Learning One Subprocedure Per Lesson

Kurt VanLehn
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Sierra is a program that learns procedures incrementally from examples. It learns by completing explanations. Whenever the current procedure is inadequate to explain (parse) the current example (a string of actions), Sierra formulates a new subprocedure whose instantiation completes the explanation. This technique is one solution to the problem of acquiring the domain theory required by learning-by-explanation systems. The key to its success lies in supplying a small amount of extra information with the examples. Instead of giving the learner an set of examples, under which conditions correct learning is provably impossible, it is given a sequence of "lessons," where a lesson is a set of examples, possibly a singleton set, that is guaranteed to introduce only one subprocedure. This permits unbiased learning, i.e., learning without a priori, heuristic preferences concerning the outcome.
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Abstract

Sierra is a program that learns procedures incrementally from examples, where an example is a sequence of actions. Sierra learns by completing explanations. Whenever the current procedure is inadequate for explaining (parsing) the current example, Sierra formulates a new subprocedure whose instantiation completes the explanation (parse tree). The key to Sierra's success lies in supplying a small amount of extra information with the examples. Instead of giving it a set of examples, under which conditions correct learning is provably impossible, it is given a sequence of "lessons," where a lesson is a set of examples that is guaranteed to introduce only one subprocedure. This permits unbiased learning, i.e., learning without a priori, heuristic preferences concerning the outcome.
Table of Contents

1 Introduction
   1.1. The particular learning task solved by Sierra 2
   1.2. Why lesson sequences simplify procedure induction 4
   1.3. How Sierra works 8
2 The representation of procedures 10
3 Disjunction: an inherent problem for induction 15
4 Subprocedures 20
5 Skeleton induction
   5.1. Prior solutions to the skeleton induction problem 28
6 Pattern induction 29
7 Function induction
   7.1. Prior solutions to the invisible objects problem 33
   7.2. The show-work felicity condition 36
   7.3. Parallels between the invisible objects problem and the disjunction problem 36
8 Task generality 39
9 Concluding remarks
   9.1. Some speculations on manner constraints and felicity conditions 42
   9.2. Summary 43
   9.3. Discussion 44
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1. Introduction

Much research in machine learning has concentrated on induction, i.e., learning from examples. Understanding induction is certainly one of the great intellectual challenges of our times. Induction stands at the center of both the psychology of learning and the philosophy of science.

More recently, induction has been heralded as a potential solution to the knowledge acquisition problem of expert systems. However, AI has not had much success at applying induction to practical problems. The difficulty is this: in practical settings, only a finite set of examples can be presented to the learner, but the knowledge representation language usually has enough expressive power that there are always infinitely many representable generalizations that are consistent with the examples that the learner has seen. (The formal results that justify this assertion will be reviewed later.) That is, the examples just do not contain enough information to correctly identify the target generalization. Consequently, the learner must guess. AI research on induction has consisted largely of studying the efficacy and domain independence of various heuristics for guessing. In this respect, the induction problem is quite different than other kinds of AI problems, e.g., recognition or synthesis, where correct answers are computable in principle, but non-AI techniques would take astronomically long to compute them.

The inherent uncertainty of induction suggests studying forms of quasi-inductive learning, where a teacher gives some extra information to the learner in addition to the examples. Winston argues for "the importance of good training sequences prepared by good teachers. I think it is reasonable to believe that neither machines nor children can be expected to learn much without them [47, pg. 6]." As Winston [48, 49, 50] has pointed out, there is a spectrum defined by how much extra information is provided. On the one end of the spectrum is induction, where all the learner receives is examples. On the other end is learning by being told, where the learner is given a complete description of the target generalization in some language.
The kind of learning studied here falls closer to the induction end of the spectrum end than the learning-by-being-told end, because very little extra information is supplied. It could be called learning from lesson sequences because the extra information given to the learner is embedded in the way that the examples are partitioned into lessons and the way the lessons are sequenced. In the last section of this article a variant of learning from lesson sequences will be discussed wherein lessons are omitted and the example sequence alone carries all the extra information.

One of the most important areas for applying induction is in the learning of procedures. Procedure learning is the central problem in programming by examples [1 7 8 40 41], in psychological modelling of skill acquisition [2 3 33 44 45], and in automating protocol analyses [6]. Machine learning of procedures has been suggested as one way to solve the knowledge acquisition problem for expert systems [31], and as a technique for modelling students in intelligent tutoring systems [27, 25]. Undoubtedly, there are many other applications for procedure learning systems that have not yet been studied.

An obvious testbed for the new approach of learning from lesson sequences is a system to learn procedures. This article presents Sierra, a system that learns procedures from lesson sequences.

1.1 The particular learning task solved by Sierra

Sierra was built for a specific application, a psychological study of the acquisition of mathematical skills, arithmetic in particular [44, 45]. Although it would make the reader's job easier if Sierra's techniques were displayed in a toy problem domain, this article presents an unsimplified, accurate picture. First, the kinds of examples given to Sierra will be described followed by a description of lesson sequences.

Two general kinds of procedure induction problems have been addressed in the literature. The harder one is learning a procedure from input-output pairs [1]. The one studied here is learning from action sequences. It is assumed that the agent that executes procedures is like a human or a robot in that its procedures manipulate both (1) an external world that all agents have access to and (2) an internal state which...
is private. The internal state might include a stack, for instance. An action sequence consists of a sequence of state changes to the external world (or equivalently, a sequence of world states). The learner cannot see the internal state of the teacher during an action sequence. Action sequences are "examples" of the target procedure's execution. The induction task is to infer a procedure from such examples. There are two kinds of examples: positive and negative. A positive example is an action sequence that an induced procedure should generate when it is run on the problem that appears as the initial state in the action sequence. A negative example is an action sequence that an induced procedure should not generate.

Sierra's input is an ordered sequence of lessons, where a lesson is an unordered set of examples. Lessons contain only positive examples. Each lesson is marked with a single bit. The bit is 1 if the lesson is a "normal" lesson, and 0 if it is an "optimization" lesson. In a moment, the semantics of lessons and their marks will be described. The point to note here is that the extra information given to Sierra over and beyond the set of examples consists only of (1) the partition of the examples into lessons, (2) the ordering of the lessons, and (3) the binary mark on lessons. Although this is very little extra information, it vastly simplifies Sierra's induction task. The next section discusses why.

Sierra learns procedures incrementally. Each lesson builds on the procedure learned in a previous lesson. This can be best illustrated with a familiar procedure, such as the procedure for ordinary multicolumn subtraction. Figure 1-1 shows a lesson sequence for subtraction. There are six lessons. Although a typical Sierra lesson has about five examples, the figure shows only the first example from each lesson. Each example is shown as a sequence of states. Figure 1-2 shows a corresponding sequence of procedures, where the procedures are sketched as Augmented Transition Nets (ATNs). The procedures correspond to the lessons as follows: P0 is induced from lesson L0; P1 from P0 and L1; P2 from P1 and L2, etc. The last procedure

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2 A larger vocabulary of markers could be used, e.g., 1 indicates a lesson in negative examples.
P5 is a complete correct subtraction procedure.  

1.2 Why lesson sequences simplify procedure induction  
An earlier discussion put learning from lesson sequences on the same scale as learning by being told. This may seem strange. Information encoded in the lesson sequence does not seem like a linguistic expression. However, it functions in exactly the same way. Linguistic expressions have meaning for the learner only under interpretation. The conventions governing the interpretation are known by both the teacher and the learner. The teacher generates explanations in such a way that the learner will, ideally, interpret them the way the teacher wants them to be interpreted. The same kind of convention-driven interpretation underlies the use of lesson sequences. The formatting information (i.e., the partition, the order, the marks) is generated by the teacher who understands the interpretation that that learner will place on that information. Two interpretive conventions are explored here:

1. A normal lesson introduces at most one subprocedure. Roughly put, a subprocedure is like one \texttt{COND} clause in Lisp: a test, which if true causes an implicit \texttt{PROGN} of function calls to be executed. A precise definition of "subprocedure" will be given after the knowledge representation language for procedures is described.

2. A normal lesson introduces material that will allow the learner to solve problems that it could not solve before. An optimization lesson shows the learner more efficient ways to solve the same class of problems that it could solve before the lesson. Furthermore, a normal lesson may not introduce optimized methods, wherein some of the procedure’s calculations are performed internally and do not appear in the action sequence.

These two conventions are not arbitrary. They directly address two of the worst combinatorial problems in induction. The first convention allows the learner to solve the disjunction problem, which involves deciding when and where to place disjunctions (The disjunction problem will be discussed at length in section 3.) For procedures, a disjunction is a branching (e.g. a \texttt{COND}) in the flow of control. Convention 1. above informs the learner that there will be at most one new disjunct per lesson. This cuts...
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**Figure 1-1:** A simple lesson sequence for subtraction
Figure 1-2: Procedures produced while traversing the lesson sequence of fig 1-1
down the possible places for disjunctions to a finite, small set and thereby significantly reducing the learner's search space of possible procedures. This makes the learner's job much easier. In fact, it makes it possible as opposed to impossible. The first convention is named one-disjunct-per-lesson.

