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REARRANGEMENT OF SUBPROGRAMS ON TAPE TO REDUCE SELECTION TIME

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ABSTRACT

A number of subprograms are read from a magnetic tape into the storage of an electronic computer during the execution of a computing job. The order in which these subprograms are arranged on the tape is being studied. A number of jobs are executed during a certain period of time using different groups of subprograms from the tape. It is our aim to re-arrange the subprograms on the tape in order to reduce the time it takes to read the subprograms into storage.
FOREWORD

An expression for the total time of selection of the subprograms is determined. It is not the aim of this paper to find the ideal distribution in order to minimize the selection time. Instead a method of trial and error is used to improve the distribution sufficiently for practical purposes. Exchanges between two subprograms at a time are made. If an exchange does not reduce the time of selection a re-exchange is made. A program in FORTRAN for the IBM 7090 is supplied to complete the task. A practical example is given.
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A subprogram on magnetic tape consists of a number of records with record gaps between them. Each record contains a certain number of "words". A word has a fixed length on the tape. Hence, the length of a subprogram may be measured in terms of the number of words it contains within each record, allowing for record gaps. Let us suppose we have n subprograms \( s_1, s_2, \ldots, s_n \) of lengths \( t_1, t_2, \ldots, t_n \) words on the tape. The "head", which reads the words from the tape, is located at the beginning of the tape, and this is also the beginning of \( s_1 \). To select the subprogram \( s_m \) \((m > 1)\) the head moves over \( (t_1 + t_2 + \ldots + t_{m-1}) \) words. After the first subprogram \( s_m \) has been read into storage, the head is located at the beginning of \( s_{m+1} \). To select the next subprogram \( s_r \), the head moves over

a) \((t_{m+1} + t_{m+2} + \ldots + t_{r-1})\) words if \( r > m + 1 \)

b) \((t_m + t_{m-1} + \ldots + t_r)\) words if \( r < m \).

It takes no time to select \( s_{m+1} \) \((r = m + 1)\). If the same subprogram is used again \((r = m)\) we can ignore the corresponding selection time, since it cannot be reduced by changing the distribution. In this fashion, the head moves forwards and backwards over the tape to select all the subprograms for a job.

We assume that the head moves forwards and backwards over the tape with the same constant speed. Hence, the total selection time may be measured in terms of the number of words passed by the head without reading.
TOTAL TIME OF SELECTION

There are \( n \) subprograms \( s_i (i = 1, \ldots, n) \) on the tape of lengths \( t_i (i = 1, \ldots, n) \). During a certain period of time, a number of jobs have been executed and the numbers \( f_{1k} (i, k = 1, 2, \ldots, n) \) and \( f_i (i = 1, 2, \ldots, n) \) have been recorded. \( f_i \) is the number of times the subprogram \( s_i \) was selected first during the execution of a job. \( f_{1k} \) is the number of times the selection of \( s_i \) was followed by the selection of \( s_k \). To find the total time of selection, \( T \), we form the following expressions:

\[
f_i (t_1 + t_2 + \ldots, + t_{i-1}), i > 1
\]

\[
f_{1k} (t_{i+1} + t_{i+2} + \ldots, + t_{k-2} + t_{k-1}), k > i + 1
\]

\[
f_{k1} (t_1 + t_{i+1} + \ldots, + t_{k-1} + t_k), k > i
\]

and find the sum of all possible terms:

\[
T = \sum_{i=2}^{n} f_i \sum_{m=1}^{i-1} t_m + \sum_{i=1}^{n-2} \sum_{k=i+2}^{n} f_{1k} \sum_{m=i+1}^{k-2} t_m + \sum_{k=2}^{n} \sum_{i=1}^{k-1} f_{k1} \sum_{m=i}^{k-1} t_m
\]

We change this expression for faster computation. The third term may be changed to

\[
\sum_{k=2}^{n} f_{k1} \sum_{m=1}^{k} t_m + \sum_{i=2}^{n-1} f_{n1} \sum_{m=i}^{n-1} t_m + \sum_{k=3}^{n} \sum_{i=2}^{k-1} f_{k1} \sum_{m=i}^{k-1} t_m
\]

The second term may be changed to

\[
\sum_{i=1}^{n-2} f_{1,i+2} \sum_{i=1}^{n-3} f_{i+1} \sum_{m=i+1}^{n-3} t_m
\]

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The first term may be written as

\[ f_2 \cdot t_1 + \sum_{k=2}^{n-1} f_{k+1} \sum_{m=1}^{k} t_m. \]

