for associations. Optional multiplicity (nulls allowed) is the permissive case as for a given record you may store an actual value or store a null; a lower limit of one (or another number) is more restrictive.

Note maximum multiplicity for associations. Many multiplicity is the permissive case as a collection can store a single value or many values; an upper limit of one (or another number) is more restrictive.

Apply semantic understanding and restate some associations as aggregations. (Aggregation is the “a-part-of” relationship.)

When discovering associations be aware to the following kind of implementations that one may encounter [2]:

- **Double-buried associations.** This is when an association was buried in both participating classes as shown in Figure 3. This construct complicates reverse engineering, since these double-buried associations look like two separate associations. Data analysis can detect redundancy between the dual pointers, but semantic understanding is required to resolve this situation.
### Model as implemented

<table>
<thead>
<tr>
<th>A table</th>
<th>B table</th>
</tr>
</thead>
<tbody>
<tr>
<td>A primary key</td>
<td>B primary key</td>
</tr>
<tr>
<td>B foreign key</td>
<td>A foreign key</td>
</tr>
<tr>
<td>other A attributes</td>
<td>other B attributes</td>
</tr>
</tbody>
</table>

#### Logical intent

- **Optional qualified association.** Figure 4 shows an optional qualified association. A *cluster* contains many *Tables*. A *Table* may belong to at most one *Cluster*. The combination of a *Cluster* and a *table#* yields a specific *Table*. This association was implemented by burying *cluster_id* as a foreign key in *Table*. Because of the optional membership in a cluster, the foreign key can be null, and the combination of *cluster_id* and *table#* is not a candidate key of *Table*. Therefore it is difficult to detect this qualified association.

- **Alternate qualifier.** In Figure 5 *Column* derives its identity from a *Table* plus a qualifier, either *column name* or *column number*.

---

*Figure 3: Double buried association*
Step 6. Perform transformation.

Various optimizations may have been employed in preparing the original RDB schema to improve time and/or space performance. Some transformations are listed here [1].
• Convert a class to a link class as needed. A link class is an association whose links can participate in associations with other classes. An association has derived, rather than intrinsic, identity.

• Lightweight one-to-one associations (they have no attributes) should be more simply represented as an attribute. For example, it is unnecessary to represent city as a class, when city-name is the only attribute of interest.

• Nonatomic n-ary associations should be decomposed into their constituent associations of lesser order. Binary associations are most common and easier to understand. We may find ternary association, but never an association of higher order.

• Consider shifting associations via transitive closure. For example associations from A to B and B to C could possibly be restated as associations from A to B and A to C. In general, multiplicity constrains derivation of association, but the vague multiplicity limits often obtained through reverse engineering allow more latitude.

• Double-buried associations should be merged into a single association. For example, an association between A and B may have been buried in both the A and B classes.

• You may need to insert an intermediate class in a generalization hierarchy to recognize common semantics, attributes, and associations.

• Transitive closure also arises through the combination of generalization and association. Where possible, eliminate an imprecise association to a
superclass in favor of a more restrictive association to a subclass. For example, in Figure 6 if our semantic knowledge is that X only associates with B and never with C, then we can eliminate the association between X and A and the association between X and C.

![Figure 6: Transitive closure involving generalization and association](image)

- Similarly, eliminate associations to subclasses by recognizing patterns of commonality. In Figure 6, if all instances of B partition across classes D and E, we can eliminate the association between X and D and the association between X and E.

**Step 7. Prepare a functional model.**

The functional model describes computations within a system, and specifies the results of this computation without specifying how or when they are computed. Database system often have a trivial function model, since their purpose is to store and organize data, not to transform it [11:123].
One can prepare the functional model only using the user manual, the forms, and, if necessary, interviewing the users.

3.3 Summary

This chapter has presented the methodology for reengineering a relational database to an object-oriented database. This methodology is heavily based on Blaha papers [1] and [2], except for the first step, that was introduced to facilitate the transition from relational to an object-oriented view, and the last step, that was introduced to give more information about the functionality of the system. It showed each step to be followed with some discussion of typical implementation techniques. The next chapter presents the application of this methodology using AFITSIS as a test case.
4 Application of Methodology

4.1 Introduction

This chapter presents the application of the proposed methodology using AFITSIS as a test case. It is divided into three sections: the first section shows how the ER model was obtained. The next section presents the transformations that were made to the ER model to obtain the object model. The last section shows how the functional model was drawn.

Following direction of the sponsor (AFIT/SC), this analysis is restricted to those tables and forms that have some relationship to the Person table. This restriction does not invalidate the work, since about 66 of 294 tables from the entire AFITSIS are considered.

4.2 ER Model

To accomplish the first step of the methodology, which is to draw the ER model, we used ERwin (an ER diagram editor developed by Logic Works [12]). Since AFITSIS was developed for Oracle version 5, which does not support foreign-key clauses, and migrated to Oracle version 6 without changes, ERwin was able to capture only the tables and its attributes (all 294 tables from AFITSIS). If you are reverse engineering a RDB that supports foreign-key clauses, ERwin can recover not only the tables and their attributes, but also foreign-keys, the relationships between tables, and can draw the entire ER model.
We started our work identifying the Person’s primary key (SSAN). Next we selected all tables that have this primary key as an attribute by querying the Oracle data dictionary (Figure 7). This query resulted in 66 tables (Appendix A).

```
SELECT table_name
FROM accessible_columns
WHERE column_name
like "%SSAN%";
```

*Figure 7: SQL statement to find tables with SSAN as an attribute*

The next step was to determine the candidate keys for each of the selected tables. We looked for unique indexes in the data dictionary, and for each one that we found we scanned the data to confirm the correctness. For the other tables for which we could not find a unique index, we had to scan the data.

During the process of scanning the data to look for the primary key, we were able to find many tables that have no data, tables that have not been used for many years, and tables that were used as a temporary files. After we confirmed that they were not being used by any form, we eliminated these tables from our diagram.

We looked for foreign key groups by matching attribute names and types. We did not have any complication in this step, especially because the names are very suggestive and we did not find any homonyms nor synonyms.

Next, we were able to generate the relationships between the entities by checking every possible foreign key against every candidate key, and linking the
related entities using ERwin. In our model we did not consider as a relationship the link of a table with a validation table, via its foreign key. For example: Address table has as foreign keys the attributes Address_type_code, Street_type_code, Address_room_type_code, and country_code; which are the primary key of the following validation tables (table look up): Address_type_valid, Street_type_code_valid, Address_room_type_code_valid, and Country_valid, respectively.

To finish the ER model (see Appendix B) we queried the data to determine the multiplicity of each relationship, doing the following:

- **One-to-one association.** To determine a one-to-one association we verified if for each row in one table of the relationship there was only one entry into the other table.

- **One-to-one-or-more association.** For this type of association we verified whether for each row in one table we found at least one or more entries into the other table.

- **One-to-zero-or-one association.** In this case we verified whether for each row in one table we could find zero or exactly one entry into the other table.

- **One-to-zero-one-or-more association.** Now we verified whether for each row in one table we found zero or more entries into the other table.
4.3 Object Model

To draw the object model we started preparing an initial object diagram (step 2) based on the ER diagram, where we represented each entity as a tentative class and each relationship as a tentative association. We transformed all columns of the related entity into the attributes of the class.

We started refining the object model (step 3) by looking for horizontally partitioned classes (classes with the same schema) and representing them as a single class. This is what we found:

- The classes *Term_entry* and *Term_entry_history* had the same schema and the same semantic intent. We merged them into a single class;

- The classes *Selected_student*, *Selected_student_91*, and *Selected_student_new_91*, had the same schema but, after checking the data, we determined that the classes *Selected_student_91* and *Selected_student_new_91* were used as temporary files. We retained only the class *Selected_student* and eliminated the others.

Then we looked for tables that could have been representing functions and/or constraints, but we did not find any.

We started the process of discovering generalization (step 4) by looking for large foreign-key groups. Although we found a couple of tables in this case we
realized that these tables were not involved in a generalization but in a binary or ternary association.

Generalization was found when we started looking for any class that had its primary key entirely composed of a foreign key of another class. We took these classes apart and analyzed their relationship.

To do a good analysis of this kind of relationship we had to improve our semantic knowledge of the system. We did that by making some queries and analyzing its results, by looking up the forms, and by interviewing the Database Administrator (DBA). After that, we were able to take these classes and select those involved in a generalization from those involved in an association. For example: the tables Spouse_info, Emergency_data, AFIT_user_name, Recall_roster, Dependent_information, and Graduation_name all have their primary key entirely composed of Person’s primary key, but they have no inheritance relationship with this table.

