Modified Cheeger and Ratio Cut Methods Using the Ginzburg-Landau Functional for Classification of High-Dimensional Data

Ekaterina Merkurjev*, Andrea Bertozzi**, Xiaoran Yan***, and Kristina Lerman***

*Department of Mathematics, University of California, San Diego  
**Department of Mathematics, University of California, Los Angeles  
***Information Sciences Institute, University of Southern California, Marina del Rey

Email: emerkurj@ucsd.edu, bertozzi@math.ucla.edu, xiaoran@isi.edu, lerman@isi.edu.

Abstract

Recent advances in clustering have included continuous relaxations of the Cheeger cut problem and those which address its linear approximation using the graph Laplacian. In this paper, we show how to use the graph Laplacian to solve the fully nonlinear Cheeger cut problem, as well as the ratio cut optimization task. Both problems are connected to total variation minimization, and the related Ginzburg-Landau functional is used in the derivation of the methods. The graph framework discussed in this paper is undirected. The resulting algorithms are efficient ways to cluster the data into two classes, and they can be easily extended to case of multiple classes, or used on a multiclass data set via recursive bipartitioning. In addition to showing results on benchmark data sets, we also show an application of the algorithm to hyperspectral video data.

Keywords: classification, spectral clustering, Cheeger cut, ratio cut, graphs, Ginzburg-Landau functional, total variation, graph Laplacian, Nyström extension method.

1 Introduction

1.1 General Problem and Background

Total variation has become increasingly more central to inverse problems in imaging. The classic paper by Rudin, Osher and Fatemi [32] sparked interest in total variation by applying it very successfully to denoising problems. Later, similar techniques were applied to areas such as inpainting, image restoration and blind deconvolution. For example, in [11], Chan et al. present a blind deconvolution method using total variation (TV), in which they use an alternating minimization implicit iterative scheme. In addition, while [9] proposes a nonlinear primal-dual method for TV-based image restoration, [10] formulates a higher order TV-based restoration algorithm.
More recently, total variation has been used successfully in a graphical framework. In [20], the authors outline the graphical framework, and apply it to image inpainting and regularization using the total variation functional. Another example is [27], in which Lou et al. describe a nonlocal method for image recovery using total variation. In [19], Gilboa and Osher not only formulate a nonlocal TV method for image regularization, but also extend it to supervised segmentation, where the user marks small regions of known information, and the desired class is recovered. Since total variation is becoming more popular in machine learning problems in which data is represented in a graphical framework, the minimum cut problem, which becomes minimizing total variation of an indicator function in a graph setting, becomes a natural mechanism with which to proceed to achieve accurate data classification.

We thus consider the problem of minimizing the cut as a way of partitioning the set $X$ into two clusters. Let $G = (V, E)$ represent an undirected graph with the set of vertices $V$ and set of edges $E$, and consider a target set $X$ of size $n$ embedded in a graph $G$. A weight function is defined on the set of edges, and represents a measure of similarity between the vertices it is connecting. As usual, the degree of a vertex $x$ is formulated as $d(x) = \sum_{y \in V} w(x, y)$. Denote by $D$ the diagonal matrix containing the degrees of vertices as diagonal entries, and let $W$ denote the similarity matrix containing the weight function values. The minimum cut problem is to find the set $S \subset V$ such that the following value is minimized:

$$\text{cut}(S, \bar{S}) = \sum_{x \in S, y \in \bar{S}} w(x, y).$$

Here, $\bar{S}$ indicates the complement. Intuitively, the weights between vertices of different clusters should be small. However, minimizing the cut leads to an undesirable solution, first and foremost because the trivial solution would be to cluster all the points into one set, thus giving the cut the value of zero. Even if a non-trivial partition condition is enforced, the solution tends to isolate individual vertices from the rest of the graph. Usually, what is wanted is for the two clusters to be relatively close in size, and one solution is to include a normalization of the cut.

One normalization is the ratio cut. The problem is to find a subset $S$ of $V$ such that

$$\text{RatioCut}(S, \bar{S}) = \text{cut}(S, \bar{S}) \left( \frac{1}{|S|} + \frac{1}{|\bar{S}|} \right)$$

is minimized. This is a NP hard discrete problem [38]. One way to simplify it would be to allow the solution to take arbitrary values in $\mathbb{R}$. A clever reformulation of the original task and relaxation of the binary constraint leads to the problem of finding the argument of the following optimization task:

$$\min_{u \in \mathbb{R}^n} \langle u, Lu \rangle, \quad \text{such that } u \perp 1, \quad ||u||^2 = n, \quad (1)$$

where $L$ is the graph Laplacian $D - W$. We emphasize the fact that the above problem obtains a real-valued solution instead of a discrete-valued one. To solve (1.1), one can apply the Raleigh-Ritz theorem, and the solution is given by the second eigenvector of the graph Laplacian. To obtain a binary solution, which indicates the binary partition, one can use several methods, the simplest of which is thresholding.

Another normalization of the cut is the normalized cut. If we let $\text{vol}(S) = \sum_{x \in S} d(x)$, where $d(x)$ represents the degree of vertex $x$, the problem is modified to find a subset $S$ of $V$ such that

$$\text{Ncut}(S, \bar{S}) = \text{cut}(S, \bar{S}) \left( \frac{1}{\text{vol}(S)} + \frac{1}{\text{vol}(\bar{S})} \right).$$
is minimized. This is a NP hard discrete problem [38] as well. We simplify it by allowing the solution to take arbitrary values in \( \mathbb{R} \). A reformulation of the normalized cut problem and relaxation of the binary constraint lead to the following task of finding the argument of the minimum of the task:

\[
\min_{u \in \mathbb{R}^n} \langle u, Lu \rangle, \quad \text{such that } Du \perp 1, \quad \langle u, Du \rangle = vol(V),
\]

where \( vol(V) = \sum_x d(x) \). One can again apply the Raleigh-Ritz theorem, and the solution is given by the second eigenvector of the random walk Laplacian \( D^{-1}L \). Thresholding can be used to obtain a binary solution which represents the desired partition.