The second convention fulfills a similar function with regards to a second combinatorial problem, the invisible objects problem. If two visible objects in a state can be related by arbitrarily long chains of calculations with arbitrarily many intermediate results, then induction is combinatorially infeasible. The intermediate results are invisible objects because they don't appear in the examples, if they could be seen, the combinatorics would be substantially reduced. The second convention provides this by mandating that normal lessons explicate such chains by showing all the intermediate results. Optimization lessons may come along later and show how to suppress the intermediate results and perform the calculations "in one's head." The convention is called show-work.

To evaluate the efficacy of this approach, or any approach to learning from material prepared by a teacher, one must evaluate burdens placed on both the teacher and the learner. One would expect there to be some work required of each, because learning-from-lesson-sequences lies halfway between induction, where the learner does most of the work, and learning-by-being-told, where the teacher does most of the work.

The teacher's job is to generate a lesson sequence for a given target procedure that satisfies one-disjunct-per-lesson and show-work. This task can be accomplished mechanically if the teacher writes down the target procedure in an appropriate procedural language e.g., Lisp. But writing procedures can be quite a bit of work. A more interesting possibility is that experienced teachers generate such lesson sequences naturally, without even realizing that their lesson sequences obey the two conventions. This is exactly what my research on naturally occurring lesson sequences in mathematics shows [44, 45]. Educators tend to generate well-formed lesson sequences even though they probably are not aware of the conventions. Apparently, they have an intuitive unconscious appreciation of the conventions. This allows them to generate appropriate lesson sequences without going through the work of explicating and formalizing the procedures taught by those curricula.
This raises an interesting possibility for applications where a computer system learns from a human user, such as programming-by-example systems (e.g., [8]) or learning apprentice systems (e.g., [30]). Such systems usually assume that there is no meaningful structure in the example sequence that is presented to the system. However, if the users view themselves as teaching the system, they may order their examples in certain ways. At the very least, they will present easy cases before hard cases. If we had a precise definition of the ordering criterion that users tend to employ, and if the learning system were designed to take advantage of these tacit constraints on the instructional material, then it could recover information that is latent in the sequential ordering. This latent information might allow it to converge faster and more reliably on the knowledge that the user is trying to teach it. One-disjunct-per-lesson and show-work are exactly such constraints, and they do make Sierra a more effective learner. Further research is needed to see how domain-general they are and to see whether there are other constraints like them.

It might seem strange that teachers should obey conventions like one-disjunct-per-lesson and show-work without being aware of them. Looked at a different way, it would seem strange if they didn’t. The teacher-learner situation is an extended communication act. We know that people naturally, unconsciously obey many conventions on natural language communication acts (see, e.g., [37]). It seems entirely likely that the teacher-learner discourse should also be formatted by conventions. In the hope that this analogy is approximately correct, the conventions that govern learning will be named felicity conditions, an early name for certain natural, unconsciously-followed language conventions [4].

The point of felicity conditions is to make the learner’s job easier without burdening the teacher too much. One-disjunct-per-lesson and show-work make Sierra’s job rather simple, although not trivial. The following is a quick sketch of how it works. The details will be presented later.

13 How Sierra works

Sierra’s algorithm is called learning by completing explanations. It begins by trying to parse an action sequence, using the procedure as if it were a grammar but a grammar with data flow, conditional tests etc. If parsing succeeds, the resulting parse tree is a trace (i.e., subroutine calling hierarchy). Looked at in a different way
the parse tree constitutes an "explanation" for the action sequence. For instance, a partial explanation for an individual action in a sequence can be read off the parse tree by walking upwards from the action, i.e., the action was done in order to satisfy the goals of the subprocedure that called it (which is the next node above it in the parse tree), and the caller was executed in order to satisfy its caller, and so on. A complete explanation for an action involves taking into account data flows and side-effects, so explicit links for these effects are included in Sierra's parse trees. Such links are analogous to the causal links that thread through the hierarchical structures of explanations of, e.g., kidnapping stories [10]. So explaining an action sequence is just parsing it. Such parsing is a form of plan recognition (e.g., [16, 20, 36]).

The conditional tests and data flows of a procedure are used to guide Sierra's parser, significantly narrowing its search for a parse tree. However, the parser may choose to relax such tests or ignore them entirely. This may allow it to find a parse tree when it could not do so otherwise. If so, then a simple form of learning can be performed. The relaxations made by the parser are edited into the procedure. For instance, if certain predicates in a conditional test must be ignored by the parser, then Sierra removes them from the conditional. This generalizes the condition test. Now if the parser redoes the narrow search, the one that obeys the constraints imposed by conditional tests, etc., it will find the parse tree. Sierra has generalized the procedure allowing it to explain examples that it could not explain before. This learning technique is similar to one form of explanation-based learning [10, 11, 14, 32, 39]. A more interesting kind of learning occurs when it is impossible to complete a parse, no matter how much the procedure is generalized. In this case, the learner's procedure is fundamentally incomplete. One or more new subprocedures must be invented. Learning by completing explanations is one approach to accomplishing the learning required in this situation.

Sierra uses a straightforward technique. A similar approach was employed in three independent investigations [16, 18, 33]. The first step is to parse the action sequence bottom-up as far as possible and top-down as far as possible. The candidate solutions to the learning task consist of any new subprocedure (or set of new subprocedures) that links the top-down parse to the bottom-up parse in such a way that a complete parse tree is yielded. Even for short action sequences, there can be
millions of candidates. The challenge is to cope with this large space of candidates. The solution used by the three independent investigators [ibid.] is to place such strong constraints on the parsing that only one (or a few) of the possible candidates are generated. My solution is to (1) use unconstrained parsing, (2) assume that only one new subprocedure will be acquired, and (3) use a factored data structure (similar to LUNAR's well-formed substring table [51] and GSP's chart [21]) to efficiently represent the space of possible candidates. The selection of a candidate from this space is accomplished by a collection of simple filters.

This technique for learning by completing explanations makes it simple to perform induction across several action sequences. Each action sequence yields a space of candidate solutions represented as a GSP-style chart. Induction amounts to intersection of these spaces.

In principle, this technique can be used in any domain which learns hierarchical knowledge structures from sequential examples. Thus, it should extend to learning grammars from strings, learning story understanding schemata from stories, and learning device models from the operation of machines.

Many of the basic intuitions behind Sierra have been presented. The remainder of this article presents the details of how Sierra accomplishes its task. The first section presents the knowledge representation language used for procedures. The next section reviews the induction problem, and isolates disjunction as a key difficulty. One-disjunct-per-lesson is proposed as a central simplifying constraint for inducing procedures. It in turn leads to the definition of "subprocedure," in section 4, in terms of three parts: its skeleton, its patterns and its functions. This reduces the problem of inducing a subprocedure to three subproblems, one for each part. The subsequent sections discuss the Sierra algorithms, for, respectively, skeleton induction, pattern induction and function induction. The final sections discuss the generality of Sierra and speculate on the origins and applications of felicity conditions.

2. The representation of procedures

It is convenient to use a mixture of nomenclature from production systems and And-
Or graphs (AOGs). The latter equivalence class of representations includes the one used by Sierra. Although any of the formalisms could be used to describe Sierra's representation the mixture of production systems and AOG is used here.) The production systems nomenclature is good for showing details, the AOG view is good for showing the overall structure.

Figure 2-1a sketches the AOG view of a subtraction procedure learned by Sierra. The nodes are called goals, and links are called rules. Rules are directed and are always drawn running downward. The nodes just beneath a goal are called its subgoals. Currently, there are two types of goals: AND and OR. To execute an AND goal, all the subgoals are executed. To execute an OR goal, just one of the subgoals is executed. AND goals are drawn with boxes around their labels. Drawings of AOGs abbreviate goals whenever they appear more than once. For instance, OVRWRT is called from several places in the AOG of figure 2-1a, but its subgoals are drawn only for one of these occurrences. Although abbreviation makes this AOG look like a tree, it is really a cyclic directed graph due to the recursive calls of MULTI and REGROUP.

AOG drawings do not indicate several kinds of information. This information is readily visible in the production system view. Figure 2-1b shows the definitions for the non-primitive goals in the AOG of figure 2-1a. Goals have arguments. For instance, SUBICOL has three arguments, T, B and A. Arguments have the substitution semantics of lambda calculus. That is, the AOG language is applicative. There are no assignment statements. The only side-effect operators are those that change the external state, i.e., writing a digit in an answer. The applicative property has important consequences that will be discussed later.