Semantic knowledge was especially important to incorporate some abstract classes. For example, in the object diagram in Figure 15 (Appendix C), the abstract classes Civilian and Military were introduced after we discovered that Faculty, Student, and Administrative people could be either military or civilian, but only military people have a relationship with Rank_history and Recall_roster. So, we decided to introduce these two abstract classes to increase code reuse and to organize features common to these subclasses.
Another generalization chain was encountered when we analyzed Student, CI_student, Resident_student, and INTL_student tables. We found out that every instance of CI_student and Resident_student was in the Student table, and that every instance of INTL_student was in the Resident_student table.

We introduced the class Part-time_student to represent another kind of student, after we discovered that this class was implemented as an attribute of Resident_student class called program_code. One of the valid values is ‘PTE’, meaning ‘part-time student’.

Data analysis was very important to increase confidence in the discovery of generalization. Figure 8 shows our initial object diagram where we made some assumptions based on our understanding of the system. But, after we analyzed the data in the Eligible, Selected_students, and Student tables, we discovered that, contrary to our assumption, not all instances of Student could be found in the Selected_student table, and that not all instances of Selected_student could be found in the Eligible table. This data analysis led us to change our object diagram to the one shown in Appendix C.

We continued our work of drawing the object model by discovering other associations (step 5). We started this step by looking for candidate keys composed of two or more foreign keys, and we converted it into an association. Figure 9 shows one binary association that we found. All the other associations can be seen in Appendix C.
Figure 8: Initial object diagram before the data analysis

Figure 9: Binary association
We introduced a qualified association when we found a candidate key that combined a foreign key with a non-foreign key attribute. Figure 10 shows some of the qualified associations that we introduced to reduce the effective multiplicity of the association (from many to one), to improve semantic accuracy, and to increase the visibility of navigation paths.

![Diagram of qualified associations](image)

**Figure 10: Qualified associations**

In our research we were not able to find the following cases of association:

- Association that could be representing an aggregation;

- Double-buried associations;

- Optional qualified association;
• Alternate qualifier.

We started our final step (step 6-perform transformation) to get the object model by transforming each lightweight one-to-one association into an attribute. We found a one-to-one association of Person with AFIT_user_name and Person with Emergency_data and we transformed these associations into attributes of Person (Person Structure Definition, Appendix C). We did the same with a one-to-one association of Resident_student with Graduation_name (see Resident Student Structure Definition, Appendix C).

We did not find any class that could be better represented as a link class. This does not mean that none of our classes is a link class; the only restriction is that our view is limited, since our research is concerned with only part of AFITSIS. We did not have to do any work to decompose n-ary associations into their constituent associations of lesser order, especially because we had only binary associations. Our complete object model, including the object structure definition, can be seen in Appendix C.

4.4 Functional Model

In order to get the functional model done we initially made use of the STARS User's Guide [3] to select the boundaries of what we were going to implement, and then limited our work in doing the functional model to this specific part. We used SQLForms to extract the name of the tables that each form can read or update, what
actions the form is doing, and all the other information needed to draw the functional model. The complete functional model is shown in Appendix D.

4.5 Summary

This chapter has presented the application of the proposed methodology using AFITSIS as a test case. It showed how the ER model was obtained, the transformations that were made on the ER model to get the object model, and how the functional model was drawn. The next chapter discusses the implementation of part of AFITSIS using an OODBMS.
5 Implementation Issues

5.1 Introduction

This chapter presents the implementation issues concerning the development of part of AFITSIS using an OODBMS. It is divided into three sections: the first section discusses how we analyzed several OODBMS and why we chose Microsoft Visual Foxpro (Foxpro) version 3 to be used in this implementation. The next section shows how part of the object model was implemented. The last section discusses some limitations encountered when using Foxpro as an OODBMS.

5.2 Analysis and Choice of the OODBMS

During the process of choosing one OODBMS to implement part of AFITSIS we took the following considerations:

- Based on the interview [4], AFIT/SC wanted to migrate AFITSIS to an OODBMS. However, before making any decision, they will wait for Oracle Company to release version 8, expected to be an extension of the RDBMS with some object capabilities;

- Since Oracle version 8 is not expected to reach the market until the end of 1996 or beginning of 1997, and with the intention of doing a useful implementation, we looked for an available RDBMS that would have some similarities with the expected Oracle version 8. These similarities that we were concerned about are: the product should be able to use all the power
and flexibility of RDBMS, like Data Definition Language (DDL) and Data Manipulation Language (DML), share the basic relational tables, and incorporate some concept of "object," and have the ability to store procedures as well as data in the database.

With these considerations in mind we started analyzing some available OODBMS. The first two OODBMS that we analyzed were ITASCA and Objectstore. Even though we concluded that each is a very good OODBMS, we decided not to use either of them because they are heavily based on some language like C or C++, and they have no compatibility and similarity with any RDBMS.

The next product that we analyzed was Foxpro. Once we had some experience in using an old version of Foxpro and knowing that it is a RDBMS, we concentrated our analysis to see if the new object-oriented features would be compatible with what we wanted. After we read the Foxpro Developer's Guide [13] and used it for two weeks, we were convinced that this product could give us a good means of comparison and insight of what we could expect when we have the Oracle version 8 available.

5.3 Implementation of the Object Model

We started the implementation of our object model by creating a new project and inserting a new database that we called Stars. In Foxpro the terms database and table are not synonymous. The term database refers to a relational database that stores information about one or more tables or views [13]. The database is where we can
create stored procedures (that can be used as field- and record-level rules) and persistent table relationships (to enforce referential integrity).

After we created the *Stars* database we created the definition of the tables that we were going to use, their primary key and indexes, and we added these tables to the database. Then we linked the tables to set up the relationships (Figure 33, Appendix E), so that we do not need a code program to check the referential integrity every time an application tries to modify the database. The database manager system takes care of it whenever the database is opened and used.

The next step was to create the forms, one for each table. After that, we were ready to start creating the definitions of the classes. To implement the Person's Object Model (Figure 15, Appendix C) we did the following:

- We created the **Person's** class based on the **Person's** form (Figure 34, Appendix E);
- We created the **Military's** class based on the **Person's** class and adding the **Military's** form (Figure 35, Appendix E);
- We created the **Military_student's** class based on the **Military's** class and adding the **Student's** form (Figure 36, Appendix E);
- We created the **Military_faculty's** class based on the **Military's** class and adding the **Faculty's** form;
• We created the Military\_resident\_student’s class based on the Military\_student’s class and adding the Resident\_student’s form (Figure 37, Appendix E);

• We created the Military\_INTL\_student’s class based on the Military\_Resident\_student’s class and adding the INTL\_student’s form (Figure 38, Appendix E);

• We created the Civilian’s class based on the Person’s class and adding the Civilian’s form (Figure 39, Appendix E);

• We created the Civilian\_student’s class based on the Civilian’s class and adding the Student’s form (Figure 40, Appendix E);

• We created the Civilian\_faculty’s class based on the Civilian’s class and adding the Faculty’s form;

• We created the Civilian\_resident\_student’s class based on the Civilian\_student’s class and adding the Resident\_student’s form (Figure 41, Appendix E);

• We created the Civilian\_INTL\_student’s class based on the Civilian\_resident\_student’s class and adding the INTL\_student’s form (Figure 42, Appendix E);

The reason we created the classes this way was to give more flexibility and to make the classes easier to maintain. For example: if we need to make some change in
person's class we do not have to modify all the other classes that use it. Because of the inheritance feature, the changes that we make in the parent class reflect over the subclasses automatically; and if we need to overload some parent method (function, procedure, trigger, or event) so that it takes a different action when running the subclass, it can be easily achieved by just creating a method in the subclass with the same name as the parent class. This way the subclass method will have precedence over the parent class method.

We implemented binary associations by first creating a view with the related tables and then creating a form based on this view. For example: the association between Person and Address (Figure 16, Appendix C) was implemented by creating a view with Person and Address tables, and then creating a form using this view, the Person class, and the Address form (Figure 43, Appendix E).

We implemented a binary association with link attributes by creating a new table with these link attributes and the primary key of the two associated tables as foreign keys. We created a form based on this new table, and a class based on this form. For example: the Dropped Courses Association (Figure 22, Appendix C) was implemented by creating a Dropped_courses table having the primary key of the Course and Resident_student tables as foreign keys, plus the link attributes. Then we created a form base on the Dropped_courses table and a class Dropped_courses based on this form.
5.4 Limitations of Foxpro Encountered During Implementation

During the implementation of our object model we faced two major problems when using Foxpro as an OODBMS. The first one is that in Foxpro we can’t define one class based on a table by only adding its methods. Instead, we have to define the table, define one form based on this table, and then define the class based on this form and add the methods. Actually, this peculiarity does not cause much of a problem (when you need to modify some attribute, you have to change the table structure and its related form), but it is a little different from what we learned in theory [11].