The last normalization to be discussed is the Cheeger cut, which involves finding a subset \( S \) of \( V \) that minimizes

\[
h(S, \bar{S}) = \frac{\text{cut}(S, \bar{S})}{\min(|S|, |\bar{S}|)}.
\]

If volume of a set is used instead, one obtains the normalization

\[
h^v(S, \bar{S}) = \frac{\text{cut}(S, \bar{S})}{\min(\text{vol}(S), \text{vol}(\bar{S}))},
\]

which is also known as conductance. Let \( \phi(G) = \min_{S \subseteq V} h^v(S, \bar{S}) \). The Cheeger cut and the second eigenvalue \( \lambda_2 \) of the matrix \( D^{-\frac{1}{2}}LD^{-\frac{1}{2}} \) can be related via the following Cheeger inequality [12]:

\[
\frac{\phi(G)^2}{2} \leq \lambda_2 \leq 2\phi(G).
\]

The authors of [18] present a generalized version of it.

We have seen that the ratio cut and the normalized cut problems can be formulated in a relaxed setting, with non-binary solutions given by the second eigenvectors of the Laplacian and the random walk Laplacian, respectively, with the binary answer obtained by thresholding. However, for an arbitrary number of classes, this technique is too simple.

The method of spectral clustering computes the \( k \) clusters using the steps:

1. Formulate the data set in a graph setting.
2. Compute either the Laplacian \( (L = D - W) \) or the random walk Laplacian \( (L_{rw} = D^{-1}L) \).
3. Compute the first \( k \) eigenvectors \( \{v_1, v_2, ..., v_k\} \) of the Laplacian (or the random walk Laplacian).
4. Let \( U \) be the matrix containing the vectors \( \{v_1, v_2, ..., v_k\} \) as columns.
5. Cluster the rows of the matrix \( U \) with the \( k \)-means algorithm into \( k \) clusters.

Note that the \( k \)-means algorithm for finding \( k \) clusters proceeds iteratively by first choosing \( k \) means and then assigning each point to the mean that is the “closest” to the point. The mean of each cluster is then recalculated. The iterations continue until there is little change from one iteration to the next.

There is a large literature on the binary partitioning problem on graphs. Here, we review work that relates to solving the cut problem under some normalization condition. In [8], the authors present a generalized version of spectral clustering using the graph p-Laplacian. They show that, as \( p \to 1 \), the cut resulting from thresholding the second eigenvector of the graph p-Laplacian is the solution to the Cheeger cut problem. An efficient scheme to calculate the eigenvector is then introduced. In [23], Hein et al. show that some constrained optimization problems
can be formulated as nonlinear eigenproblems. The authors then describe a generalization of the inverse power method which converges to nonlinear eigenvectors. This method is applied to spectral clustering and sparse PCA. Moreover, in [13], Chung describes the heat kernel as the pagerank of a graph, and uses it and the local Cheeger cut inequality to establish a fast graph partitioning algorithm. In [33], Spielman et al. present nearly-linear time algorithms for graph partitioning, graph sparsification, and solving linear systems.

Recent work by Bresson, Szlam, Laurent, von Brecht, et al. includes several important papers related to clustering. In [35], Szlam et al. present a relaxation of the Cheeger cut problem and then show similarities of the energy of the relaxed problem and well studied energies in image processing. An algorithm based on the split-Bregman method [21] is then developed to minimize the proposed energy. Authors of [3] describe two procedures solving the relaxed Cheeger cut problem. The first algorithm is a novel steepest descent approach, and the second one is a modified inverse power method. Some convergence results are also given. In [5], Bresson et al. develop another version of the method shown in [3] using a new adaptive stopping condition. The result is an algorithm that is monotonic and much more efficient than before. Note that the mentioned methods are unsupervised, which do not incorporate any known data. As opposed to these kinds of algorithms, semi-supervised optimization approaches involve the inclusion of a fidelity term, a fitting term to known data.

Multiclass extensions relating to the cut have also been proposed. One approach is to use recursion and thus solve a collection of binary problems. However, other ideas have also been introduced. The authors of [4] present a framework for multiclass total variation clustering that does not use recursion. They formulate the Cheeger energy in a multiclass setting, and then relax the energy in a continuous setting. This results in an optimization problem involving total variation, which is then solved using the proposed proximal splitting algorithm. In [6], a multiclass extension of the result of [35] is derived. The method deals with learning several classes simultaneously with a set of labels, and is made even more efficient by the usage of fast $L^1$ solvers, designed for the total variation semi-norm. Finally, we mention the work of Hu et al. [25] in which the authors formulate modularity optimization as a minimization problem using an energy functional that contains a total variation term.

In addition to using the cut as a basis for our method, one technique used in this work is one of spectral methods. Spectral methods are a type of approach used to solve differential equations. Their idea is to write the solution of the differential equation, and perhaps other relevant parts of the problem, as a sum of basis functions. The coefficients of the sum are then calculated so that the differential equation is satisfied (or satisfied to the best ability). In this work, we use the eigenfunctions of a Laplacian as basis functions.