A goal's rules (i.e., the rules leading from it to its subgoals) are listed in its definition. SUBICOL has three rules. Each rule has a pattern and an action. Patterns are large, so in figure 2-1b, most patterns have been replaced by English

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Footnotes:

1. It is well known that context-free grammars, push-down automata and basic transition nets are equivalent in the same way. Attribute grammars [22], affix grammars [23] and ATNs are equivalent, provided that side-effect operators (e.g., SETO) are not used in the ATNs.

2. In order to accommodate certain empirical data, a third goal type, FOREACH will be added as a representation for iteration across a sequence of objects. This goal type appears in figure 2-1b as a loop in the ATN's graph. Currently, Sierra represents iteration with two invocations.
Figure 2-1: A subtraction procedure shown as (a) an ACG, and (b) a production system.
glosses. A pattern is a conjunction of literals (i.e., predicates or negated predicates). Predicate arguments may be either arguments from the enclosing goal or pattern variables. As an example, the pattern

\[(\text{Column C})\&(\text{Top C T})\&(\text{Digit T})\&(\text{Bottom C B})\&(\text{Digit B})\&(\text{LessThan T B})\]

matches columns that require borrowing. The empty pattern always matches.

A rule's action is a form, in the Lisp sense, which calls the rule's subgoal. The action may pass arguments to the subgoal, often by evaluating functions. For instance, SUBICOL's third rule has \((\text{Write A} \ (\text{Sub} \ (\text{Read T}) \ (\text{Read B})))\) as its action. This action writes the difference of the top and bottom digits of a column in the column's answer.

An OR goal's rules are tested in left-to-right order. The first rule whose pattern matches is executed. The learner adds new rules at the left. Hence, the left-to-right ordering convention corresponds to a common conflict resolution strategy in production systems called "recency in long term memory" [26]. Because the patterns of OR rules test whether to execute a rule, they are called test patterns. Although AND rule patterns have the same syntax as OR rule patterns, they are not used to control which rules are executed. The order of execution of AND rules is fixed: the rules are executed in left-to-right order. AND rule patterns are used to retrieve information in the current problem state so that the information can be passed to the rule's subgoal. AND rule patterns are called fetch patterns.

Any learning model that describes how knowledge is constructed from smaller units is open to questioning about its set of primitives: what are the units that are assumed to be present when learning begins? For completeness, table 2-1 lists the kinds of primitives used by Sierra, and the particular ones employed to learn the procedures discussed in this article. In addition to these primitives, the initial knowledge state may contain non-primitive procedures as well. For instance, the initial procedure from which the procedure of figure 2-1 was learned contained the non-primitive goal \text{CURWRT}, which crosses out a symbol and writes another symbol over it. The multiplication procedure's initial knowledge state included an addition procedure.

The procedural representation language has been presented. The remainder of this section is a "walk through" of the procedure of figure 2-1 which some reader may
1. Primitive actions cause a change in the current problem state. The only primitive actions used in mathematics are ones that write a given alphanumeric symbol at a given position (Write), or ones that write special kinds of symbols (CrossOut puts a slash over a symbol; Bar writes a bar under a group of symbols).

2. Fact functions return a number without changing the problem state. The following fact functions were employed: Add, Sub, Add1, Sub1, Mult, Quotient, Remainder. One (which always returns 1), Zero (which returns 0), and Concat (which concatenates two numbers, e.g., (Concat 1 4) returns 14).

3. Fact predicates return true or false without changing the problem state. The fact predicates used were LessThan?, Equal?, and Divisible?.

4. The primitive function Read returns the symbol written at a given place.

Table 2-1: The four kinds of primitives used by Sierra

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find helpful as a way of cementing their understanding of the representation.

The root goal START, and its subgoal SUB, initialize column traversal to start with the units column. IMUL chooses between three subgoals. MULTI is for multiple column problems. REGROUP is for "regrouping" exercises that don't involve any subtraction at all. Regrouping is the part of borrowing where one digit is reduced by one and an adjacent digit is increased by ten. This subgoal is left over from learning regrouping separately from multi-column subtraction. (The lesson sequence of figure 1-1 lacks a regrouping lesson, but most textbooks include one. This procedure was learned from Heath's subtraction lesson sequence [13], which has a separate regrouping lesson.) Normally, SUB never calls REGROUP. The third goal, Write, is for single column subtraction problems. The "main loop" of multi-column traversal is expressed by MULTI as a tail recursion. MULTI calls itself via its subgoal SUB/REST. SUB/COL processes a column. It chooses between three methods for doing so if the bottom of the column is blank, it copies the top of the column into the answer via the subgoal SHOW2. If the top digit of the column is less than the bottom, it calls BORROW. Otherwise, it writes the difference of the two digits in the answer. BORROW has two subgoals. 1/BORROW calls REGROUP, and 2/BORROW just takes the difference in the column and writes it in the answer. REGROUP is a conjunction of borrowing into the
column that originates the borrow (BORROW/INTO) and borrowing from the adjacent column (BORROW/FROM). In this procedure, BORROW/FROM occurs before BORROW/INTO. It would be equally correct to reverse their order, but that is not the way that Heath teaches them. Borrowing into a digit is just adding ten to it. Borrowing from the next column is also easy when its top digit is non-zero: the digit is decremented. If the digit is zero, it calls BFZ. BFZ regroups, which causes the zero to be changed to ten, then it decrements the ten to nine.

3. Disjunction: an inherent problem for induction

For many kinds of induction tasks, there are proofs that the task has no algorithmic solution. In such proofs, the induction problem is defined by specifying a class U of all possible generalizations that the learner can output and a class T of all possible trainings that the learner can receive. (The "U" stands for the learner's universe of generalizations.) There are a variety of theorems for various U and T. For instance, one such theorem is: If U is the class of recursive functions and T is the set of all possible training sequences that contain all possible positive and negative examples, then there is no Turing machine that can learn any given generalizations from U [35, Proposition 5]. Such negative results guarantee that there is no straightforward solution to induction.

The standard attack is to incorporate biases into the inducer [28]. There are two kinds. An absolute bias is a unary predicate on generalizations that says whether or not the generalization should ever be output by the inducer. A relative bias is a binary predicate on generalizations that says which of the two generalizations is preferred for output in case both generalizations are consistent with the given training. Often, absolute biases are implemented by representing generalizations in a limited representation language. If the generalization can not be expressed in the language (say, because the language lacks the appropriate primitives), then it will never be output by the inducer. A relative bias, on the other hand, is usually defined by comparing two formal expressions that represent the generalizations. Simplicity metrics are a common relative bias. AI inducers often implement relative biases implicitly by the order in which they search. Because they stop when they get to the first generalization that is consistent with the training, the search strategies act as relative biases. In short, biases correspond to two obvious kinds of constrains: unary and binary predicates on generalizations.
A non-standard approach is to employ a third kind of constraint, which could be called a manner constraint. A manner constraint relates a generalization to the manner in which the training is presented. A manner constraint is a binary predicate: one argument is a generalization and the other is the form (syntax) of the training. Both one-disjunct-per-lesson and show-work are manner constraints\(^5\).

Manner constraints are a known loophole to most formal learnability results. For instance, a major result \([17]\) is that it is impossible to learn when (1) \(T\) employs only positive examples and (2) \(U\) contains a generalization for every finite set of examples and a generalization for at least one infinite set of examples, unless (3) the examples sequences in \(T\) are ordered by some primitive recursive function. That is, conveying information with the order of the example sequences allows a learner to succeed where it could not otherwise. The manner of example presentation is a factor that hasn't been studied much, but is potentially quite important.

There are dozens of theorems on the learnability of certain \(U\) given certain \(T\). In order to obtain such results, it is necessary to be quite specific about what the \(U\) and \(T\) are. Rather than review all these specific results, it seems more profitable for this article to sacrifice formality in order to uncover the key components of \(Us\) and \(Ts\) that cause induction to be impossible. In spirit, this strategy is like Newell's Knowledge Level strategy of stepping back from the details of countless AI problems and representations in order to analyze them from a single perspective, which he suggests should be the perspective of first order logic \([34]\).

In that spirit, I suggest that one of the inherent problems of induction is disjunction. Whenever \(U\) allows a generalization to be built from the disjunction of any two other members of \(U\), then induction is infeasible.

More specifically, suppose that \(g\) and \(g'\) are two generalizations from \(U\). Even without knowing how they are represented, we can define their disjunction. Each generalization has an extension, i.e., the set of all possible examples (instances)

\(^5\) The terminology could use a little clarification here. Felicity conditions are defined as interpretive constraints on skill acquisition that people obey without being aware of them. Presumably, one could explicitly tell students a manner constraint, in which case it wouldn't qualify as a felicity condition. Presumably, there could be felicity conditions that aren't manner constraints, but say, relative biases.
consistent with the generalization. Let \( x \) and \( x' \) be the extensions of \( g \) and \( g' \), respectively. The disjunction of \( g \) and \( g' \) is any generalization whose extension is the union of \( x \) and \( x' \). This is the definition of "disjunction" that will be used in this article.