Another problem encountered is that Foxpro did not appear to support multiple inheritance. To implement the Person’s Object Model (Figure 15, Appendix C), instead of defining only one class for Student, Faculty, and Administrative, we had to define the classes Military_student, Military_faculty, Military_administrative, Civilian_student, Civilian_faculty, Civilian_administrative, and so on. This restriction may cause some problems if the subclass has other relationships. For example: the relationship Advises between Faculty and Resident_student (Figure 15, Appendix C), has to be implemented by defining one relationship Advises from Civilian_faculty to Civilian_resident_student and Military_resident_student, and the same relationship Advises from Military_faculty to Civilian_resident_student and Military_resident_student.
5.5 Summary

This chapter has presented the implementation issues of part of AFITSIS. Some OODBMS were analyzed and Foxpro was chosen because of its similarity to what we are expecting for Oracle version 8. We showed the implementation techniques that we used to implement inheritance and some associations. During the implementation we found two limitations in using Foxpro as an OODBMS. One is that we can’t define a class directly from a table and the other is that it does not support multiple inheritance. With this implementation done, we had the last piece of information necessary to make an analysis and conclusion of our research. That is presented in the next chapter.
6 Analysis, Conclusions, and Recommendations

6.1 Analysis of The Results

In Chapter III we presented a methodology for reengineering a relational database to an object-oriented database. To validate this methodology we applied it to reengineering AFITSIS as a test case. As we presented in Chapter IV, this methodology is easy to use in practice. We did not have any difficulty when following its steps.

Our methodology has the purpose of reengineering a relational database, independent of the kind of the RDBMS and its version. With our test case, we had the opportunity to verify this applicability, especially because AFITSIS comes from an old version of Oracle RDBMS. This way we could apply most of the steps of what we proposed in the methodology. For example: to accomplish the first step of the methodology, which is to draw the ER model, we used ERwin. Since AFITSIS was developed for Oracle version 5, which does not support foreign-key clauses, ERwin was able to capture only the tables and their attributes. Since ERwin was not able to draw the entire ER model and facilitate this job, we really had to apply the methodology and follow its steps to recover the foreign-keys and the relationships between tables.

When applying our methodology to draw the object model we found out that semantic logic can play an important role, especially to discover generalization and to
incorporate some abstract classes. Another important factor that we found that increased our confidence in the discovery of generalization was data analysis. After we analyzed the data we could change our first object diagram to the one shown in Appendix C. Even though in our test case we were not able to apply our methodology to exemplify the discovery of all kinds of associations, we were able to find some of them.

The most important result of this analysis is that it demonstrated that the proposed methodology can be easily used for reengineering any relational database to an object-oriented database, filling the lack of a robust process that can be applied in all cases.

6.2 Conclusion

The life span of an information system consists of specification, design and maintenance. The maintenance phase dominates in time and often with respect to resources as well. During this phase the system is subjected to a number of changes and additions. The gap between the older technology in the system and the new technology that becomes available increases successively. Changes in the activities of an organization also mean that systems grow old.

Gradually the system approaches a limit where it no longer is cost-efficient or even technically feasible to continue the maintenance. But the cost of enforcing the required changes is usually very high [8]. A possible way out of this dilemma is to
define well delimited system parts that are candidates for modernization. This is where reengineering can help.

We have described a practical method for reengineering. The method is based on object-oriented modeling. We have described how the work can be divided into a number of steps so that the method can be performed in a systematic manner.

We have used AFITSIS as a test case and have shown that with this method we can model an existing system in a simple manner and with limited effort. The new model is object-oriented and can serve as a basis for a future development plan.

We have implemented part of AFITSIS using Foxpro, one OODBMS that we have chosen because of its similarities with the expected Version 8 of Oracle. From this experience we were able to see that our object model can be easily mapped to be implemented using another OODBMS.

The six research objectives, as stated in Chapter I, were:

1. Define an appropriate reverse-engineering methodology;
2. Determine an appropriate database application to be a test case;
3. Analyze and reverse engineer the test case using this methodology;
4. Redesign the test case using object-oriented methods;
5. Implement a portion of the new design in an Object-Oriented Database Management System prototype system;
6. Analyze the methodology based on this experience.
The research was successful in all the original objectives. We presented a practical methodology that can be applied for reengineering any relational database system. We chose AFITSIS as a test case, we applied our methodology for reengineering it, we obtained an object and functional model from this work, and we implemented this model using an OODBMS. Finally, one of the most important lessons that we have learned when working with reengineering is that in general, the mapping between object models and schemes are many-to-many. Various optimizations and design decisions can be used to transform an object model into a database schema. Similarly, when reverse-engineering a database, alternate interpretations of the structure and data can yield different object models. Usually, there is no obvious, single correct answer for reverse engineering. Multiple interpretations can yield plausible results [1].

6.3 Recommendations

For those who intend to use the object model obtained from our test case, we recommend that you revise this model making another data analysis. This is because we had restricted access to AFITSIS tables, since they record confidential information. Doing this you can have more confidence on the model, and may find some important information that we were not able to uncover.

Even though Foxpro was demonstrated to be a good OODBMS to be used in our test case, I recommend further analysis concerning its security of the data. This is because security is an important aspect to be considered for adopting an OODBMS to
implement AFITSIS, and Foxpro does not appear to have any mechanism to restrict the access to the databases (for example: password with level of access.)
Appendix A: List of Tables with SSAN

(Pk): Primary key; (Fk): Foreign key.

Address: SSAN (Pk) (Fk), Address_Type_Code (Pk), Address_Line_1, Address_Line_2, Address_Line_3,
City, State_Code, Zipcode, Zipcode_Extension, Country_Code, Area_Code, Phone_Number,
Address_Effective_Date, DSN_Prefix, Login_Name, Firm_Name_Office_Symbol,
Additional_Address_Information, Street_Address, Street_Type_Code,
Address_Room_Type_Code, Address_Room_Type_Number, Revision_Name, Revision_Date,
Country, Login_Date, Phone_Number_Ext.

Address_Data_Final: SSAN (Pk) (Fk), Address_Type_Code (Pk), Firm_Name_Office_Symbol,
Additional_Address_Information, Street_Address, City, State_Code, Zipcode,
Zipcode_Extension, Street_Type_Code, Address_Room_Type_Code,
Address_Room_Type_Number, Area_Code, Phone_Number, Address_Effective_Date,
DSN_Prefix, Login_Name, Revision_Name, Revision_Date, Country.

AFITnet_User_Name: SSAN (Pk) (Fk), Login_Name, Input_Date, User_Name, User_UID,
Host_Accnt_Created

Ci_Student: SSAN (Pk) (Fk), Major00career_Pointer_Code, Academic_Status_Code,
Accounting_Status_Code, Book_Payment_Authorize_Code, Ci_Student_Comment,
Civilian_Institution_Code, Corps_Code, Course_Of_Study, Current00ASC_Code,
Current00ewi_Option_Code, Email_Address, Grad_Date, Gre_Status_Code,
Hpsp_Medical_Code, Ida_Date, Kit_Sent_Date, MAJCOM_Abbrev, Major00ASC_Code,
Motorcycle_Status_Code, Motorcycle_Train_Date, Office_Code, Overseas_Indicator,
Program_Entry_Date, Quota_Program_Code, Report_No_Earlier_Than_Date,
Report_No_Later_Than_Date, Residence00State_Code, Residency_Status_Code,
Selected00ASC_Code, Selected00career_Pointer_Code, Selected00ewi_Option_Code,
Selected_Date, Selected_Quota_Year, Thesis_Diss_Required_Code,
Thesis_Program_Complete_Date, Type_Degree_Code, Input_Date, Login_Name,
Ewi00occupation_Series_Code, Scholarship_Type_Code, DLI_Language_Code,
Af_Acad_Sponsor_Dpt_Code, No_Cost_Indicator, Esp_Code, Book_Qtrs_Paid,
DLI_Entry_Date, School_In_Civ_Ins_Code, Remarks, Advance_Flag.