1.2 Our Contributions

Specifically, our contributions are:

1. We show how to use the graph Laplacian to solve the fully nonlinear ratio cut problem. We call this algorithm the Modified Ratio Cut Method (RCCM), and present it in Section 3. Two approaches are presented; one approach links a modified ratio cut formulation to an optimization problem involving total variation, and the other one deals with solving the ratio cut formulation directly.
2. We show how to use the graph Laplacian to solve the fully nonlinear *Cheeger cut* problem. We call this algorithm the Modified Cheeger Cut Method (MCCM), and present it in Section 4.

3. A main advantage of these methods is efficiency, while being able to maintain the accuracy of that of the state-of-the-art. The efficiency is partly achieved through the use of the Nyström extension method [15,16]. The methods can be easily extended to case of multiple classes, or used on a multiclass data set via recursive bipartitioning.

2 Graphical Framework

2.1 Fundamentals

In this paper, we consider a graphical framework with an undirected graph \( G = (V, E) \), where \( V \) is the set of vertices and \( E \) is the set of edges. Earlier in Section 1, we introduced some important concepts related to the framework, such as the weight function, the degree and the similarity matrix. One advantage of using a graphical framework is that it allows one to be non-local and take into account the relationship between any two nodes in the data set. In the case of image processing, this makes it easier for one to capture repetitive structure and texture, as shown in [29]. Furthermore, the framework is also very general, and can be easily applied to any data set.

It is possible to introduce some common mathematical operators in a graphical setting. In this section, we will only be concerned with the graph version of the differential Laplace (\( \Delta \)) operator. Although many versions exist, three popular ones are the unnormalized Laplacian \( L = D - W \), symmetric Laplacian \( L_s = D^{-\frac{1}{2}}LD^{-\frac{1}{2}} \), and the random walk Laplacian \( L_{rw} = D^{-1}L \). The graph Laplacians have the following easily shown equations:

1) \( L \) and \( L_s \) are symmetric.
2) \( L, L_s \) and \( L_{rw} \) are positive semi-definite.
3) \( L, L_s \) and \( L_{rw} \) have non-negative, real-valued eigenvalues \( 0 = \lambda_1 \leq \lambda_2 \leq ... \leq \lambda_n \).
4) The smallest eigenvalue of \( L \) is 0; the corresponding eigenvector is just a constant vector.
5) \( \lambda \) is an eigenvalue of \( L_{rw} \) with eigenvector \( u \) if and only if \( \lambda \) is an eigenvalue of \( L_s \) with eigenvector \( w = D^{\frac{1}{2}}u \).

Note that the multiplicity of eigenvalue 0 equals the number of connected components in the graph.

2.2 Graphical Framework, Extended

We outline the graphical framework in more detail here, giving general definitions of other operators. The operators are defined similarly to [22,37], where the justification for these choices is explained.

Let \( n \) be the number of vertices in the graph and let \( V \cong \mathbb{R}^n \) and \( E \cong \mathbb{R}^{\frac{n(n-1)}{2}} \) be Hilbert spaces (associated with the set of vertices and edges, respectively) defined via the inner products:

\[
\langle u, \gamma \rangle_V = \sum_x u(x) \gamma(x) d(x)^r,
\]

\[
\langle \psi, \phi \rangle_E = \frac{1}{2} \sum_{x,y} \psi(x,y) \phi(x,y) w(x,y)^{2q-1},
\]

\[
\langle u, \gamma \rangle_{E^2} = \sum_x u(x) \gamma(x),
\]
∀u, γ ∈ V and ∀ψ, φ ∈ E, for some r ∈ [0, 1] and q ∈ [1/2, 1]. Let us also define:

\[ \|u\|_V = \sqrt{\langle u, u \rangle_V} = \sqrt{\sum_x u(x)^2 d(x)^r}, \]

\[ \|\phi\|_E = \sqrt{\langle \phi, \phi \rangle_E} = \sqrt{\frac{1}{2} \sum_{x,y} \phi(x,y)^2 w(x,y)^{2q-1}}, \]

\[ \|u\|_E = \sqrt{\langle u, u \rangle_E} = \sqrt{\sum_x u(x)^2}, \]

\[ \|\phi\|_{E,\infty} = \max_{x,y} |\phi(x,y)|. \]

The gradient operator \( \nabla : V \rightarrow E \) is formulated as:

\[ (\nabla u)_w(x, y) = w(x, y) (1 - q)(u(y) - u(x)). \]

The Dirichlet energy does not include parameters \( r \) or \( q \):

\[ \frac{1}{2} \|\nabla u\|_E^2 = \frac{1}{4} \sum_{x,y} w(x,y)(u(x) - u(y))^2. \]

The divergence \( \text{div}_w : E \rightarrow V \) definition is based on the adjoint of the gradient:

\[ (\text{div}_w \phi)(x) = \frac{1}{2d(x)^r} \sum_y w(x,y)^q (\phi(x,y) - \phi(y,x)), \]

where we define the adjoint using the following definition: \( \langle \nabla u, \phi \rangle_E = -\langle u, \text{div}_w \phi \rangle_V. \)

We now have a set of graph Laplacians \( \triangle_r = \text{div}_w \nabla : V \rightarrow V \):

\[ (\triangle_w u)(x) = \sum_y w(x,y)\frac{d(x)^r}{d(x)^r} (u(y) - u(x)). \]

In matrix form, we can write the set of Laplacians as

\[ -\triangle_w = D^{1-r} - D^{-r} W. \]

The matrices resulting from \( r = 0 \) and \( r = 1 \) are called the unnormalized Laplacian and the random walk Laplacian, respectively. Using divergence, a set of anisotropic total variations \( TV_w : V \rightarrow \mathbb{R} \) can now be formulated:

\[ TV_{w,q}(u) = \max \left\{ \langle \text{div}_w \phi, u \rangle_V : \phi \in E, \|\phi\|_{E,\infty} \leq 1 \right\} = \frac{1}{2} \sum_{x,y} w(x,y)^q |u(x) - u(y)|. \]