Disjunctions often correspond to syntactical constructions in representational language. In AOG representations, disjunctions correspond to OR goals. In context-free grammars, a disjunction is present when two or more rules reduce the same non-terminal category. In production systems, a production is potentially disjunctively related to all the other productions; it requires careful analysis to uncover the actual disjunctive relationships.

Induction's trouble occurs when the class of all possible generalizations admits free disjunction. That is, the disjunction of \( g \) with \( g' \) is in the class whenever \( g \) and \( g' \) are. When this is the case, induction acquires some strange properties that make it seem quite unlike anything that one would want to call "learning."

Free use of disjunction allows the learner to generate absurdly specific generalizations. One such absurdity is the trivially specific generalization: a disjunction whose disjuncts are exactly the positive examples that the learner has received. Thus, if the learner received positive examples \( a \), \( b \), and \( c \), then the disjunction \((\text{OR} \ a \ b \ c)\) is the trivially specific generalization. The trivially specific generalization is not really a generalization at all. Its extension is just \( \{a \ b \ c\} \). The learner didn't really learn it; just remembered. This problem could be solved with an arbitrary prohibition against trivially specific generalizations, if nothing else. But there is another problem that is much worse.

When disjunctions are unconstrained, the learner has to be given the complete extension of the generalization being taught before it can reliably discriminate that generalization from the others. To see this, first assume that for each example there is a generalization in \( U \) whose extension is that example and only that example. This is equivalent to assuming that the examples can be represented in the representation language used for generalizations. For instance, a grammar consisting only of the rule \( S \rightarrow w \) is such a generalization where \( S \) is the root category and \( w \) is a string of terminals. This grammar's extension is the singleton set \( \{w\} \). Using such singleton
generalizations and disjunction, any finite set of examples can be described by some generalization. To get the generalization for \( \langle w_1, w_2 \rangle \), one finds the generalization for \( \langle w_1 \rangle \) and for \( \langle w_2 \rangle \), then forms their disjunction. Since all finite sets of examples correspond to generalizations, the learner can't tell which generalization is correct until it is told exactly what the target generalization's extension is. This means it must be shown all possible examples and be told which are positive examples and which are negative examples. Such conditions, where a "learner" is shown a complete extension of a generalization and asked to identify the generalization, hardly qualify as learning.

Suppose learning is viewed as the following rather naive search for generalizations. This will provide another perspective on the trouble that disjunction causes. Suppose that the learner receives an initial example and generates a single generalization. Suppose further that the next example is a positive one that the learner's current generalization is not consistent with. The learner has two choices: (1) to modify the current generalization enough so that it becomes consistent with the new example, or (2) to create a generalization specifically for the new example, then disjoin that generalization with the current one. Roughly speaking, these two choices are available at every step, so after \( N \) examples, there will be roughly \( 2^N \) possible generalizations consistent with the examples. The point is simply that the set of consistent generalizations grows as the learner is given more positive examples. It doesn't shrink, as one would intuitively expect of learning from examples. On the other hand, if disjunction is barred from \( U \), then there is only one choice at each step, and the set of generalizations does not grow unboundedly. So this learner's failure to learn can be blamed squarely on disjunction.

One-disjunct-per-lesson would constrain this naive learner's search while allowing generalizations to contain disjunctions. Most of the time, the learner would have a single choice. However, on the first example of each lesson, it would have two choices. It can either disjoin or not. If it chooses to disjoin then it may not disjoin again until the next lesson. If it chooses not to disjoin then the twofold choice is again available on the next example. One-disjunct-per-lesson is indeed a...
straightforward solution to the disjunction problem

One-disjunct-per-lesson solves the disjunction problem by modifying the relationship between $T$ and $U$. Absolute and relatives biases solve the disjunction problem by modifying $U$. A relative bias partially orders the elements of $U$ by preferring, e.g., generalizations with fewer disjunctions. An absolute bias removes from $U$ all generalizations that contain disjunctions (or contain more than 13 disjunctions, etc.). In general, biases modify $U$, and manner constraints modify the relationship between $T$ and $U$.

Biases are appropriate for a learning task when the learner can make strong a priori assumptions about $U$. Manner constraints are appropriate when the learner can't make assumptions about what's going to be learned, but it can make assumptions about how it's going to be taught. For Sierra's task domain, it is inappropriate to use biases to solve disjunct problem. There is no reason for the learner to believe that a procedure should have no conditionals (or less than 13 conditionals). So an absolute bias against disjunction is inappropriate. There is no reason for the learner to believe that a procedure with the fewer conditionals is better, so a relative bias is also inappropriate. However, given the felicity conditions hypothesis, there is reason to believe that $T$ and $U$ are related. So a manner constraint such as one-disjunct-per-lesson is appropriate.

The preceding comments were meant to motivate the pragmatic utility of manner constraints in general, and one-disjunct-per-lesson in particular, by considering how one-disjunct-per-lesson helps solve the disjunction problem, one of the inherent problems of induction. Later, a similar motivation will be presented for the show-work manner constraint, based on another inherent problem of induction the invisible objects problem. Manner constraints are a new technique for solving inherent problems in induction. Previous AI learners have employed either absolute biases or relative biases. Although any addition to AI's toolkit of techniques is welcome, manner constraints seem particularly welcome, for they are remarkably general as the preceding discussion argued, and they are quite effective in reducing the complexity of programs for learning procedures, as the remainder of this article shows.
4. Subprocedures

In order to make one-disjunct-per-lesson easy to implement, the AOG representation permits disjunctions in only one place: the OR goals. Thus, a disjunct is an OR goal's rule plus, roughly speaking, whatever that rule calls. Such fragments of AOGs are called subprocedures. A subprocedure consists of several components:

1. A new OR rule that is placed beneath an existing OR goal. The existing OR goal is called the parent.

2. A new AND goal, which is called by the new OR rule.

3. The new AND has one or more rules. Each rule calls a new OR goal that has just one rule. These ORs are merely a convenience. They provide a place for later subprocedures to attach.

4. Each such OR has a single rule that calls some existing AND goal. These existing AND goals are called the kids.

Figure 4-1 illustrates these components of a subprocedure by showing an AOG before and after a subprocedure has been added. This subprocedure was acquired from a lesson that teaches how to borrow across zeros. The pre-lesson AOG (figure 4-1a) can borrow only from non-zero digits; the post-lesson AOG (figure 4-1b, which is the same as figure 2-1) can borrow across zeros. BORROWFROM is the subprocedure's parent. The new OR rule connects BORROWFROM to BFZ. The new AND is BFZ. The kids are REGROUP and QVART. BFZ and 2BFZ are the new OR goals that are added as places to attach future subprocedures.

One-disjunct-per-lesson takes us as a long way towards solving the whole procedure induction problem. Inducing a procedure is reduced to a series of subprocedure induction problems: one per lesson. A subprocedure induction problem reduces to three subproblems:

- **Skeletons**: Skeleton induction determines the parent and the kids of the new subprocedure. This establishes the topology of the new subprocedure (i.e., the connectivity of the goals and rules). Because it doesn't determine the conditions and action arguments of the new rules, skeleton induction is like inducing only the bones and not the flesh of the subprocedure.

- **Patterns**: One-disjunct-per-lesson entails that a rule's conditions have no disjunctions. This means that they can be induced by standard disjunction-free pattern induction techniques.
Figure 4-1: An ADG before (A) and after (B) a lesson. The new subprocedure is circled.
Functions: One-disjunct-per-lesson entails that disjunctions are not permitted in the nests of functions that express the data flow in AOGs. This makes inducing those function nests easier. However, section 7 shows that function nest induction is still infeasible unless more constraints are added. The show-work felicity condition is the key to Sierra’s solution.

These three induction tasks will be discussed serially in the following sections. In Sierra, skeleton induction happens first using one pass over the lesson’s examples. Pattern and function induction occur together on a second pass over the examples.

5. Skeleton induction

To see what skeleton induction involves, a computer science fixture is needed: the trace of a procedure’s execution. Figure 5-1 shows the trace tree for a correct subtraction procedure (the procedure is shown in figures 2-1 and 4-1b) solving a BFZ (i.e., borrow from zero) problem. Each call is shown as a tree node, with its arguments abbreviated. A trace tree is just a parse tree for the action sequence, using the procedure as the grammar.

Roughly speaking, a skeleton is a hole in a trace tree. If the procedure is missing the BFZ goal, then the trace tree would have a hole in the middle of it, as in figure 5-2. The gap is right where the BFZ node would be. From the figure, one can see that a skeleton can be characterized by the link coming into it from above and the links leaving it from below. Thus, a skeleton is uniquely specified by the parent and the kids.