Class_Leader: SSAN (Pk) (Fk), Leader_Code (Pk), Program_Graduation00Term_Code, Class_Code,
Program_Year_Prefix

Degree_Awards: AFIT_Degree_Code (Pk), SSAN (Pk) (Fk), Career_Pointer_Code, Grad_Status_Code,
Pse_Code, Grade_Rank_Abbrev, Name_Prefix, Name_Suffix, First_Name, Last_Name,
Middle_Initial, Birth_Date, Sex_Code, Race_Code, Marital_Status_Code, Religion_Code,
Blue_Chip_Indicator, Aka_Fname, Aka_Lname, Prior_AFIT_Months, Tafms_Date,
Ethnic_Group_Code, Aero_Rating_Code, Manning_Code, Deros_Date, Separation_Date,
Commission_Code, Grade_Rank_Date, Citizenship00Country_Code, Department_Code,
Duty_Title, Duty_Phone, Duty_Area_Code, Badge_Number, Academic_Action_Code,
Overdue_Indicator, Classification_Code, Part_Record_Indicator, Admin_Hold_Indicator,
Major00ASC_Code, Academic_Specialty, Major00ed_Level_Code, Ed_Level, Program_Code,
Program, Program_Graduation00term_Code, Class_Code, Program_Year_Prefix,
Selected_Type_Code, Selected_Type, AFIT_Degree, Graduation00term_Code,
Graduation00quarter_Code, Graduation_Year_Prefix, Graduation_Date, Departure_Date,
Box_Number, Card_Number, Encoded_CARD_NUMBER, Library_Number, Locker_Number,
Admit_Date, Student_Sponsor_SSAN, Entry00term_Code, Entry00quarter_Code,
Entry_Year_Prefix, Admission_Type_Code, Admission_Action_Code, Gaining00AFSC_Code,
Faculty_Advisor_SSAN, Registration00department_Code, Program_Effective00term_Code,
Effective00quarter_Code, Effective_Year_Prefix, Leader_Code, Program_Section_Number,
Majors: Career_Pointer_Code (Pk), Major (Pk), SSAN (Pk) (Fk), Login_Name, Input_Date.
Name_History: SSAN (Pk) (Fk), Name_Change_Date (Pk), First_Name, Last_Name, Middle_Initial, Name_Suffix, Name_Prefix, Login_Name, Marital_Status_Code.
New_AFSC: SSAN (Pk) (Fk), AFSC_Code (Pk), Prefix.
OER_Data: SSAN (Pk) (Fk), Last_OER_Date, OER_Due_Date.
PCE_Grade: SSAN (Pk) (Fk), PCE_Course_Prefix (Pk), PCE_Course_Number (Pk), PCE_Course_Letter (Pk), PCE_Course_Year, PCEO0Grade_Code.
PCE_Std: SSAN (Pk) (Fk), PCE_Stay_Begin_Date (Pk), PCE_Stay_End_Date, PCEO0Billing_Code, MAJCOM_Code.
Person: SSAN (Pk), Grade_Rank_Abbrev, Name_Prefix, Name_Suffix, First_Name, Last_Name, Middle_Initial, Birth_Date, Sex_Code, Race_Code, Marital_Status_Code, Religion_Code, Blue_Chip_Indicator, Aka_Fname, Aka_Lname, Prior_AFIT_Months, Tafsms_Date, Ethnic_Group_Code, Aero_Rating_Code, Manning_Code, Deros_Date, Separation_Date, Commission_Code, Grade_Rank_Date, Citizenship00country_Code, Department_Code, Duty_Title, Duty_Phone, Duty_Area_Code, Badge_Number, Branch_Service_Code, Login_Name, Input_Date, Duty_Phone_Ext.
Personnel: SSAN (Pk) (Fk), Personnel00Department_Code, Personnel_Hire_Date, Personnel_Duty_Title, Phone_Number.
PHD: SSAN (Pk) (Fk), PHD_Major_Remark, PHD_Minor_Remark.
Planes_Flown: SSAN (Pk) (Fk), Plane_Name.
Program_History: Program_Code (Pk), Program_Graduation00term_Code (Pk), Program_Effective00term_Code (Pk), SSAN (Pk) (Fk), Input_Date, Faculty_Advisor_SSAN, Class_Code, Program_Year_Prefix, Ed_Level_Code, ASC_Code, Login_Name, Effective00quarter_Code, Effective_Year_Prefix, Career_Pointer_Code, AFIT_Degree_Code.
Program_STD_Section: SSAN (Pk) (Fk), Section_Number (Pk).
Rank_History: SSAN (Pk) (Fk), Grade_Rank_Abbrev (Pk), Grade_Rank_Date, Login_Name, Input_Date, Manning_Code, Branch_Service_Code.
Recall_Roster: SSAN (Pk) (Fk), Home_Phone_Number, Next_In_Chain_SSAN.
Registration_Verification: SSAN (Pk) (Fk), Term_Code, Quarter_Code, Year_Prefix, Registration_Notice.
Resident_Student: SSAN (Pk) (Fk), Academic_Action_Code, Overdue_Indicator, Classification_Code, Part_Record_Indicator, Admin_Hold_Indicator, Major00ASC_Code, Major00ed_Level_Code, Program_Code, Program_Graduation00term_Code, Class_Code, Program_Year_Prefix, Selected_Type_Code, AFIT_Degree_Code, Graduation00term_Code, Graduation00quarter_Code, Graduation_Year_Prefix, Grad_Status_Code, Departure_Date, Box_Number, Card_Number, Encoded_Card_Number, Library_Number, Locker_Number, Admit_Date, Student_Sponsor_SSAN, Entry00term_Code, Entry00quarter_Code, Entry_Year_Prefix, Admission_Type_Code, Admission_Action_Code, Career_Pointer_Code, Gaining00AFSC_Code, Faculty_Advisor_SSAN, Registration00department_Code, Program_Effective00term_Code, Effective00quarter_Code, Effective_Year_Prefix, Leader_Code, Program_Section_Number, Gain00MAJCOM_Abbrev, Gain00duty_Station, Losing00MAJCOM_Abbrev, Pse_Code.
Section_Leaders: Leader_Code (Pk), SSAN (Pk) (Fk), Program_Code, Class_Code, Program_Section_Number.
Selected_Comments: SSAN (Pk) (Fk), Selected_Comment.
Selected_Projection: SSAN (Pk) (Fk), Gain00MAJCOM_Code, Gain00MAJCOM_Abbrev, Gain_MAJCOM_Supervisor, Gain_MAJCOM_Supervisor_Phone, Gain_MAJCOM_DSN_Prefix, Gain_MAJCOM00department_Code, Position_Number_Projecteed.
Selected_Student: SSAN (Pk) (Fk), Selected00ASC_Code, Selected00ed_Level_Code, Selected_Quota_Year, Selected00ewi_Option_Code, Pca_Indicator, Quota_Program_Code, List_Number, Report,No_Earlier_Than_Date, Report,No_Later_Than_Date, Projected_Start_Date, Assign_Avail_Date, Assign_Reason_Code, MAJCOM_Abbrev, Projected_Entry_Class, Selected_Date, Input_Date, Login_Name, Selected_Type_Code, Selected00AFSC_Code, Reselection_Code, MPC_School_Code, Pse_Code, Assign00department_Code.
Selected_Student_Archive: Selected_Quota_Year (Pk), SSAN (Pk) (Fk), Selected00ASC_Code,
Selected00ed_Level_Code, Selected00ewi_Option_Code, Pca_Indicator, Quota_Program_Code,
List_Number, Report_No_Earlier_Than_Date, Report_No_Later_Than_Date,
Projected_Start_Date, Assign_Avail_Date, Assign_Reason_Code, MAJCOM_Abbrev,
Projected_Entry_Class, Selected_Date, Input_Date, Login_Name, Selected_Type_Code,
Selected00AFSC_Code, Gain00MAJCOM_Code, Gain00MAJCOM_Abbrev,
Gain_MAJCOM_Supervisor, Gain_MAJCOM_DSN_Prefix, Gain_MAJCOM_Supervisor_Phone,
Gain_MAJCOM0department_Code, Position_Number_Projected, Reselection_Code, Pse_Code,
MPC_School_Code, Assign00department_Code.

Sponsors_Country_Prefer: SSAN (Pk) (Fk), Ce, Preferred00country_Code (Pk).

Spouse_Info: SSAN (Pk) (Fk), Spouse_Birth_Date, Spouse_Fname, Spouse_Lname,
Spouse__At_AFIT_Indicator, Spouse_Nickname, Spouse_In_Military_Indicator,
Spouse_Occupation, Spouse_Remarks.

Student: SSAN (Pk) (Fk), Last_Name, First_Name, Department_Code, Grade_Rank_Abbrev,
Program_Code, Class_Code, GradTerm, Date_Entered, Revision_Date.