The last operator to be defined will involve the Ginzburg-Landau functional that our methods are based upon:

\[ GL_r(u) = \|\nabla u\|_E^2 + \frac{1}{\epsilon} \sum_x W(u(x)), \]

where \( W \) is the double well potential \( W(u) = u^2(u-1)^2 \). Note that, while the first term in the continuous Ginzburg-Landau functional

\[ GL_c(u) = \epsilon \int |\nabla u|^2 dx + \frac{1}{\epsilon} \int W(u) dx. \]
is scaled by $\epsilon$, the first term of $GL_\epsilon$ does not include $\epsilon$. This is because, unlike the case with the continuous functional, the difference terms of $GL_\epsilon$ are finite even for binary functions, and no rescaling of the first term is necessary.

In relation to parameters, we choose $q = 1$ since in [37], it is shown that for any $r$ and $q = 1$, $TV_w$ is the $\Gamma$-limit (Gamma convergence) of a sequence of graph-based Ginzburg-Landau (GL)-type functionals:

**Theorem 1.** $GL_\epsilon \Gamma \rightarrow GL_0$ as $\epsilon \rightarrow 0$, where

$$GL_\epsilon(u) = \|\nabla u\|_E^2 + \frac{1}{\epsilon} \sum_x W(u(x)) = \frac{1}{2} \sum_{x,y} w(x,y)(u(x) - u(y))^2 + \frac{1}{\epsilon} \sum_x W(u(x)),$$

$$GL_0(u) = \begin{cases} TV_{w,1}(u) & \text{for } u \text{ s.t. } u(x) \in \{0,1\}, \\ \infty & \text{otherwise.} \end{cases}$$

**Proof.** See Theorem 3.1 of [37].

Thus, to be consistent with their work, use $q = 1$. The parameter $r$ is allowed to change, since we experiment with different Laplacians. With the parameters so defined, total variation is now formulated as

$$TV_w(u) = \max \{ \langle \operatorname{div}_w \phi, u \rangle_V : \phi \in \mathcal{E}, \|\phi\|_{E,\infty} \leq 1 \} = \frac{1}{2} \sum_{x,y} w(x,y)|u(x) - u(y)|. \quad (10)$$

This is the definition used in the paper.

3 Modified Ratio Cut Method

3.1 Derivation

Let us consider the binary partitioning problem involving the target set $X$ (of size $n$) embedded in a graphical framework. One way to classify the set into two classes would be to find the partition that minimizes some normalization of the cut. In this section, our derivations are based on

$$\text{RatioCut}(S, \bar{S}) = \frac{\text{cut}(S, \bar{S})}{\frac{1}{|S|} + \frac{1}{|\bar{S}|}}.$$

The ratio cut is related to the Cheeger cut $h(S, \bar{S})$ via the following inequality:

$$h(S, \bar{S}) \leq \text{RatioCut}(S, \bar{S}) \leq 2h(S, \bar{S}),$$

which can be easily shown using the definitions. We now present two different approaches to the ratio cut problem.

3.1.1 Approach I

In this section, we do not solve the exact ratio cut formulation, but instead focus on a slightly modified version. Earlier, we discussed the general problem of minimizing the cut to find a desired partition:

$$\min_{S \subseteq V} \left( \text{cut}(S, \bar{S}) = \sum_{x \in S, y \in \bar{S}} w(x, y) \right).$$
Since minimizing the cut is equivalent to minimizing its square, we can consider that slightly modified problem. However, some sort of normalization is needed to create a bias towards partitions where the two clusters are of similar size. Using the regular ratio cut normalization, one obtains the modified ratio cut problem:

$$\min_{S \subseteq V} \text{cut}(S, \bar{S})^2 \left( \frac{1}{|S|} + \frac{1}{|\bar{S}|} \right),$$

which is equivalent to

$$\min_{S \subseteq V} \text{cut}(S, \bar{S}) \sqrt{\frac{1}{|S|} + \frac{1}{|\bar{S}|}}. \quad (12)$$

Before giving more details, we note that if $\chi_S : G \to \{0, 1\}$ is the indicator function of a subset $S$ of $G$, then

$$\text{TV}_w(\chi_S) = \text{cut}(S, \bar{S}). \quad (13)$$

In addition, if the mean of a function $u : V \to \mathbb{R}$ is defined as $\frac{1}{n} \sum_{i \in V} u_i$, we also have

$$\left\| \chi_S - \text{mean}(\chi_S) \right\|^2_{L^2} = \sum_i \left( \chi_S(x_i) - \frac{|S|}{n} \right)^2 = |S| \left( 1 - \frac{|S|}{n} \right)^2 + |\bar{S}| \left( \frac{|S|}{n} \right)^2 = \frac{|S||\bar{S}|}{|S|+|\bar{S}|}, \quad (14)$$

where $x_i$ represents node $i$.

Therefore, due to (13) and (14), problem (11) (and (12)) is equivalent to the following problem:

$$\min_{u : u(x) \in \{0, 1\}} \frac{\text{TV}_w(u)}{||u - \text{mean}(u)||_{L^2}}.$$

This is an NP-hard problem, and one way to simplify it would be to relax the binary condition to include a larger set of functions; in our case, we minimize over all functions $u : V \to \mathbb{R}$ with values in $\mathbb{R}$. We now show that there exists a binary minimizer of the relaxed problem, and it represents the solution to the modified ratio cut problem.