Almost all action sequences, including the example of figure 5-2, admit more than one skeleton. Most of the ambiguity is due to the fact that one can almost always make a skeleton bigger. The kids can be lower in the tree (e.g., figure 5-3), the parent can be higher (e.g., figure 5-4). Any node that would complete an otherwise incomplete trace tree is a legitimate skeleton.

Sierra uses two context-free grammar parsing algorithms to enumerate the skeletons: A top-down recursive descent parser is used to find all possible parents (it is actually just a non-deterministic version of the regular AOG interpreter). A bottom-up, breadth-first parser is used to find all possible kids. For each parent, all possible tuples of kids are collected where a kid tuple is a sequence of adjacent kids that together span the same part of the action sequence as the parent. This generates the set of
Figure 5-1: Trace tree for solution of a BFZ problem
Figure 5-2: The skeleton is right where the SFC node and its daughters were
Figure 5-3: The kids of the skeleton can be lower
Figure 5-4: The parent of the skeleton can be higher
all possible skeletons. In general, there can be thousands of possible skeletons (e.g.,
(50 possible parents) x (3 kids per kid tuple on average) x (10 possible kids for each
tuple position on average) = 4500 possible skeletons.) Sierra represents this set
implicitly in order to save space.

Notice that there would be many more possibilities if it weren't for one-disjunct-per-
lesson. For instance, if two new subprocedures were allowed, then Sierra would have
to collect all possible pairs of parents, each with their possible kids, etc.

Parsing Sierra's AOGs is equivalent to parsing attribute grammars, for which there are
many algorithms [46]. Such parsing is quite simple because the language is
additive (i.e., no side effects, no assignment statements). Once a goal's arguments
are bound by its caller, those values never change. The arguments act as a
subcategorization of the goal. This makes parsing nearly as simple as context-free
grammar parsing. If side-effects were allowed, parsing would still be possible, but it
would be combinatorially costly, because the left (preceding) context of a goal would
have to be included in the goal's subcategorization.

Bottom-up parsing requires inverse execution of AOG code. Sierra must be able to
figure out the arguments of a caller from the arguments of its callees. This requires
matching fetch patterns "backwards" and inverse evaluation of functions. The first is
easy. Matching patterns backwards is the same as matching them forwards.
Backwards evaluation of arithmetic functions, such as (Add x y), is accomplished by
hand-coded inverse functions that produce sets of tuples that represent possible input
values. This technique would collapse if it were not feasible to assume that examples
only use small numbers. If numbers could be arbitrarily large, then other techniques
(e.g., symbolic execution, followed by solution of a system of polynomial equations)
would have to be employed.

The techniques just mentioned generate one set of candidate skeletons per action
sequence. One-disjunct-per-lesson entails that skeleton induction can be performed
simply by intersecting these sets. Because a lesson may introduce just one
subprocedure, all the skeletons' parents must be the same. Because the new
subprocedure is disjunction-free, each skeleton's list of kids must be equal to each
other skeleton's list of kids. In particular, two kid tuples A B C and A B D cannot be
merged by using disjunction on the middle kid to form something like \(A \lor B \lor C\).

Skeleton intersection is powerful enough that it is usually possible to devise a lesson that yields a unique skeleton when its examples' skeletons are intersected. However, in some cases, this is not possible. In fact, the BFZ skeleton that is our running illustration cannot be uniquely specified by examples\(^3\) in such cases, there is no choice but to guess. This means that relative biases for skeletons are required.

Sierra has been used to test various relative biases in order to find the ones that explain the skeleton choices the people make. Sierra first generates all consistent skeletons using parsing and skeleton intersection. In manual mode, it displays them in a menu and allows the user to choose one. This is useful for exploration. In automatic mode, Sierra partially orders the skeletons using the biases that are being tested. Usually there is just one skeleton that is maximal in the partial order. If so, Sierra just takes it and goes on. If there is more than one, Sierra chooses the first one, then stores the learner's state so that it can come back later and take the other choices. There is nothing new about this architecture, but it is remarkable how easy it makes it to search the space of hypotheses. The current best hypothesis is that people choose the smallest, most deeply embedded skeleton. See [44], chapters 18 and 19, for a complete discussion.

5.1 Prior solutions to the skeleton induction problem

Skeleton induction determines the procedure's goal-subgoal calling hierarchy. Inducing such hierarchies has proved to be a tricky problem in machine learning. Neves [33] used hierarchical examples to get his procedure learner to build hierarchy. However, subtraction teachers rarely use such examples. Badre [5] recovered hierarchy by assuming examples are accompanied by a written commentary. Each instance of the same goal is assumed to be accompanied by the same verb (e.g., "borrow"). This is a somewhat better approximation to the kind of input that students

---

\(^3\)The data relies on the distribution of branches in the ACG that is induced in the earlier. The ACG is shown in figure 5.4. If the examples are designed so that either \(A\) or \(B\) or \(C\) has \(e\) as parent, then \(e\) in \(A\cdots\lor B\cdots\lor C\cdots\) is always possible. If a new name \(v\) is given, and \(v\) is always the first input parameter in the example, it is possible to substitute \(v\) for \(e\) in the middle. This is an example of how the intersection can lead to the next generation of examples.
actually receive, but again it rests on delicate and often violated assumptions. Anzai
and Simon [3] used production compounding (chunking) to build hierarchy. However
to account for which of many hierarchies would be learned, Anzai used domain-specific
features, such as the pyramids characteristic of subgoal states in the Tower of Hanoi
puzzle. Sierra's technique, learning by completing explanations, is less domain specific
than Anzai and Simon's technique, and requires less information from the teacher than
Neves and Badre's technique.

6. Pattern induction

The rules in the new subprocedure must be given appropriate conditions. The new
OR rules require test patterns, and the new AND rules require fetch patterns. Because
patterns have no disjunctions nor other representational devices that trouble induction
patterns can be induced using standard techniques. This section presents the ones
that Sierra uses. Although pattern induction is not particularly interesting from a
theoretical standpoint, it has turned out to be the bottleneck in Sierra's performance.

Parsing an action sequence with the new subprocedure installed will pair the new OR
rule with a problem state where its test pattern would have to be true in order for the
parse to go through. Such states are positive instances for the test pattern. Similarly, parsing can uncover states where the test pattern would have to be false in
order for the parse to go through. These states are the negative instances. The test
pattern induction problem is to find a test pattern that matches all the positive
instances and none of the negative instances. Parsing also collects positive instances
for the fetch patterns, together with the values that the fetch patterns should return.
The fetch pattern induction problem is to find a pattern that matches all the positive
instances and returns the appropriate values each time.

Both induction problems are solved using version spaces [29]. A version space is
represented by a pair <S,G>, where S is the set of maximally specific patterns
consistent with the instances received so far and G is the set of maximally general
patterns consistent with the instances. To use version spaces, several application-
specific functions must be defined. The most important two are Update-S and
Update-G. Given a version space and a positive instance, Update-S generalizes the
patterns in S so that they match the instance and remain maximally specific. To
implement Update-S, Sierra uses an algorithm that finds the largest common subgraph
of two labelled directed graphs. Given a version space and a negative instance, Update-G augments the patterns in \( G \) so that they do not match the negative instance. Update-G is implemented with an algorithm for generating minimal covers of a set.

Sierra is designed so that pattern induction never finishes. A rule keeps the version space of its pattern so that induction can continue whenever the pattern requires more refinement. It is almost always the case that introducing a new subprocedure will cause the patterns in older subprocedures to be modified. Those older patterns will be "seeing" problem states that they have never "seen" before, namely, the ones that trigger the new subprocedure and the ones that the new subprocedure produces. In order to continue to function properly, the older patterns must be generalized to match these new states. The generalization of older patterns to match new situations is a simple form of explanation-based learning, as the term is used by some authors (c.f. [11]).

Since pattern induction occurs so often, it needs to be fairly efficient. However, the Update-S computation is NP-hard, and the Update-G routine calls it as a subroutine. More specifically, if pattern \( P \) has \( n \) variables, and it is matched against a problem state with \( m \) objects, then Update-S takes \( O(m^n) \). These combinatorics reflect the usual AI matching problem: Each variable in \( P \) can be paired with any object in the problem state.