Student_Address: SSAN (Pk) (Fk), Address_Type (Pk), Address_Type_Desc, Address_Line_1,
Address_Line_2, City, State, ZIP, ZIP_Ext.

Student_Courses: Course_Prefix_Code (Pk) (Fk), Course_Number (Pk) (Fk), Course_Section (Pk) (Fk),
SSAN (Pk) (Fk), Term_Code, Hours, Grade, Calculated_Field.

Student_Duty_History: Duty_Sequence_Number (Pk), SSAN (Pk) (Fk), Duty_Title, Duty00AFSC_Code,
Duty_Organization, Duty_Station, Duty_Assigned_Date, Login_Name.

Student_Sequences: Program_Sequence_Code (Pk), SSAN (Pk) (Fk).

Term_Entry: SSAN (Pk) (Fk), Entry00term_Code, Entry00quarter_Code, Entry_Year_Prefix,
Admission_Type_Code, Admission_Action_Code.

Term_Entry_History: SSAN (Pk) (Fk), Entry00term_Code, Entry00quarter_Code, Entry_Year_Prefix,
Admission_Type_Code, Admission_Action_Code.

Test_Scores: Test_Type_Code (Pk), SSAN (Pk) (Fk), Test_Taken_Date, Test_Score, Login_Name,
Input_Date.

TDY_Attendees: SSAN (Pk) (Fk), Left_For_TDY_Date (Pk), Returned_From_TDY_Date, 
TDY_Destination_Code, TDY_Purpose.

Thesis_Diss_Book_Allowance: SSAN (Pk) (Fk), Allowance_Code, ASC_Code, ED_Level_Code.

Transcript_SSANS_105643: SSAN (Pk) (Fk), Last_Name, First_Name, Middle_Initial, Name_Suffix,
Grade_Rank_Abbrev, Program_Code, Selected_Type_Code, Table_Indicator.

Transcript_Sent: SSAN (Pk) (Fk), Transcript_Sent_Date (Pk), Num_Transcript_Sent.

Transfer_Transcript: SSAN (Pk) (Fk), Course_Prefix_Code (Pk), Course_Number (Pk), 
Transfer_Course_Prefix (Pk), Transfer_Course_Number (Pk), Course_Section,
AFIT00Credit_Hours, Earned_Hours_Indicator, GPA_Indicator, Transfer_Credit_Hours,
Transfer00Grade_Code, Transfer_Start_Date, Transfer_End_Date, MPC_School_Code,
Soche_Indicator, Tmstr00Term_Code, Tmstr00Quarter_Code, Tmstr_Year_Prefix,
Tmstr00Career_Pointer_Code, Transfer_Course_Title, Transcript_Course_Title.

Wait_List: SSAN (Pk) (Fk), Course_Prefix_Code (Pk), Course_Number (Pk), Course_Section (Pk).

Waived_Course: Waived00course_Prefix_Code (Pk), Waived00course_Number (Pk),
Course_Prefix_Code (Pk) (Fk), Course_Number (Pk) (Fk), Course_Section (Pk) (Fk), SSAN (Pk)
(Fk), Waived00grade_Code, Waived_Date.
Appendix B: The Entity Relationship Diagram

Figure 11: ER diagram (Person)
Figure 12: ER diagram (Resident Student 1)
Figure 13: ER diagram (Resident Student 2)
Figure 14: ER diagram (Resident Student 3)
Appendix C: The Object Model

Figure 15: Person's Object Model

This model does not show all the associations related to this object
This model does not show all the associations related to this object

Figure 16: Person Object Model (cont.)
This model does not show all the associations related to this object

Figure 17: Resident Student Object Model
This model does not show all the associations related to this object

Figure 18: Resident Student Object Model (cont.)

Figure 19: Grade History Association
Figure 20: Grade Change History Association

Figure 21: Student Courses Association

Figure 22: Dropped Course Association

Figure 23: Waived Course Association
Figure 24: Education History Association
Person Structure Definition

Object Name: Person
Object Number:
Object Description: General model of a person

Author: Maj Pedro Arthur Linhares Lima
Date: 03/25/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:

- **SSAN**: SSAN type Social Security Account Number.
- **lastname**: String Person’s last name.
- **firstname**: String Person’s first name.
- **name_prefix**: Prefix type Name prefix.
- **name_suffix**: Suffix type Name suffix.
- **middle_initial**: Character Name middle initial.
- **gender**: {male, female} Birth date.
- **birth_date**: Date type Birth date.
- **marital_status**: {single, married, divorced, widow} Person’s race.
- **race**: Race type Person’s religion.
- **religion**: Religion type Person’s ethnicity.
- **ethnic_group**: Ethnic type Badge number.
- **badge_number**: Badge number type Electronic mail address.
- **email_address**: String Academic education status.
- **academic_ed_status**: Academic ed type Duty title.
- **duty_title**: String Duty phone number.
- **duty_phone**: Phone type Person’s department.
- **department_code**: Department type Person’s country.
- **citizenship_country**: Country type Login name of who made changes.
- **login_name**: String Date of last change.
- **input_date**: Date type

Constraints:

Z Static Schema:

Let SSAN_TYPE be the set of all Social Security Account Numbers.
Let DATE_TYPE be the set of all possible dates.
Let PREFIX_TYPE be the set of all possible name prefixes.
Let SUFFIX_TYPE be the set of all possible name suffixes.
Let RACE_TYPE be the set of all possible races.
Let RELIGION_TYPE be the set of all possible religions.
Let ETHNIC_TYPE be the set of all possible ethnic groups.
Let BADGE_NUMBER_TYPE be the set of all possible badge numbers.
Let ACADEMIC_ED_TYPE be the set of all possible academic education types.
Let PHONE_TYPE be the set of all possible phone numbers.
Let DEPARTMENT_TYPE be the set of all possible departments.
Let COUNTRY_TYPE be the set of all possible country codes.
Person

SSAN : SSAN_TYPE
lastname : String
firstname : String
name_prefix : PREFIX_TYPE
name_suffix : SUFFIX_TYPE
middle_initial : Character
gender : {male, female}
birth_date : DATE_TYPE
race : RACE_TYPE
marital_status : {single, married, divorced, widow}
religion : RELIGION_TYPE
e-mail_address : String
academic_ed_status : ACADEMIC_ED_TYPE
ethnic_group : ETHNIC_TYPE
citizenship_country : COUNTRY_TYPE
department_code : DEPARTMENT_TYPE
duty_title : String
duty_phone : PHONE_TYPE
badge_number : BADGE_TYPE
login_name : String
input_date : DATE_TYPE
Military Structure Definition

Object Name: Military
Object Number:
Object Description: General model of a military

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: Person
Components: None
Context: None

Attributes:
- rank
  - Rank type
  - Military rank.
- branch
  - Branch type
  - Branch of service.
- date_of_rank
  - Date type
  - Date of rank.
- AFSC
  - AFSC type
  - AFSC code.
- date_of_commission
  - Date type
  - Date of commission.
- date_of_separation
  - Date type
  - Date of separation.
- manning_code
  - Manning type
  - Manning code.
- DEROS_date
  - Date type
  - DEROS date.
- duty_effective_date
  - Date type
  - Date of effective duty.
- aero_rating_code
  - Aero_rating type
  - Aero rating code.
- MAJCOM
  - MAJCOM type
  - MAJCOM code.
- base
  - Base type
  - Base code.
- blue_chip_indicator
  - Character
  - Blue chip indicator.
- NCO_indicator
  - Boolean
  - NCO indicator.
- MPC_code
  - MPC type
  - MPC code.
- recall_roster
  - Person pointer
  - Pointer to another military.