**Theorem 2.** Let $u : V \to \mathbb{R}$. Consider the problem

$$\lambda = \min_u \frac{\text{TV}_w(u)}{||u - \text{mean}(u)||_{L^2}}. \quad (15)$$

There is binary valued minimizer, and

$$\lambda = \min_{S \subseteq V} \text{cut}(S, \bar{S}) \sqrt{\frac{1}{|S|} + \frac{1}{|\bar{S}|}}. \quad (16)$$

**Proof.** The techniques of the proof are similar to the ones in [35]. Suppose that $u^*$ is the minimizer of (15). We have $\text{TV}_w(u^*) \neq 0$, otherwise $u^*$ would be constant, and the ratio in (15) would be undefined. Note that the functional in (15) is scale invariant; we can thus rescale $u^*$ so that $\text{TV}_w(u^*)=1$. In this case, $u^*$ is a maximizer of the denominator, constrained to the set of functions $u$ such that $\text{TV}_w(u) \leq 1$. Let $A_1$ be the set of indices where $u^*$ is non-positive, and $A_2$ be the set where it is positive. Define $Z$ to be the set of functions that are non-positive on the first set and non-negative on the second one, and $C$ to be the TV norm unit ball on $Z$. Clearly, $u^*$ is a solution to $\max_C ||u - \text{mean}(u)||_{L^2}$. Since the set $C$ is convex, and $E(u) = ||u - \text{mean}(u)||_{L^2}$ is a convex functional on $C$, a maximum is also attained on at least one of the extreme points. By Lemma 2.2 of [35], any extreme point of $C$ takes the
form of a (scaled) indicator function of some set $S^\ast$. Therefore, there exists an extreme point $v = \alpha\chi_{S^\ast}$ of $C$, which is also a solution to $\max_C \|u - \text{mean}(u)\|_2$. Note that due to the binary property of $v$, and since $\|u^\ast - \text{mean}(u^\ast)\|_2 = \|v - \text{mean}(v)\|_2$ and $TV_w(v) \leq TV_w(u^\ast) = 1$,

$$\min_u TV_w(u) \|u - \text{mean}(u)\|_2^2 = \lambda = \frac{TV_w(u^\ast)}{\|u^\ast - \text{mean}(u^\ast)\|_2} \geq \frac{TV_w(v)}{\|v - \text{mean}(v)\|_2} = \text{cut}(S^\ast, S\bar{\ast})\sqrt{\frac{1}{|S^\ast|} + \frac{1}{|S\bar{\ast}|}} \geq \min_{S \subset V} \text{cut}(S, S\bar{\ast})\sqrt{\frac{1}{|S|} + \frac{1}{|S\bar{\ast}|}} \geq \min_u TV_w(u) \|u - \text{mean}(u)\|_2,$$

Since

$$\min_{S \subset V} \text{cut}(S, S\bar{\ast})\sqrt{\frac{1}{|S|} + \frac{1}{|S\bar{\ast}|}} = \min_u TV_w(u) \|u - \text{mean}(u)\|_2 \geq \min_u TV_w(u) \|u - \text{mean}(u)\|_2,$$

we get equality in (17).

Therefore, we consider the problem

$$\min_u \frac{TV_w(u)}{\|u - \text{mean}(u)\|_2^2}. \tag{18}$$

Due to the relationship shown in Theorem 1, the idea now is to interchange total variation, which is directly related to the cut, with the Ginzburg-Landau functional. This will allow for “cleaner” minimization techniques due to the $L^2$ term of the functional, and will later allow us to express all the calculations via the Laplacian, for which we use a spectral approach. We later introduce the use of the Nyström extension method [15, 16] - a clever technique that approximates the eigenvectors of a Laplacian using only a small portion of the graph weights. This technique is very efficient, and especially useful for very big data sets. Replacing the total variation in (18) by the Ginzburg-Landau functional, we obtain the problem of

$$\min_u F(u) = \min_u \frac{GL_w(u)}{\|u - \text{mean}(u)\|_2^2}. \tag{19}$$

We denote the function by which $GL_w(u)$ is multiplied by $C_1(u)$.

To solve (19), we compute the modified Allen-Cahn equation using the $L^2$-gradient flow associated with (19). Since

$$\frac{d}{dt} F(u + tv) = (\frac{d}{dt} GL_w(u + tv))(C_1(u + tv)) + GL_w(u + tv)(\frac{d}{dt} C_1(u + tv)),$$

evaluating the derivative at zero gives

$$\frac{d}{dt} F(u + tv)|_{t=0} = -2(C_1(u)\Delta_w u, v)_{L^2} + \frac{1}{\varepsilon}(C_1(u)W'(u), v)_{L^2} + (C_2(u)GL_w(u), v)_{L^2},$$

where

$$C_1(u) = \frac{1}{\|u - \text{mean}(u)\|_2^2} \tag{20}$$

and

$$C_2(u) = -\frac{u - \text{mean}(u)}{\|u - \text{mean}(u)\|_2^3}. \tag{21}$$

We obtain the modified Allen-Cahn equation:

$$\frac{\partial u(x)}{\partial t} = 2C_1(u)(\Delta_w u)(x) - \frac{C_1(u)W'(u(x))}{\varepsilon} - C_2(u)GL_w(u),$$
or equivalently
\[
\frac{\partial u(x)}{\partial t} = 2C_1(u) \Delta_w u(x) - \frac{C_1(u) W'(u(x))}{\epsilon} - C_2(u) ||\nabla u||_2^2 - \frac{C_2(u) \sum_i W(u(x))}{\epsilon}.
\] (22)

We will describe how we solve it shortly.