Hence,

\[ T = f_2 \cdot t_1 + \sum_{k=2}^{n-1} (f_{k+1} + f_{k+1}) \sum_{m=1}^{k} t_m + f_m \sum_{m=1}^{n} t_m \]

\[ + \sum_{i=2}^{n-1} f_{m+1} \sum_{m=1}^{n} t_m + \sum_{i=2}^{n-1} f_{i-1,i+1} \cdot t_i \]

\[ + \sum_{k=3}^{n-1} \sum_{i=3}^{k} f_{k,i-1} \sum_{m=1}^{k} t_m + \sum_{i=3}^{n-1} \sum_{k=1}^{n-1} f_{i-2,k+1} \sum_{m=1}^{k} t_m. \]

The last two terms add up to

\[ \sum_{i=3}^{n-1} \sum_{k=1}^{n-1} (f_{k,i-1} + f_{i-2,k+1}) \sum_{m=1}^{k} t_m. \]

This form of T has been used in the subroutine FINDT to compute T.
III  REDUCTION OF THE TIME OF SELECTION

To reduce the time of selection $T$ we change the order in which the
subprograms are on the tape. We compute at first $T_0$ for the original di-
tribution. Then we exchange two subprograms $s_i$ and $s_k$ and compute $T_1$. If
$T_1 \geq T_0$, we exchange $s_i$ and $s_k$ again. Otherwise we make the next exchange.
We let $i = 1, 2, \ldots, n-1$ and $k = i + 1, i + 2, \ldots, n$ until a permanent change
has been made. Then we start again with $i = 1$ and $k = 2$. This method will
lead to a distribution which cannot be improved by the exchange of any two
subprograms. For $n > 50$ it may take hours of computer time before the final
distribution has been found, because the changes become more and more
insignificant and harder to find as time goes on. It is advisable to run
the program for 10 minutes at a time and break it off as soon as the
improvement becomes insignificant. In an example with $n = 65$ the time $T$
was reduced by 26% during the first 10 minutes of computer time, by 7%
during the next 30 minutes run, by 4% during the next 30 minutes run and
by 1% during the next 20 minutes run. In an example with $n = 71$ the time
$T$ was reduced by 5.5% in 7 minutes and by 1.6% in another 30 minutes. In
this case the program worked on an improved distribution of subprograms.

To enable a restart of the program without a change of the input data
$f_{ik}, f_i, \text{and } t_i$, the program transforms the matrices $f_{ik}, f_i, \text{and } t_i$ to any
given order of distribution $s_a, s_b, s_c, \ldots, s_x, s_y, s_z$. We simply type the
integers $a, b, c, \ldots, x, y, z$ on certain data cards, and the program
establishes this order and continues making improvements.
A program in FORTRAN for the IBM 7090 is provided. The arrays \( F(N,N), \) \( T(N) \) and \( S(N) \) contain the matrices \( f_{i,j}, t_i \) and \( f_i \), where \( N = n \). The array \( MU(N) \) contains the integers \( a, b, c, \ldots, x, y, z \) mentioned in 3. The array \( NU(N) \) is initially made to agree with \( MU(N) \) and then subjected to changes, indicating the current order of distribution,

\[ S_{NU(1)}, S_{NU(2)}, \ldots, S_{NU(N)}. \]

The total time of selection is denoted by \( TT \), the change in time due to an exchange of subprograms by \( DT \). The following subroutines are used:

- **READ**: This subroutine reads from the data cards the arrays \( T, F, MU \) and \( S \). \( F \) must be typed columnwise on the cards, e.g. \( f_{11}, f_{21}, \ldots, f_{n1}, f_{12}, f_{22}, \ldots, f_{nn}, \ldots \), where \( f_{11} = f_{n2} = \ldots = f_{nn} = 0 \).
- **EXCH**: This subroutine exchanges the subprograms \( S_{L_1} \) and \( S_{L_2} \). It makes the necessary changes in the matrices \( F, T, NU \) and \( S \).
- **HEAD**: This subroutine prints the total time \( TT \) and the order of distribution \( NU \) as a heading for each page.
- **ORDER**: This subroutine prints the total time \( TT \) and the order of distribution \( NU \) whenever 10 changes have been made.
- **INIT**: This subroutine transforms to the order of distribution \( NU = MU \).
- **FINDT**: This subroutine evaluates \( TT \).
- **PLAN**: This subroutine exchanges two subprograms at a time and prints for each permanent change the numbers \( L_1, L_2 \) and \( DT \), indicating that the \( L_1 \)th and \( L_2 \)nd subprograms have been exchanged and that \( TT \) has been increased by \( DT \).

The main program calls \( INIT, PLAN, ORDER \) and \( EXIT \).