One way to deal with the complexity of pattern induction is to use a small \( n \) and/or \( m \). For instance, blocks-world inducers (e.g., [12, 48]) typically have an \( n \) and \( m \) of less than 5. Using a small \( n \) and \( m \) is impossible in Sierra's case. Its task domain

\[^{9}\text{Sierra's patterns are conjunctions of literals, where a literal is a predicate or a negated predicate. Such patterns are similar to labelled directed graphs, where the variables are nodes and the relations are labelled arcs. If all the predicates in a pattern are binary, the correspondence is exact.}\]

\[^{10}\text{Given a } g \text{ from } G \text{ that matches the negative instance } N, \text{ and an } s \text{ from } S \text{ that doesn't match } N \text{ Update-G needs to find relations in } s \text{ to add to } g \text{ such that the revised } g \text{ doesn't match } N. \text{ Update-G first generates all possible mappings of the variables of } s \text{ into the variables of } N. \text{ Each such mapping becomes an element of the set of maps, } M, \text{ to be covered. Each relation in } s \text{ is paired with the subset of } M \text{ that contains the maps under which the relation is not a member of } N. \text{ The relation is said to "cover" that subset of } M. \text{ A key fact is that a conjunction of relations covers at least the union of their individual covers of } M. \text{ The main goal is to find the smallest conjunction of relations that covers all of } M. \text{ Such a conjunction is not a member of } N \text{ under any map, so adding it to } g \text{ will yield a pattern that doesn't match } N, \text{ which is what we want. So the algorithm for finding minimal covers of } M \text{ yields candidates for the revised } g. \text{ Some of these candidates may generalize others, so Update-G has one more task, which is to filter out the candidates that are not maximally general.}\]
requires problem states with 10 to 50 objects. The larger patterns in the version
spaces have about the same number of variables as objects in the problem states
they were induced from. The combinatorics for straightforward pattern inducer can go
as high as 50^50.

A second solution is to impose constraints on which variable-object mappings will be
considered. Sierra uses two constraints. First, pattern variables have an implicit
inequality relationship between them. That is, distinct variables must match distinct
objects. This lowers the combinatorics to the binomial coefficient function \( \binom{n}{m} \) \( (n^m)/(m^m) \). Second, the patterns and problem states are split into two components: a
part-whole tree and the rest. In a problem state, the part-whole tree is simply the
usual parse tree for the mathematical notation. For instance, a subtraction problem's
parts are its columns and a column's parts are its digits. In patterns, there are
variables for each of the components (i.e., a variable for the problem for each column
and for each digit), and their part-whole relationships are kept separately from the
main pattern. Pattern induction (and pattern matching, too) considers only variable-
object mappings that do not violate the tree topologies. This cuts the complexity
down to \( O(B^{100s}) \), where \( B \) is the branching factor of the part-whole trees, typically
about three. When this constraint is turned off in Sierra, an Update-S that normally
takes 10 seconds takes hours. This constraint or something like it is a practical
necessity. Even with it, most of Sierra's time is spent running the pattern induction
algorithms. In a typical run, about 70% of the time is spent doing pattern
inductions.

---

1 Pattern variable \( x \) is paired with object \( o \) in a state, and \( \text{Part of } x \text{ is } s \) in the pattern (which
means that \( x \) is a part \( s \)). Then \( x \) can only be paired with objects that are parts \( s \) in the state, and
there are \( m \) such objects, then \( x \) is left unpaired.

2 Sierra uses a further trick. It avoids as much pattern matching and pattern induction as possible by
using action sequences in two passes. The first pass ignores the subgoals generated by goal
substitutes and therefore does no pattern matching. The second pass takes the same trees produced by
the first pass and instantiates them by doing all the requisite pattern matching and then evaluating
in case some parse trees are pruned. In the first pass, the pattern matching that they would have
been carried out.

7. Function induction

Some of the rules in the new subprocedure may require function nests to be
induced for their actions. Functions are used to represent number facts such as
(Sub 17 8) = 9. A typical function nest is (Sub (Add 7 10) 8). This section discusses
how function nests can be learned. This learning task is called function induction.
Function induction involves discovering which function or nest of functions will yield the
numbers shown in the examples. For instance, suppose the learner already knows
how to do single column subtraction problems and it is taking a lesson on two-column
subtraction. After seeing examples a and b.

\[
\begin{array}{c}
\text{a.} & 72 \\
& - 41 \\
& \hline
& 31
\end{array} \\
\begin{array}{c}
\text{b.} & 74 \\
& - 21 \\
& \hline
& 53
\end{array}
\]

there are many function nests that explain where the tens column answer comes from.
Here are three candidates:

1. \( A_{10} = T_1 T_2 B_{10} \)
2. \( A_{10} = T_1 + B_1 \)
3. \( A_{10} = (T_{10} - T_1)(B_{10} + B_1) - A_1 \)

where the subscripts indicate the column, and \( T, B, \) and \( A \) stand for the top, bottom
and answer. The first generalization is the correct one. The second generalization is
that the ten's answer is the sum of the units columns' digits. This second
generalization, although consistent with examples a and b, is inconsistent with c:

\[
\begin{array}{c}
\text{c.} & 36 \\
& -12 \\
& \hline
& 24
\end{array}
\]

Many such accidental generalizations can be eliminated by giving lots of examples.
However, generalization 3 can never be eliminated that way. It is true of any
subtraction problem. This may seem like a peculiarity of this case, but it isn't. There
are infinitely many polynomials consistent with any finite set of input-output number
tuples. In particular, there are infinitely many \( n \)-degree polynomials consistent with any
\( n \) points.

The underlying problem has nothing to do with the fact that functions and
polynomials are the representation language of generalization. Any functional
expression can be easily converted to a relational one. For instance, generalization 3 above could be expressed as

\[
\text{AND} \quad (\text{PLUS} \ X \ T.. \ T.) \\
(\text{PLUS} \ Y \ B.. \ B.) \\
(\text{PLUS} \ Z \ A.. \ A.) \\
(\text{MINUS} \ Z \ X \ Y) \\
(\text{INVISIBLE} \ X) \\
(\text{INVISIBLE} \ Y) \\
(\text{INVISIBLE} \ Z))
\]

where \((\text{PLUS} \ u \ v \ w)\) means \(u = v + w\). The special relation \text{INVISIBLE} is needed because \(X, Y\), and \(Z\) do not match any of the visible objects in the examples. Under normal confirmation conventions for relational descriptions [19], the variables match only visible objects, so variables that designate invisible objects must be specially marked, and that is what \text{INVISIBLE} does.

Looked at this way, the underlying induction problem is clear: if the representation allows invisible object designators, then there will always be far too many generalizations consistent with any finite set of examples. Some constraint must be placed on the use of invisible objects in examples. This induction problem will be called the invisible objects problem.

7.1 Prior solutions to the invisible objects problem

AI's most common solution to the invisible objects problem is to ban invisible object designators from the representation. For instance, Winston's blocks world representation language [48] could have employed an elegant expression of the arch concept if it allowed invisible objects:

\[
\text{AND} \quad (\text{ISA} \ \text{LINTEL} \ '\text{PRISM}) \\
(\text{ISA} \ \text{LEG1} \ '\text{BRICK}) \\
(\text{ISA} \ \text{LEG2} \ '\text{BRICK}) \\
(\text{ISA} \ \text{GAP} \ '\text{BRICK}) \\
(\text{INVISIBLE} \ \text{GAP}) \\
(\text{SUPPORTS} \ \text{LEG1} \ \text{LINTEL}) \\
(\text{SUPPORTS} \ \text{GAP} \ \text{LINTEL}) \\
(\text{SUPPORTS} \ \text{LEG2} \ \text{LINTEL}) \\
(\text{ABUTTS} \ \text{LEG1} \ \text{GAP}) \\
(\text{ABUTTS} \ \text{GAP} \ \text{LEG2}))
\]

This says that the lintel rests on three abutting bricks, and the middle one is invisible. Using differently shaped invisible blocks for the gap is a simple way to describe
pyramidal, trapezoidal and circular arches as well as the rectangular arch above. However, the invisible objects problem makes it impossible to induce such descriptions. Once invisible blocks are allowed, they could be anywhere. The inducer would have no way of knowing whether there was just one invisible block, the gap, or dozens lying around all jumbled up. Winston avoids the problem by omitting invisible object designators from the representation, and employing the relationship \texttt{NOT TOUCHING LEG1: LEG2:} to express the gap between the arch's legs.

Banning invisible object designators is one way to solve the invisible objects problem. But it won't work in the mathematical domain. Invisible object designators are needed for representing procedures such as multi-addend addition. In the problem $1 - 3 + 5 = 9$, the intermediate result, either a 4, 8 or 6, is invisible.