Constraints:

Z Static Schema:
- Let RANK_TYPE be the set of all possible rank types.
- Let BRANCH_TYPE be the set of all possible branch types.
- Let DATE_TYPE be the set of all possible dates.
- Let MANNING_TYPE be the set of all possible manning types.
- Let AERO_RATING_TYPE be the set of all possible aero rating types.
- Let MAJCOM_TYPE be the set of all possible MAJCOM types.
- Let BASE_TYPE be the set of all possible base types.
- Let MPC_TYPE be the set of all possible MPC types.
- Let PERSON_POINTER_TYPE be a pointer to a particular person.
Military

rank: RANK_TYPE
branch: BRANCH_TYPE
date_of_rank: DATE_TYPE
AFSC: AFSC_TYPE
date_of_commission: DATE_TYPE
date_of_separation: DATE_TYPE
manning_code: MANNING_TYPE
DEROS_date: DATE_TYPE
duty_effective_date: DATE_TYPE
aero_rating_code: AERO_RATING_TYPE
MAJCOM: MAJCOM_TYPE
base: BASE_TYPE
blue_chip_indicator: Character
NCO_indicator: Boolean
MPC_code: MPC_TYPE
recall_roster: PERSON_POINTER_TYPE
Student Structure Definition

Object Name: Student
Object Number:
Object Description: General model of student

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: Person
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Part_Record_Indicator</td>
<td>Boolean</td>
<td>If a student is resident or part time.</td>
</tr>
<tr>
<td>Ed_Level</td>
<td>EDUCATION_LEVEL_TYPE</td>
<td>Major education level.</td>
</tr>
<tr>
<td>Selected_Type</td>
<td>SELECTED_TYPE</td>
<td>School that has been selected.</td>
</tr>
<tr>
<td>Gaining_AFSC</td>
<td>AFSC_TYPE</td>
<td>The AFSC that he is going.</td>
</tr>
<tr>
<td>Gain_MAJCOM</td>
<td>MAJCOM_TYPE</td>
<td>The MAJCOM that he is going.</td>
</tr>
<tr>
<td>Gain_duty_Station</td>
<td>Duty_Station_TYPE</td>
<td>Duty station he is going.</td>
</tr>
<tr>
<td>Losing_MAJCOM</td>
<td>MAJCOM_TYPE</td>
<td>The MAJCOM that he is losing.</td>
</tr>
<tr>
<td>PSE</td>
<td>PSE_TYPE</td>
<td>The professional specialized educ.</td>
</tr>
<tr>
<td>Last_Year_Attended</td>
<td>YEARTYPE</td>
<td>The last year he attended a school.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

Let EDUCATION_LEVEL_TYPE be the set of all education level types.
Let SELECTED_TYPE be the set of all possible selected types.
Let AFSC_TYPE be the set of all possible AFSC types.
Let MAJCOM_TYPE be the set of all possible MAJCOM types.
Let DUTY_STATION_TYPE be the set of all possible duty stations.
Let PSE_TYPE be the set of all possible PSE types.
Let YEARTYPE be the set of all possible years.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Part_Record_Indicator</td>
<td>Boolean</td>
<td>If a student is resident or part time.</td>
</tr>
<tr>
<td>Ed_Level</td>
<td>EDUCATION_LEVEL_TYPE</td>
<td>Major education level.</td>
</tr>
<tr>
<td>Selected_Type</td>
<td>SELECTED_TYPE</td>
<td>School that has been selected.</td>
</tr>
<tr>
<td>Gaining_AFSC</td>
<td>AFSC_TYPE</td>
<td>The AFSC that he is going.</td>
</tr>
<tr>
<td>Gain_MAJCOM</td>
<td>MAJCOM_TYPE</td>
<td>The MAJCOM that he is going.</td>
</tr>
<tr>
<td>Gain_duty_Station</td>
<td>Duty_Station_TYPE</td>
<td>Duty station he is going.</td>
</tr>
<tr>
<td>Losing_MAJCOM</td>
<td>MAJCOM_TYPE</td>
<td>The MAJCOM that he is losing.</td>
</tr>
<tr>
<td>PSE</td>
<td>PSE_TYPE</td>
<td>The professional specialized educ.</td>
</tr>
<tr>
<td>Last_Year_Attended</td>
<td>YEARTYPE</td>
<td>The last year he attended a school.</td>
</tr>
</tbody>
</table>

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Resident Student Structure Definition

Object Name: Resident Student
Object Number: 
Object Description: General model of resident student

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: Student
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Academic_Action</td>
<td>Academic Action type</td>
</tr>
<tr>
<td>Overdue_Indicator</td>
<td>Boolean</td>
</tr>
<tr>
<td>Classification</td>
<td>Classification type</td>
</tr>
<tr>
<td>Major_ASC</td>
<td>ASC type</td>
</tr>
<tr>
<td>Program</td>
<td>Program type</td>
</tr>
<tr>
<td>Class</td>
<td>Class type</td>
</tr>
<tr>
<td>Program_Year</td>
<td>Year type</td>
</tr>
<tr>
<td>AFIT_Degree</td>
<td>Degree type</td>
</tr>
<tr>
<td>Career_Pointer_Code</td>
<td>Career pointer type</td>
</tr>
<tr>
<td>Departure_Date</td>
<td>Date type</td>
</tr>
<tr>
<td>Box_Number</td>
<td>Box type</td>
</tr>
<tr>
<td>Card_Number</td>
<td>Card type</td>
</tr>
<tr>
<td>Library_Number</td>
<td>Library_Number type</td>
</tr>
<tr>
<td>Locker_Number</td>
<td>Locker type</td>
</tr>
<tr>
<td>Student_Sponsor</td>
<td>Person pointer</td>
</tr>
<tr>
<td>Faculty_Advisor</td>
<td>Person pointer</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

Let ACADEMIC_ACTION_TYPE be the set of all possible academic actions.
Let CLASSIFICATION_TYPE be the set of all possible classification types.
Let ASC_TYPE be the set of all possible ASC types.
Let PROGRAM_TYPE be the set of all possible programs.
Let CLASS_TYPE be the set of all possible class types.
Let YEAR_TYPE be the set of all possible years.
Let DEGREE_TYPE be the set of all possible degree types.
Let CAREER_POINTER_TYPE be the set of all possible career pointer types.
Let DATE_TYPE be the set of all possible dates.
Let BOX_TYPE be the set of all possible boxes.
Let CARD_POINTER_TYPE be the set of all possible card pointer types.
Let LIBRARY_NUMBERTYPE be the set of all possible library numbers.
Let LOCKER_TYPE be the set of all possible locker numbers.
Let PSE_TYPE be the set of all possible PSE types.
Let PERSON_POINTER_TYPE be a pointer to a particular person.
<table>
<thead>
<tr>
<th>Resident student</th>
</tr>
</thead>
<tbody>
<tr>
<td>Academic_Action : ACADEMIC_ACTION_TYPE</td>
</tr>
<tr>
<td>Overdue_Indicator : Boolean</td>
</tr>
<tr>
<td>Classification : CLASSIFICATION_TYPE</td>
</tr>
<tr>
<td>Major_ASC : ASC_TYPE</td>
</tr>
<tr>
<td>Program : PROGRAM_TYPE</td>
</tr>
<tr>
<td>Class : CLASS_TYPE</td>
</tr>
<tr>
<td>Program_Year : YEAR_TYPE</td>
</tr>
<tr>
<td>AFIT_Degree : DEGREE_TYPE</td>
</tr>
<tr>
<td>Career_Pointer_Code : CAREER_POINTER_TYPE</td>
</tr>
<tr>
<td>Departure_Date : DATE_TYPE</td>
</tr>
<tr>
<td>Box_Number : BOX_TYPE</td>
</tr>
<tr>
<td>Card_Number : CARD_TYPE</td>
</tr>
<tr>
<td>Library_Number : LIBRARY_NUMBER_TYPE</td>
</tr>
<tr>
<td>Locker_Number : LOCKER_TYPE</td>
</tr>
<tr>
<td>Student_Sponsor : PERSON_POINTER_TYPE</td>
</tr>
<tr>
<td>Faculty_Advisor : PERSON_POINTER_TYPE</td>
</tr>
</tbody>
</table>
INTL-Student Structure Definition

Object Name: INTL-Student
Object Number:
Object Description: General model of INTL-student

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: Resident Student
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>WSCN</td>
<td>WSCN type</td>
<td>Work Sheet Control Number.</td>
</tr>
<tr>
<td>ITO</td>
<td>ITO type</td>
<td>ITO number.</td>
</tr>
<tr>
<td>Case_number</td>
<td>Case_number type</td>
<td>Case number.</td>
</tr>
<tr>
<td>DLI_request_indicator</td>
<td>Boolean</td>
<td>If student has to attended DLI.</td>
</tr>
<tr>
<td>DLI_indicator</td>
<td>Boolean</td>
<td>If student attended DLI.</td>
</tr>
<tr>
<td>evaluation_req_date</td>
<td>Date type</td>
<td>Date evaluation was requested.</td>
</tr>
<tr>
<td>requested_program</td>
<td>Program type</td>
<td>Requested student’s Program.</td>
</tr>
<tr>
<td>eval_forward_date</td>
<td>Date type</td>
<td>Evaluation forward date.</td>
</tr>
<tr>
<td>forward_to_dept</td>
<td>Department type</td>
<td>Department it was forwarded.</td>
</tr>
<tr>
<td>Eval_returned_date</td>
<td>Date type</td>
<td>Date the evaluation returned.</td>
</tr>
<tr>
<td>admission_status</td>
<td>Admission type</td>
<td>Admission status code.</td>
</tr>
<tr>
<td>eval_remarks</td>
<td>String</td>
<td>Evaluation remarks.</td>
</tr>
<tr>
<td>country_notified_date</td>
<td>Date type</td>
<td>Date the country was notified.</td>
</tr>
<tr>
<td>AFSAT_notified_date</td>
<td>Date type</td>
<td>Date that AFSAT was notified.</td>
</tr>
<tr>
<td>AFSAT_quota_indicator</td>
<td>Boolean</td>
<td>If student fills a country’s quota.</td>
</tr>
<tr>
<td>first_sponsor</td>
<td>Person pointer</td>
<td>Pointer to sponsor.</td>
</tr>
<tr>
<td>second_sponsor</td>
<td>Person pointer</td>
<td>Pointer to sponsor.</td>
</tr>
<tr>
<td>source_of_funds</td>
<td>Funds type</td>
<td>Source of funds code.</td>
</tr>
<tr>
<td>AFSAT_country</td>
<td>Country type</td>
<td>The home country of a student.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