\( N \) is taken to be 71 in the following listing of the program.
ORDER

MAIN PROGRAM
1 CALL INIT
2 CALL PLAN
3 CALL ORDER
4 CALL EXIT
END(1,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0)

SUBROUTINE INIT
DIMENSION F(71,71),T(71),NU(71),MU(71),S(71),LU(71)
COMMON F,T,T1,D1,T1,L1,L2,MU,S

N=71
2 CALL READ
17 DO 19 I=1,N
18 NU(I)=I
19 LU(I)=I
20 N1=N-1
21 DO 27 I=1,N1
22 L1=I
23 L2=MU(I)
24 L2=LU(L2)
25 IF(L2-L1)24,27,26
26 CALL EXCH
27 LU(I)=L2
28 CALL FINDT
31 CALL HEAD
RETURN
END(1,0,0,0,0,0,1,0,0,0,0,0,0,0,0,0,0)
SUBROUTINE EXCH

DIMENSION F(71,71),T(71),NU(71),MU(71),S(71)
COMMON F,T,TT,DT,NU,LN,N,L1,L2,MU,S

5 DO6 I=1,N
   A=F(L1,I)
   F(L1,I)=F(L2,I)
   F(L2,I)=A
   DO10 I=1,N
   A=F(I,L1)
   F(I,L1)=F(I,L2)
   F(I,L2)=A
10   A=T(L1)
   T(L1)=T(L2)
   T(L2)=A
   A=NU(L1)
   NU(L1)=NU(L2)
   NU(L2)=A
15   A=S(L1)
   S(L1)=S(L2)
   S(L2)=A
19   RETURN
END(1,0,0,0,0,0,1,0,0,0,0,0,0,0)

SUBROUTINE READ

DIMENSION F(71,71),T(71),NU(71),MU(71),S(71)
COMMON F,T,TT,DT,NU,LN,N,L1,L2,MU,S

1 READ INPUT TAPE 5,2,T
2 FORMAT(16F5.0)
3 READ INPUT TAPE 5,4,F
4 FORMAT(24F3.0/24F3.0/23F3.0)
5 READ INPUT TAPE 5,6,MU
6 FORMAT(40I2)
7 READ INPUT TAPE 5,4,S
RETURN
END(1,0,0,0,0,0,1,0,0,0,0,0,0,0,0)
SUBROUTINE FINDT
DIMENSION F(71,71),T(71),NU(71),MU(71),S(71)
COMMON F,T,TT,DT,NU,LN,N,L1,L2,MU,S

1 N1=N-1
2 TT=0.0
3 A=T(1)
   TT=TT+A*S(2)
4 DO 6 I=2,N1
5 A=A+T(I)
6 TT=TT+A*(F(I,1)+S(I+1))
   A=A+T(N)
   TT=TT+A*F(N,1)
7 DO 9 I=2,N1
8 A=A-T(I-1)
9 TT=TT+A*F(N,I)
10 DO 14 I=3,N1
11 A=T(I-1)
12 DO 14 J=1,N1
13 A=A+T(J)
14 TT=TT+A*(F(I-2,J+1)+F(J,I-1))
15 DO 16 I=2,N1
16 TT=TT+T(I)*F(I-1,I+1)
17 RETURN
END(1,0,0,0,0,0,0,0,0,0,0,0,0)

SUBROUTINE HEAD
DIMENSION F(71,71),T(71),NU(71),MU(71),S(71)
COMMON F,T,TT,DT,NU,LN,N,L1,L2,MU,S
1 WRITE OUTPUT TAPE 6,2,TT,NU
2 FORMAT(1H1,9X,5HTIME=,E16.8,12HTHE ORDER IS/(1X,I2,39I3))
3 LN=0
4 RETURN
END(1,0,0,0,0,0,0,0,0,0,0,0)
SUBROUTINE PLAN

DIMENSION F(71,71), T(71), NU(71), MU(71), S(71)

COMMON F, T, TT, DT, NU, LN, N, L1, L2, MU, S

A=TT

1  N2=N-1
2  DO 16 I=1, N2
3     L1=I
4  L3=L1+1
5  DO 16 J=L3, N
6     L2=J
7    CALL EXCH
8    CALL FINDT
9    DT=TT-A
10   IF(DT) 7, 16, 16
11    WRITE OUTPUT TAPE 6, 9, L1, L2, DT
12    FORMAT(10X, 1OHCHANGE L1=, 14, 8H AND L2=, 14, 4H DT=, E16.8)
13    A=TT
14   10   LN=LN+1
15   11   IF(50-LN) 12, 12, 13
16    CALL HEAD
17   12   IF(XEQDF(LN, 10)) 2, 14, 2
18    CALL ORDER
19   13   GO TO 2
20   14   CALL EXCH
21    RETURN

END(1, 0, 0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0)

SUBROUTINE ORDER

DIMENSION F(71,71), T(71), NU(71), MU(71), S(71)

COMMON F, T, TT, DT, NU, LN, N, L1, L2, MU, S

1 WRITE OUTPUT TAPE 6, 2, TT, NU
2 FORMAT(8X, 13HNEW ORDER. TT=, E16.8/(1X, I2, 39I3))
3 LN=LN+1
4 RETURN

END(1, 0, 0, 0, 0, 0, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0)

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