3.1.2 Approach II

In this section, we consider the exact ratio cut problem. Let \( u \) be a binary function (taking values 0 or 1) denoting the class of each of the nodes. Then one can write the ratio cut problem as the task
\[
\min_{u: u(x) \in \{0, 1\}} TV_w(u) \left( \frac{1}{\sum_x u(x)} + \frac{1}{n - \sum_x u(x)} \right).
\] (23)

Similarly to the procedure of the previous section, we relax the binary constraint and replace the total variation term in (23) by the GL functional, with motivation described in Section 3.1.1, so the following problem is formulated:
\[
\min_u F(u) = GL_\epsilon(u) \left( \frac{1}{\sum_x u(x)} + \frac{1}{n - \sum_x u(x)} \right).
\] (24)

Let us denote the second term in the product by \( C_1(u) \).

To solve (24), we compute the modified Allen-Cahn equation using the \( L^2 \)-gradient flow associated with (24).

Since
\[
\frac{d}{dt} F(u + tv) = (\frac{d}{dt} GL_\epsilon(u + tv))(C_1(u + tv)) + GL_\epsilon(u + tv) (\frac{d}{dt} C_1(u + tv)),
\]
evaluating the derivative at zero gives
\[
\frac{d}{dt} F(u + tv) \bigg|_{t=0} = -2\langle C_1(u) \Delta_w u, v \rangle_{L^2} + \frac{1}{\epsilon} \langle C_1(u) W'(u), v \rangle_{L^2} + \langle C_2(u) GL(u), v \rangle_{L^2},
\]
where
\[
C_1(u) = \frac{1}{\sum_x u(x)} + \frac{1}{n - \sum_x u(x)}.
\] (25)

and
\[
C_2(u) = \frac{-1}{(\sum_x u(x))^2} + \frac{1}{(n - \sum_x u(x))^2}.
\] (26)

Again, we obtain the modified Allen-Cahn equation
\[
\frac{\partial u(x)}{\partial t} = 2C_1(u) \Delta_w u(x) - \frac{C_1(u) W'(u(x))}{\epsilon} - C_2(u) ||\nabla u||_2^2 - \frac{C_2(u) \sum_i W(u(x))}{\epsilon}.
\] (27)

3.2 The Procedure of Modified Ratio Cut Method (MRCM)

We now present two different versions of the algorithm solving (22) (or (27)). The first version computes the gradient term of the Ginzburg-Landau functional explicitly using the graph weights via equation (4). The second version uses
\[
\|\nabla u\|_2^2 = -(\Delta u, u)_V
\] (28)
and spectral composition of the Laplacian to compute the Dirichlet's energy. Since the Laplacian is the only other operator in (27), one can use the Nyström extension method to efficiently approximate the eigenvectors that will then
be used in (28) to compute the Dirichlet’s energy. Only a small portion of the graph weights have to be computed in the method’s procedure, which results in more efficiency, and this is the advantage of the second version of the algorithm. In fact, for most data sets, we only needed to compute a very small fraction of the weights, around 0.3 percent.

3.2.1 Version 1

Discretizing (22) and calculating the Laplacian implicitly, one obtains
\[
\frac{u_i^{n+1} - u_i^n}{dt} = 2C_1(u^n)(\Delta_w u_i^{n+1})_i - \frac{C_1(u^n)W'(u^n)}{\epsilon} - C_2(u^n)||\nabla_w u^n||^2_{\mathcal{E}} - \frac{C_2(u^n)\sum \mathcal{W}(u^n)}{\epsilon}.
\] (29)

To solve the equation, we use the eigendecomposition of a Laplacian and write terms in the above equation as series involving eigenvectors. Let
\[
u^n = \sum_k a_k^n \phi_k(x),
\] (30)
\[
-C_2(u^n)||\nabla_w u^n||^2_{\mathcal{E}} - \frac{C_1(u^n)W'(u^n)}{\epsilon} - \frac{\sum \mathcal{W}(u^n)C_2(u^n)}{\epsilon} = \sum_k b_k^n \phi_k(x),
\] (31)
where \(\{\phi_k(x)\}_k\) are the eigenfunctions of a Laplacian. Here, we calculate the Dirichlet’s energy term explicitly using (4). After plugging (30) and (31) into (29), one obtains
\[
a_k^{n+1} = \frac{a_k^n + dtb_k^n}{1 + 2dtC_1(u^n)\lambda_k},
\] (32)
where \(\lambda_k\) is the \(k^{th}\) eigenvalue of a Laplacian.

3.2.2 Version 2

Discretizing (22) and calculating the Laplacian implicitly, one obtains
\[
\frac{u_i^{n+1} - u_i^n}{dt} = 2C_1(u^n)(\Delta_w u_i^{n+1})_i - \frac{C_1(u^n)W'(u^n)}{\epsilon} - C_2(u^n)||\nabla_w u^n||^2_{\mathcal{E}} - \frac{C_2(u^n)\sum \mathcal{W}(u^n)}{\epsilon}.
\] (33)

To solve the above equation, we use the eigendecomposition of a Laplacian. Let \(u^n = \sum_k a_k^n \phi_k(x)\), where \(\{\phi_k(x)\}_k\) are the eigenfunctions of a Laplacian. Using (28) and the orthonormal property of eigenvectors, one can derive the following equation involving the \(||\nabla_w u^n||^2_{\mathcal{E}}\) term:
\[
||\nabla_w u^n||^2_{\mathcal{E}} = \sum_k \lambda_k (a_k^n)^2,
\] (34)
where \(\lambda_k\) are the eigenvalues of a Laplacian.