Let WSCN_TYPE be the set of all possible WSCN numbers.
Let ITO_TYPE be the set of all possible ITO numbers.
Let CASE_NUMBER_TYPE be the set of all possible case numbers.
Let DATE_TYPE be the set of all possible dates.
Let PROGRAM_TYPE be the set of all possible programs.
Let DEPARTMENT_TYPE be the set of all possible department types.
Let ADMISSION_TYPE be the set of all possible admission types.
Let PERSON_POINTER_TYPE be a pointer to a particular person.
Let FUNDS_TYPE be the set of all possible funds types.
Let COUNTRY_TYPE be the set of all possible country types.
INTL-student

WSCN : WSCN_TYPE
ITO : ITO_TYPE
Case_number : CASE_NUMBER_TYPE
DLI_request_indicator : Boolean
DLI_indicator : Boolean
evaluation_req_date : DATE_TYPE
requested_program : PROGRAM_TYPE
eval_forward_date : DATE_TYPE
forward_to_dept : DEPARTMENT_TYPE
Eval_returned_date : DATE_TYPE
admission_status : ADMISSION_TYPE
evalRemarks : String
country_notified_date : DATE_TYPE
AFSAT_notified_date : DATE_TYPE
AFSAT_quota_indicator : Boolean
first_sponsor : PERSON_POINTER_TYPE
second_sponsor : PERSON_POINTER_TYPE
source_of_funds : FUNDS_TYPE
AFSAT_country : COUNTRY_TYPE
Address Structure Definition

Object Name: Address
Object Number: 
Object Description: General model of address

Author: Maj Pedro Arthur Linhares Lima
Date: 03/25/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>address_type</td>
<td>Address type</td>
<td>Address type.</td>
</tr>
<tr>
<td>address</td>
<td>String</td>
<td>Address.</td>
</tr>
<tr>
<td>city</td>
<td>String</td>
<td>City.</td>
</tr>
<tr>
<td>state</td>
<td>State type</td>
<td>State.</td>
</tr>
<tr>
<td>country</td>
<td>Country type</td>
<td>Country.</td>
</tr>
<tr>
<td>zipcode</td>
<td>Zip type</td>
<td>Zip code.</td>
</tr>
<tr>
<td>phone</td>
<td>Phone type</td>
<td>Phone number.</td>
</tr>
<tr>
<td>address_effective_date</td>
<td>Date type</td>
<td>Date of effective address.</td>
</tr>
<tr>
<td>login_name</td>
<td>String</td>
<td>Login name of who made changes.</td>
</tr>
<tr>
<td>login_date</td>
<td>Date type</td>
<td>Date of last change.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

Let ADDRESS_TYPE be the set of all possible address types.
Let STATE_TYPE be the set of all possible states.
Let COUNTRY_TYPE be the set of all possible country codes.
Let ZIP_TYPE be the set of all possible zip codes.
Let PHONE_TYPE be the set of all phone Numbers.
Let DATE_TYPE be the set of all possible dates.

```
Address
address_type : String
address : String
city : String
state : STATE_TYPE
zipcode : ZIP_TYPE
country : COUNTRY_TYPE
phone : PHONE_TYPE
address_effective_date : DATE_TYPE
login_name : String
login_date : DATE_TYPE
```
Dependent Information Structure Definition

Object Name: Dependent Information
Object Number:
Object Description: General model of dependent information

Author: Maj Pedro Arthur Linhares Lima
Date: 03/26/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>number_children</td>
<td>Integer</td>
<td>Number of student’s children.</td>
</tr>
<tr>
<td>single_dep_chldrn</td>
<td>Character</td>
<td>Indicator if single with children.</td>
</tr>
<tr>
<td>deps_at_AFIT</td>
<td>Character</td>
<td>Indicator if dependents at AFIT.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

```
Dependent Information
number_children : Integer
single_dep_chldrn: Character
deps_at_AFIT: Character
```
Education History Structure Definition

Object Name: Education History
Object Number:
Object Description: General model of education history

Author: Maj Pedro Arthur Linhares Lima
Date: 03/26/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:
- MPCSchoolCode: MPC type
- EdLevelCode: Ed_level type
- Type_Degree_Code: Type_degree type
- ASC_Code: ASC type
- Quality_Points: Integer
- Total_Credit_Hours: Integer
- Method_Of_Obtainment: Method_of_obt_type
- Academic_Ed_Status: Academic_ed_type
- Last_Year_Attended: Date type
- ABET_Accredited: Boolean
- Ed_History_Remarks: String
- Work_ID_Processed: Work_ID type
- Login_Name: String
- Input_Date: Date type
- Operators_Initials: String
- Transcript_Career_Pointer: Transcript type
- Duty_Location: Duty_location type
- Degree_Cum_Gpa: Integer
- Degree_Title: Degree type

Constraints:

Z Static Schema:
Let MPC_TYPE be the set of all possible MPC types.
Let ED_LEVEL_TYPE be the set of all possible education level types.
Let TYPE_DEGREE_TYPE be the set of all possible degree types.
Let ASC_TYPE be the set of all possible ASC types.
Let DATE_TYPE be the set of all possible dates.
Let METHOD_OF_OBT_TYPE be the set of all possible method of obtained types.
Let ACADEMIC_ED_TYPE be the set of all possible academic education types.
Let WORK_ID_TYPE be the set of all possible work ID types.
Let TRANSSCRIPT_TYPE be the set of all possible transcript types.
Let DUTY_LOCATION_TYPE be the set of all possible duty location types.
Let DEGREE_TYPE be the set of all possible degree types.
Education History

MPC_School_Code : MPC_TYPE
Ed_Level_Code : ED_LEVEL_TYPE
Type_Degree_Code : TYPE_DEGREE_TYPE
ASC_Code : ASC_TYPE
Quality_Points : Integer
Total_Credit_Hours : Integer
Method_Of_Obtainment : METHOD_OF_OBT_TYPE
Academic_Ed_Status : ACADEMIC_ED_TYPE
Last_Year_Attended : DATE_TYPE
ABET_Accredited : Boolean
Ed_History_Remarks : String
Work_ID_Processed : WORK_ID_TYPE
Login_Name : String
Input_Date : DATE_TYPE
Operators_Initials : String
Trnscript_Career_Pointer : TRANSCRIPT_TYPE
Duty_Location : DUTY_LOCATION_TYPE
Degree_Cum_Gpa : Integer
Degree_Title : DEGREE_TYPE
Emergency Data Structure Definition

Object Name: Emergency Data
Object Number:
Object Description: General model of emergency data