As mentioned earlier, absolute and relative biases are the two customary ways to succeed at induction. An absolute bias, banning invisible objects, was just discussed. For a relative bias, the obvious candidate is to prefer generalizations with the fewest invisible objects. This is roughly what \textsc{bacon3} does [24]. It induces physical laws given tables of idealized experimental data. For instance, it can induce the general law for ideal gases when it is given "experiments" such as this one:

\[
\text{AND (MOLES 1.0) (TEMPERATURE 300.0) (PRESSURE 300000.0) (VOLUME 0.008320)}
\]

This formal representation describes the experiment in the same way that Winston's representation described a scene in the blocks world (this is not the representation that \textsc{bacon3} uses, by the way). The expression above says that there is one mole of gas at a certain temperature and pressure occupying a certain volume. The goal of \textsc{bacon3} is to find a description that is a generalization of the experiments that it is given. For experiments of this type, the generalization that it induces is:
(AND
  (MOLES N)
  (TEMPERATURE T)
  (PRESSURE P)
  (VOLUME V)
  (TIMES X1 P V)
  (INVISIBLE X1)
  (TIMES X2 N T)
  (INVISIBLE X2)
  (QUOTIENT X3 X1 X2)
  (INVISIBLE X3)
  (CONSTANT X3))

That is, \( PV/NT \) is a constant. This is one way to express the ideal gas law, which is more widely known as \( pV = nRT \), where \( R = 8.32 \). The intermediate results \( PV \), \( NT \) and \( PV/NT \) do not appear in the “scene” described earlier. This is what makes BACON3's job hard. BACON3's method for solving this induction problem is, very roughly speaking, to guess useful invisible objects descriptors and enter their values in the scenes. It might start by forming all binary function on the visible objects, e.g., \( NT \), \( P + V \), \( PP \), \( PT \), etc. Since none of these yield values (invisible objects) that are constant across all the scenes, it tries further compositions: \( NT/PV \), \( NT + V \), \( NTPV \), etc. At this level, it succeeds, since \( PV/NT \) turns out to be the same value, 8.32, in all the scenes. Essentially, BACON3 solves the invisible object problem by choosing a generalization with a minimal number of invisible object designators.

BACON3 is not an incremental inducer. It assumes that it has the total example set at the beginning. There is a reason for this. Any inducer that seeks a generalization with the fewest invisible object designators would clearly want to entertain generalizations with \( N+1 \) invisible object designators only after it had disconfirmed all the generalizations with \( N \) invisible object designators. However, adding an invisible object designator to a disconfirmed generalization won't help it a bit. That is, if \( f(g(x)) \) doesn't match a certain example, then wrapping an \( h(\cdot) \) around it won't help. This means that failure at the level of \( N \) invisible objects doesn't tell one anything about what generalizations to use at the \( N+1 \) level. If the inducer is incremental, and it is at the \( N^l \) level, and the \( M^l \) example exhausts the level then the inducer must start over at the \( N+1 \) level and re-examine all \( M \) examples. It would be better off just

\[ \text{f} \] is the identity function almost everywhere so that \( f(g(x)) = g(x) \) except for the 4 that is the union of \( f \) and \( g \).
waiting until the teacher told it that all the examples were presented, then do a non-
incremental induction. This would require a manner constraint. The teacher would
have to mark the example presentation, and the learner would have to understand
such marks as indicating that it was okay to begin non-incremental induction. As it
turns out, naturally occurring mathematical curricula do employ a manner constraint, but
it is not the one just mentioned. The next subsection describes the one that actually
occurs.

7.2. The show-work felicity condition

In almost all cases, textbooks do not require the student to do invisible object
induction. Instead, whenever the text needs to introduce a subskill that has a mentally
held intermediate result, it uses two lessons. The first introduces the subskill using
special, ad hoc notations to indicate the intermediate results. Figures 7-1 and 7-2
show some examples. Since the intermediate results are written out in the first
lesson, the students need guess no invisible objects in order to acquire the subskill.
The learning of this lesson may proceed as if invisible object designators were banned
from the representation language.

The second lesson teaches the subskill again, without writing the intermediate results.
The second lesson is almost always headed by the key phrase. "Here is a shorter
way to X" where X is the name of the skill. The students are being instructed that
they will be doing exactly the same work (i.e., the path of fact functions is the same).
They are left with the relatively simple problem of figuring out how the new material
relates to the material they learned in the preceding lesson. This kind of learning is
a kind of optimization. They learn how to do the same work with less writing. So, the
normal lessons are "show work" lessons: the learner does invisible object-free
induction. The marked lessons are "hide work" lessons: the learner does optimization
learning. The felicity condition is called "show-work / hide-work" or just show-work for
brevity.

7.3. Parallels between the invisible objects problem and the disjunction problem

The invisible object problem and the disjunction problem are similar in many
respects. (1) Both the invisible objects problem and the disjunction problem are
impossible to solve using unbiased induction if the class of all possible
generalizations allows free use of them then there are far too many generalizations
Figure 7-1: Three formats for column addition obeying the show-work principle. Exercises appear unsolved on the left, solved on the right.
Figure 7-2: Other exercise formats obeying the show-work principle. Exercises appear in normal format on left, in show-work format on right.
consistent with any finite set of examples. Hence, both the disjunction problem and the invisible object problem require extra constraints. (2) Both can be solved trivially with absolute biases: bar their representational devices from the representation language. This is not a viable option in the mathematics domain because the target procedures use both disjunctions and invisible objects. (3) In both cases, a relative bias that works is to minimize the use of the respective devices (i.e., to prefer generalizations with the fewest disjuncts and the fewest invisible object designators). (4) For empirical reasons, minimization biases are not included in the theory. More accurate hypotheses are based on felicity conditions, i.e., tacitly held manner constraints.

8. Task generality

Subtraction is the only task for which I have an extensive data from human learners, so Sierra was developed primarily to simulate the learning of subtraction. In order to make a rough assessment of its task generality, Sierra was given lesson sequences for three new task domains, whose lesson sequences were drawn from a popular elementary mathematics textbook [13]:

1. **Addition** of multi-column, multi-addend problems, such as

\[
\begin{array}{c}
307 \\
+ 620 \\
\end{array}
\]

2. **Multiplication** of multi-column multiplication problems, such as

\[
\begin{array}{c}
307 \\
\times 25 \\
\end{array}
\]

3. **Sixth grade algebra** The skill is to solve linear equations with one occurrence of one unknown, with natural number solutions. At the end of the sixth grade, students are expected to be able to solve \(5(3x + 1) = 20\) but not \(3x + 2x = 10\) or \(5(3x + 1) = 9\).

Sierra learned correct procedures for all three skills, although there are some caveats to this assertion that will be discussed in a moment. Rather than go through Sierra's learning in detail, this section describes the difficulties encountered and the kinds of revisions that would be required resolve them.

One problem is that a **FOREACH** goal type is needed. Given a sequence of objects
of the same type (e.g., a sequence of columns), a FOREACH would execute a subgoal on each object. This new goal type is needed because some naturally occurring lesson sequences are not quite right for learning the tail-recursions that Sierra currently uses to implement loops. Although the lesson sequences can be easily modified, that would harm Sierra's cognitive fidelity. To keep Sierra empirically accurate, a new relative bias towards iteration is needed for skeleton induction. This bias is most effectively implemented by including the FOREACH goal type in the representation language.

A problem was discovered during the last multiplication lesson. At the time of the lesson, the learner can solve single-digit multiplier problems, such as (a) below.

\[
\begin{array}{ccc}
\text{a. } & 5 & 7 \\
\times & 5 & \\
\hline
2 & 8 & 5 \\
\text{b. } & 5 & 7 \\
\times & 1 & 5 \\
\hline
2 & 8 & 5 \\
+ & 5 & 7 & 0 \\
\hline
8 & 5 & 5
\end{array}
\]

Note the use of the scratch marks to indicate carrying. The lesson teaches how to do two-digit multiplier problems, such as (b) above. The addition subproblem presents no difficulties for Sierra, because the initial knowledge state for multiplication includes a multicolumn addition procedure that was learned from the addition lesson sequence.

The problem with the lesson is that it does not use scratch marks to indicate carrying. This causes two difficulties for Sierra. First, the textbook does not include an optimization lesson to teach how to suppress carry marks (using examples such as (c) below).

\[
\begin{array}{ccc}
c. & 5 & 7 \\
\times & 5 & \\
\hline
2 & 8 & 5
\end{array}
\]

Perhaps teachers present such examples on their own initiative, without the guidance of a textbook lesson. A more serious problem is that even if there were an optimization lesson, there is no easy way to modify the multiplication procedure in order to suppress the scratch marks. The procedure has a loop, which iterates leftward through the top row, multiplying the row's digits by the single-digit multiplier. A carried digit is written during one invocation of the loop body and read during another invocation of the loop body (the next one, in fact). For this dataflow to...
happen applicatively without writing on the page, the second invocation must somehow be called from the first invocation. There is no way to achieve this without a complete overhaul of the procedure's calling structure. Sierra cannot do such a major overhaul. I'm not sure what students do because I have no data for multiplication. However, I suspect that they do not radically overhaul their understanding of the procedure just to suppress scratch marks. I suspect that they use their fingers to hold the carried digit between multiplies, or they use some short-term memory resource to do so. Implementing the latter possibility would entail making the representation language non-applicative, which would make parsing much more complex. On the other hand, if students use their fingers, then the hypotheses and representation can remain intact, but the representation of the state of a multiplication problem would have to have a "hand" added to it. Either modification would require significant enough programming that I simply stopped Sierra's traversal of the multiplication sequence at this lesson. It was the last one, anyhow.