Therefore, when calculating
\[
-C_2(u^n)||\nabla_w u^n||^2_{\mathcal{E}} - \frac{C_1(u^n)W'(u^n)}{\epsilon} - \frac{C_2(u^n)\sum \mathcal{W}(u^n)}{\epsilon}
\]
to find \(b^n\) (see (31)), we use (34). Thus, in this version, the coefficients \(b_k^n\) are defined as:
\[
-C_2(u^n)\sum_k \lambda_k (a_k^n)^2 - \frac{C_1(u^n)W'(u^n)}{\epsilon} - \frac{C_2(u^n)\sum \mathcal{W}(u^n)}{\epsilon} = \sum_k b_k^n \phi_k(x),
\] (35)
where we have used (34) to replace the Dirichlet’s energy term. After plugging these representations into (33), one obtains
\[ a_{k}^{n+1} = \frac{a_{k}^{n} + dtb_{k}^{n}}{1 + 2dtC_{1}(u^{n})\lambda_{k}}, \]
where \( \lambda_{k} \) is the \( k^{th} \) eigenvalue of a Laplacian. Below, we outline the steps of both versions of the method:

### Modified Ratio Cut Method (MRCM)

Here \( u \) will be a function whose thresholded value indicates the class membership. The steps of the algorithm are:

* Organize the data set in a graphical setting, compute eigenvalues and eigenvectors of a chosen Laplacian.
* For Approach I, calculate \( C_{1} \) and \( C_{2} \) via (20) and (21), respectively. For Approach II, calculate \( C_{1} \) and \( C_{2} \) via (25) and (26), respectively.
* Initialize \( u^{0} \), calculate \( a^{0} \) using (30), calculate \( b^{0} \) using (31) (for Version 1) and (35) (for Version 2).
* Repeat the following steps from \( n = 0 \) until a stopping criterion is satisfied:
  - For Version 1, calculate \( a^{n+1} \) using (32). For Version 2, calculate \( a^{n+1} \) using (36).
  - For Version 1, calculate \( b^{n+1} \) using (31). For Version 2, calculate \( b^{n+1} \) using (35)
  - Calculate \( u^{n+1} \) using (30).
* Threshold the solution to obtain a binary answer, which represents the assigned class.

### 4 Modified Cheeger Cut Method

#### 4.1 Derivation

We again consider the binary partitioning problem involving the target set \( X \) embedded in a graphical framework. We use similar techniques of the previous sections, but using the Cheeger cut. As described in Section 1, the Cheeger cut problem can be formulated as finding the subset \( S \) of \( X \) such that the following value is minimized:

\[ h(S, \bar{S}) = \frac{\text{cut}(S, \bar{S})}{\min(|S|, |\bar{S}|)} \]

If \( u \) is a binary function (taking values 0 or 1) denoting the class of each of the nodes, one can write the Cheeger problem as the task

\[ \min_{\{u: u(x) \in \{0,1\}\}} TV_{w}(u) \left( \frac{1}{\min\left( \sum_{x} u(x), n - \sum_{x} u(x) \right)} \right). \]

Unfortunately, this is an NP hard problem, and one approach is to relaxation of constraints. We relax the problem by minimizing over functions with values in \( \mathbb{R} \):

\[ \min_{u} TV_{w}(u) \left( \frac{1}{\min\left( \sum_{x} u(x), n - \sum_{x} u(x) \right)} \right). \]  

(37)
Similarly to the procedure of the previous section, we replace the total variation term in (23) by the GL functional, with motivation described in Section 3.1.1, so that the above problem is formulated as:

$$\min_u F(u) = GL_\epsilon(u) \left( \frac{1}{\min(\sum_x u(x), n - \sum_x u(x))} \right).$$

(38)

We denote the second term in the product by $B_1(u)$.

To solve (38), we compute the modified Allen-Cahn equation using the $L^2$-gradient flow associated with GL functional. Since

$$\frac{d}{dt} F(u + tv) = \left( \frac{d}{dt} GL_\epsilon(u + tv) \right)(B_1(u + tv)) + GL_\epsilon(u + tv)(\frac{d}{dt} B_1(u + tv)),$$

evaluating the derivative at zero gives

$$\frac{d}{dt} F(u + tv)|_{t=0} = -2\langle B_1(u) \Delta u, v \rangle_{L^2} + \frac{1}{\epsilon} \langle B_1(u) W'(u), v \rangle_{L^2} + \langle B_2(u) GL_\epsilon(u), v \rangle_{L^2},$$

where

$$B_1(u) = \frac{1}{\min(\sum_x u(x), n - \sum_x u(x))}$$

and

$$B_2(u) = \begin{cases} \frac{1}{(\sum_x u(x))^2} & \text{if } (\sum_x u(x))^2 \geq (n - \sum_x u(x))^2 \\ \frac{1}{(\sum_x u(x))^2} & \text{if } (\sum_x u(x))^2 < (n - \sum_x u(x))^2 \end{cases}$$

The modified Allen-Cahn equation is thus:

$$\frac{\partial u(x)}{\partial t} = 2B_1(u)(\Delta u)(x) - \frac{B_1(u) W'(u(x))}{\epsilon} - B_2(u) GL_\epsilon(u)$$

or equivalently

$$\frac{\partial u(x)}{\partial t} = 2B_1(u)(\Delta u)(x) - \frac{B_1(u) W'(u(x))}{\epsilon} - B_2(u) ||\nabla u||_E^2 - \frac{(B_2(u) \sum_i W(u(x)))}{\epsilon}.$$

(39)

In the next section, we describe how we solve this equation.