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Contact_Fname</td>
<td>String</td>
<td>Emergency contact first name.</td>
</tr>
<tr>
<td>Contact_Lname</td>
<td>String</td>
<td>Emergency contact last name.</td>
</tr>
<tr>
<td>Relation</td>
<td>Relation</td>
<td>The relationship with a person.</td>
</tr>
<tr>
<td>Address</td>
<td>String</td>
<td>Home address.</td>
</tr>
<tr>
<td>City</td>
<td>String</td>
<td>Home city.</td>
</tr>
<tr>
<td>State</td>
<td>State</td>
<td>Home state.</td>
</tr>
<tr>
<td>Zipcode</td>
<td>Zip</td>
<td>Home zip code.</td>
</tr>
<tr>
<td>Country</td>
<td>Country</td>
<td>Person’s country.</td>
</tr>
<tr>
<td>Phone</td>
<td>Phone</td>
<td>Home phone number.</td>
</tr>
<tr>
<td>Firm_Name_Office</td>
<td>Office</td>
<td>Person’s office symbol.</td>
</tr>
<tr>
<td>Additional_Address</td>
<td>String</td>
<td>Additional address information.</td>
</tr>
<tr>
<td>Street_Address</td>
<td>String</td>
<td>Additional street information.</td>
</tr>
<tr>
<td>Address_Room_Type</td>
<td>Room</td>
<td>Additional Room type.</td>
</tr>
<tr>
<td>Address_Room_Number</td>
<td>String</td>
<td>Additional room number.</td>
</tr>
<tr>
<td>Street_Type_Code</td>
<td>Street</td>
<td>Additional street type.</td>
</tr>
<tr>
<td>Revision_Name</td>
<td>String</td>
<td>Name of who made the revision.</td>
</tr>
<tr>
<td>Revision_Date</td>
<td>Date</td>
<td>Revision date.</td>
</tr>
<tr>
<td>Login_Name</td>
<td>String</td>
<td>Login name of who made changes.</td>
</tr>
<tr>
<td>Login_Date</td>
<td>Date</td>
<td>Date of last change.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:

Let RELATION_TYPE be the set of all possible relationship types.
Let STATE_TYPE be the set of all possible states.
Let ZIP_TYPE be the set of all possible zip codes.
Let COUNTRY_TYPE be the set of all possible country codes.
Let PHONE_TYPE be the set of all phone numbers.
Let OFFICE_TYPE be the set of all possible offices.
Let ROOM_TYPE be the set of all possible room types.
Let STREET_TYPE be the set of all possible street types.
Let DATE_TYPE be the set of all possible dates.
**Emergency data**

Contact_Fname : String  
Contact_Lname : String  
Relation : RELATION_TYPE  
Address : String  
City : String  
State : STATE_TYPE  
Zipcode : ZIP_TYPE  
Country : COUNTRY_TYPE  
Phone : PHONE_TYPE  
Firm_Name_Office : OFFICE_TYPE  
Additional_Address : String  
Street_Address : String  
Address_Room_Type : ROOM_TYPE  
Address_Room_Number : String  
Street_Type_Code : STREET_TYPE  
Revision_Name : String  
Revision_Date : DATE_TYPE  
Login_Name : String  
Login_Date : DATE_TYPE
Spouse Information Structure Definition

Object Name: Spouse Information
Object Number:
Object Description: General model of spouse information

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Birth_Date</td>
<td>Date type</td>
<td>Spouse birth date.</td>
</tr>
<tr>
<td>Fname</td>
<td>String</td>
<td>Spouse first name.</td>
</tr>
<tr>
<td>Lname</td>
<td>String</td>
<td>Spouse last name.</td>
</tr>
<tr>
<td>Spouse_At_AFIT</td>
<td>Boolean</td>
<td>If spouse came with him/her.</td>
</tr>
<tr>
<td>Nickname</td>
<td>String</td>
<td>Spouse nickname.</td>
</tr>
<tr>
<td>Spouse_In_Military</td>
<td>Boolean</td>
<td>If spouse in military service.</td>
</tr>
<tr>
<td>Occupation</td>
<td>String</td>
<td>Spouse occupation.</td>
</tr>
<tr>
<td>Remarks</td>
<td>String</td>
<td>Remarks.</td>
</tr>
</tbody>
</table>

Constraints:

Z Static Schema:
Let DATE_TYPE be the set of all possible dates.

---

Spouse Information

\[
\begin{align*}
Birth\_Date & : DATE\_TYPE \\
Fname & : String \\
Lname & : String \\
Spouse\_At\_AFIT & : Boolean \\
Nickname & : String \\
Spouse\_In\_Military & : Boolean \\
Occupation & : String \\
Remarks & : String
\end{align*}
\]
Student Duty History Structure Definition

Object Name: Student Duty History
Object Number:
Object Description: General model of student duty history

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:
- Title: Title type Job title.
- AFSC: AFSC type Duty AFSC.
- Organization: String Where the student works.
- Duty_Station: Duty_Station type Student’s duty station.
- Assigned_Date: Date type Date the student was assigned.
- Sequence_Number: Integer Sequence number.
- Login_Name: String Login name.

Constraints:

Z Static Schema:
Let TITLE_TYPE be the set of all possible titles.
Let AFSCTYPE be the set of all possible AFSC types.
Let DUTYSTATIONTYPE be the set of all possible duty stations.
Let DATETYPE be the set of all possible dates.

<table>
<thead>
<tr>
<th>Student duty history</th>
</tr>
</thead>
<tbody>
<tr>
<td>Title : TITLE_TYPE</td>
</tr>
<tr>
<td>AFSC : AFSCTYPE</td>
</tr>
<tr>
<td>Organization : String</td>
</tr>
<tr>
<td>Duty_Station : DUTY_STATION_TYPE</td>
</tr>
<tr>
<td>Assigned_Date : DATE_TYPE</td>
</tr>
<tr>
<td>Sequence_Number : Integer</td>
</tr>
<tr>
<td>Login_Name : String</td>
</tr>
</tbody>
</table>
Test Scores Structure Definition

Object Name: Test Scores
Object Number:
Object Description: General model of test scores

Author: Maj Pedro Arthur Linhares Lima
Date: 03/27/96
History: Thesis

Superclass: None
Components: None
Context: None

Attributes:
- Test Type: Test type
- Taken Date: Date type
- Score: Integer
- Input Date: Date type
- Login Name: String

Constraints:

Z Static Schema:
Let TEST_TYPE be the set of all possible test types.
Let DATE_TYPE be the set of all possible dates.

Test scores
Test Type : TEST_TYPE
Taken Date : DATE_TYPE
Score : Integer
Input Date : DATE_TYPE
Login_Name : String
Appendix D: The Functional Model

Figure 25: STARS Application Level 0 DFD
Figure 26: Perform Action Level 1 DFD
Figure 27: Menu Option Level 2 DFD
Figure 28: Handle Selection Level 2 DFD
Figure 29: Perform Selection Level 3 DFD
Figure 30: Perform Insertion Level 4 DFD
Figure 31: Perform Update Level 4 DFD
Figure 32: Perform Deletion Level 4 DFD
Appendix E: Implementation of the Object Model

Figure 33: Table's relationship in STARS Database
Figure 34: Person's Class
Figure 35: Military's Class
Figure 36: Military_student's Class
Figure 37: Military_resident_student’s Class
Figure 38: Military INTL student's Class
Figure 39: Civilian’s Class
Figure 40: Civilian_student's Class
Figure 41: Civilian_resident_student's Class
Figure 42: Civilian_INTL_student's Class
Figure 43: Address' Class
Appendix F: List of Abbreviations

AFIT - Air Force Institute of Technology

AFITSIS - Air Force Institute of Technology Student Information System

AFIT/SC - Air Force Institute of Technology Communication Computer System

CASE - Computer-Aided Software Engineering

DBA - Database Administrator

DBMS - Database Management System

DFD - Data Flow Diagram

ER - Entity Relationship

OODBMS - Object-Oriented Database Management System

RDB - Relational Database

RDBMS - Relational Database Management System

SQL - Structured Query Language

STARS - Student Tracking and Registration System


Vita

Major Pedro Arthur Linhares Lima was born in Rio de Janeiro, Brazil.

He entered the Air Force Preparatory School of Cadets (EPCAR) in Barbacena, Minas Gerais in 1975, and attended the Brazilian Air Force Academy, where he was graduated in December of 1981. He entered the Catholic University of Rio de Janeiro (PUC-RJ), where was awarded the degree of Bachelor in Systems Analysis in June of 1984. His first assignment was at the Air Force Computer Center of Rio de Janeiro (CCA-RJ), where he worked as a systems analyst for the Flight's Statistics System Project. In January of 1989 he was assigned to Staff and War College (ECEMAR), where he worked as a systems analyst for the War Games Project. In January of 1992 he was assigned to Air Force Computer Science and Statistics Department (DIRINFE). In June of 1994 Major Pedro Lima entered the Air Force Institute of Technology as a Master candidate in computer science.
This research proposes and evaluates a methodology for reengineering a relational database to an object-oriented database. We applied this methodology to reengineering the Air Force Institute of Technology Student Information System (AFITSIS) as our test case. With this test case, we could verify the applicability of the proposed methodology, especially because AFITSIS comes from an old version of Oracle RDBMS. We had the opportunity to implement part of the object model using an object-oriented database, and we present some peculiarities encountered during this implementation. The most important result of this research is that it demonstrated that the proposed methodology can be used for reengineering an arbitrarily selected relational database to an object-oriented database. It appears that this approach can be applied to any relational database.