The most significant task dependencies concerned notational syntax. This is not surprising, since the four tasks employ quite different notations. Sierra's treatment of problem states and their syntax has not yet been discussed in this article (see [44]). The basic idea, however, is simple to present. Problem states are represented as letters, digits and lines situated on a Cartesian plane, and a two-dimensional context-free grammar is used to parse them. This technique failed in some cases. For instance, given the expression "5-x", the minus sign must be viewed two ways: as a prefix for the term following it and as an infix operator that separates the two terms. So there are two parse trees for "5-x," four for "2+(5-x)," and so on. Sierra's context-free grammar technique is combinatorially explosive. A better solution would be to redesign the parser and pattern matcher so that they keep local ambiguities local. This might, in fact, be the first step toward an interesting theory of the interpretation of mathematical notation.

To sum up, there were two main difficulties in getting Sierra to learn other skills than subtraction. (1) The dataflow architecture is incomplete. Some globally bound resource (e.g., fingers, short-term memory) is needed to do carrying without scratch marks. (2) The notational grammars are not quite expressive enough. People seem to view the same problem state several ways, a facility that Sierra's grammar system does not adequately support.
9. Concluding remarks

Until now it seemed that to be successful, an inducer had to use either absolute or relative biases. To put it differently, successful inducers have either been partially blind or strongly prejudiced. Sierra is a demonstration that there is a third way. An inducer can be successful if it receives a well-structured example sequence whose structure it understands. That is, the example sequence obeys certain manner constraints and thereby encodes information about the structure of the target generalization. The learner takes advantage of these constraints in order to recover the “message” that is encoded in the form of the example sequence. In human learning situations, if neither teacher nor learner are aware of the manner constraints then they warrant the name felicity conditions. So a successful inducer is either partially blind, strongly prejudiced, and/or felicitiously taught.

9.1. Some speculations on manner constraints and felicity conditions

In a certain sense, manner constraints may be optimal strategies for knowledge communication. For instance, in order to solve the learner’s disjunction problem, the teacher’s optimal strategy would be to point to a node in the learner’s knowledge structure and say “disjoin that node with the following subprocedure: ...”. Clearly, this is impossible. So the teacher says the next best thing. “Disjoin some node with the following subprocedure: ...”. The learner must figure out which node to disjoin because the teacher can’t point to it. But the learner now knows that some disjunction is necessary and that the examples following the teacher’s command will determine its contents. If it were not for the exigencies of school scheduling, this would be perhaps the optimal information transmission strategy. However, lessons must be about an hour long. This means that only some of the lesson boundaries will correspond to the teacher’s command to start a new disjunction. The other lessons will finish up a previous lesson. In short, the optimal, feasible manner constraint for disjunctive information transmission could well be one-disjunct-per-lesson.

There is a great deal of complaining about the so-called knowledge acquisition bottleneck in developing expert systems. It seems to be quite difficult to get human experts to formalize their expertise as, e.g., production rules. One often heard solution is to have the system learn the knowledge on its own, e.g., by discovery or by analogy (e.g., [23, 28]). I tend to agree with Simon [42], who predicts that programming will always be the most effective way to “educate” computer. However
if Simon and I are wrong, and machine learning does hold promise as a solution to the knowledge acquisition bottleneck, then examining how human experts acquire their knowledge is a good research heuristic. Even a cursory examination shows that most human experts didn’t discover their knowledge or infer it; they learned it from a mentor, either in school or as an apprentice. A good mentor is careful about selecting tasks for the student that are appropriate for the student’s current state of knowledge. That is, the instruction is not a randomly ordered sequence of tasks, but a carefully structured one. Yet few researchers are trying to get an expert systems to learn from the kinds of structured instruction that human experts receive. Such a system would take advantage of the format that its mentor places on the instruction. The present research, in its explication of felicity conditions, should be helpful in building such a knowledge acquisition system. Such a system will be easier for human experts to educate than present systems because the experts, many of whom are experienced teachers, are more familiar with formatting their knowledge as lesson sequences than as production rules.

9.2. Summary
Putting speculation aside, I’ll review the techniques that have been presented. The formost is one-disjunct-per-lesson. Because of it, Sierra’s design is quite simple. The procedure induction problem is reduced to a series of subprocedure induction problems, one per lesson. Subprocedure induction reduces to three subproblems: skeleton induction, pattern induction and function induction.

Skeleton induction is performed by parsing the action sequences top-down as far as possible and bottom-up as far as possible. In neither case will parsing yield a complete parse tree. To complete the parse tree, a new piece of tree structure must be built to connect the top-down parse to the bottom-up parse. Any such structure is a candidate skeleton. This dual-parser calculation is done on each example in the lesson, yielding one set of skeletons per example. These sets are intersected, yielding the skeleton(s) that are consistent with all the examples in the lesson.

Footnote: Expert programmers, particularly older ones, are notorious exceptions. Many of them acquired their expertise without instruction. Perhaps this accounts for the widespread myth in the knowledge engineering community that domain experts become experts without instruction.
Pattern induction is performed by a standard technique. Mitchell's version space algorithm [29].

Function induction can employ a brute force generate-and-test algorithm because there is a manner constraint that simplifies the problem. The show-work constraint says that examples must "show all the work" when introducing a new subprocedure. That is, intermediate results must be written on the example where the learner can "see" them. Because a composition (nest) of two primitive functions has an intermediate result that is not written down, composite functions can not be introduced during a normal lesson. Consequently, the learner only has to consider primitive functions and not compositions of functions when it does function induction. The show work constraint makes function induction almost trivial.

9.3 Discussion

The "per lesson" part of one-disjunct-per-lesson is slightly misleading. It turns out that Sierra could get along just fine without lessons. The skeleton induction algorithm is the one component that controls the introduction of disjunctions. The algorithm fails only if an example introduces two or more disjuncts. In particular, just by omitting skeleton intersection, Sierra could get along fine without lesson boundaries. Its two manner constraints would become (1) at most one disjunct per example, and (2) an example is induced without invisible object designators unless it is marked as a same-work example. A per-example learner might be more appropriate for some knowledge acquisition tasks than the current per-lesson version.

It bears reiterating that the dual-parser technique that performs skeleton induction is simple and efficient because the procedure representation language is applicative. If the language allowed side-effects, such as storage of information in global buffers or variables, then parsing would be much more difficult.

As mentioned earlier Sierra consists of three induction algorithms, for respectively, skeletons, patterns and functions. This three-way decomposition may apply to other learning tasks than procedure learning. The application of the induction algorithms is limited only by the topology of the knowledge representations and not, of course, by what those knowledge representations denote. I would expect these algorithms to apply, for instance, to other learning-by-explanation tasks. Explanations of stories often
feature hierarchical structures similar to the calling structures of procedures. Such explanations are generated by instantiating and composing schemata. These schemata are analogous to subprocedures. Schemata often contain restrictions on slots that are equivalent to patterns. If these equivalences hold, then learning new schemata could be accomplished by the same techniques that are used here to learn subprocedures.

To take a second example, electronic circuits and other engineered devices are often designed to have a hierarchy of modules. This hierarchy corresponds to the calling hierarchy of procedures. Patterns and functions may have analogs in device designs as well. If so, then the techniques presented here might suffice to learn how to analyze devices. In short, the three-way decomposition of the induction problem into induction of hierarchies, induction of pattern-like constructs, and induction compositions may be quite generally applicable.

Sierra's skeleton induction is a form of context-free grammar induction, since the skeleton of a whole AOG is precisely a context-free grammar. Skeleton induction can be used to induce grammars, as long as the learner's example sequence conforms to one-disjunct-per-example. On the other hand, if one-disjunct-per-example is not an appropriate manner constraint for some domain, then some other grammar induction algorithm may be employed to perform skeleton induction, while the other two induction algorithms can remain relatively unchanged. (For reviews of grammar induction, see [15 9].)

A beneficial consequence of one-disjunct-per-lesson is that rule patterns are pure conjunctions. This means that a non-heuristic, complete induction algorithm, based on Mitchell's version space technique [29], can be employed to induce the conditions. However, it turns out to be infeasible to use just the version space technique. For empirical reasons, Sierra must use non-toy patterns. A single pattern may have 50 variables and 200 relations. For patterns of such sizes, induction is just not practical without further constraints on patterns. Fortunately, there are several well-motivated constraints available in this domain. The two mentioned at the end of section 6 seem likely to be useful outside the present domain.

The invisible objects problem was somewhat of a surprise. It took a long time to figure out that function induction had an inherent problem: i.e., that Sierra wasn't a victim of incomplete sets of examples! It took even longer to become convinced that...
there is no least-commitment, incremental algorithm to solve it, such as the version space algorithm. If an inducer cannot ban invisible object designators, there seem to be only two ways to get around the invisible-objects problem: BACON3's non-incremental induction (see section 7) or Sierra's show-work manner constraint. Perhaps more research will find other techniques. The problem of inducing invisible object designators has received little attention from machine learning.
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