### 4.2 The Procedure of the Modified Cheeger Cut Method

We present two different versions of the modified Cheeger cut method. As mentioned in the previous section, Version 2 calculates the Dirichlet’s energy term of the Ginzburg-Landau functional using a different approach than the explicit graph formula (4). This allows one to use the Nyström extension method, due to the Laplace operator, and thus only a small portion of the graph weights need to be calculated. Thus, this version is more efficient.

#### 4.2.1 Version 1

Discretizing (39) and calculating the Laplacian implicitly, one obtains

$$\frac{u_{i}^{n+1} - u_i^n}{dt} = 2B_1(u^n)(\Delta u^{n+1})_i - B_2(u^n)||\nabla u^n||_E^2 - \frac{B_1(u^n) W'(u^n)}{\epsilon} - \frac{B_2(u^n) \sum_i W(u^n)}{\epsilon}.$$

(40)
To solve the equation, we use

\[ u^{n} = \sum_{k} a_{k}^{n} \phi_{k}(x), \]  

\[ -B_{2}(u^{n}) \| \nabla_{w} u^{n} \|^2_{L^{2}} - \frac{B_{1}(u^{n}) W'(u^{n})}{\epsilon} - \frac{B_{2}(u^{n}) \sum_{i} W(u_{i}^{n})}{\epsilon} = \sum_{k} b_{k}^{n} \phi_{k}(x), \]  

where \{ \phi_{k}(x) \}_{k} are the eigenfunctions of a Laplacian. The Dirichlet’s energy term is calculated explicitly via (4).

After plugging (41) and (42) into (40), one obtains

\[ a_{k}^{n+1} = \frac{a_{k}^{n} + db_{k}^{n}}{1 + 2dtB_{1}(u^{n})\lambda_{k}}, \]  

where \lambda_{k} is the \( k^{th} \) eigenvalue of a Laplacian.

### 4.2.2 Version 2

Discretizing (39) and calculating the Laplacian implicitly, one obtains

\[ \frac{u_{i}^{n+1} - u_{i}^{n}}{dt} = 2B_{1}(u^{n})(\Delta_{w} u^{n+1})_{i} - B_{2}(u^{n}) \| \nabla_{w} u^{n} \|^2_{L^{2}} - \frac{B_{1}(u^{n}) W'(u^{n})}{\epsilon} - \frac{B_{2}(u^{n}) \sum_{i} W(u_{i}^{n})}{\epsilon}. \]  

To solve the above equation, we use the eigendecomposition of a Laplacian. Let \( u^{n} = \sum_{k} a_{k}^{n} \phi_{k}(x), \) where \{ \phi_{k}(x) \}_{k} are the eigenfunctions of a Laplacian. When calculating

\[ -B_{2}(u^{n}) \| \nabla_{w} u^{n} \|^2_{L^{2}} - \frac{B_{1}(u^{n}) W'(u^{n})}{\epsilon} - \frac{B_{2}(u^{n}) \sum_{i} W(u_{i}^{n})}{\epsilon} = \sum_{k} b_{k}^{n} \phi_{k}(x), \]  

to find \( b^{n} \) (see (42)), we use (34). Thus, in this version, the coefficients \( b_{k}^{n} \) are defined as:

\[ -B_{2}(u^{n}) \sum_{k} \lambda_{k}(a_{k}^{n})^{2} - \frac{B_{1}(u^{n}) W'(u^{n})}{\epsilon} - \frac{B_{2}(u^{n}) \sum_{i} W(u_{i}^{n})}{\epsilon} = \sum_{k} b_{k}^{n} \phi_{k}(x), \]  

where we have used (34) to replace the Dirichlet’s energy term. Here, \( \lambda_{k} \) is the \( k^{th} \) eigenvalue of a Laplacian. After plugging these representations into (44), one obtains

\[ a_{k}^{n+1} = \frac{a_{k}^{n} + dtb_{k}^{n}}{1 + 2dtB_{1}(u^{n})\lambda_{k}}. \]  

Below, we outline steps of both versions of the algorithm:

---

**Modified Cheeger Cut Method (MCCM) - Versions 1 and 2**

Here \( u \) will be a function whose thresholded value indicates the class membership. The steps of the algorithm are:

* Organize the data set in a graphical setting, compute eigenvalues and eigenvectors of a chosen Laplacian.

* Initialize \( u^{0} \), calculate \( a^{0} \) using (30), calculate \( b^{0} \) using (42) (for Version 1) and (45) (for Version 2).

* Repeat the following steps from \( n = 0 \) until a stopping criterion is satisfied:
  
  - For Version 1, calculate \( a^{n+1} \) using (43). For Version 2, calculate \( a^{n+1} \) using (46).
  
  - For Version 1, calculate \( b^{n+1} \) using (42). For Version 2, calculate \( b^{n+1} \) using (45).
  
  - Calculate \( u^{n+1} \) using (30).

* Threshold the solution to obtain a binary answer, which represents the assigned class.
5 Parametrized Laplacian Framework

In [18], Ghosh et al. formulated a general Laplacian framework by considering a dynamic process on graphs:

\[
\frac{d\Theta}{dt} = -L\Theta.
\]

Here, \(\Theta\) is a column vector containing the values of the dynamic variable for all vertices, and \(L\) is a positive semi-definite matrix called the spreading operator, which defines the dynamic process. The parameterized Laplacian family is defined as

\[
L = T^{\frac{1}{2}}D^{\frac{1}{2}}(D' - BWB)D^{\frac{1}{2}}T^{-\frac{1}{2}}.
\]

Compared with traditional Laplacian operators, the parametrized Laplacian has two additional parameters: \(T\) and \(B\). The diagonal matrix \(T\) controls the time delay factors, or local clock rate of a random walk, at each vertex. The bias factors form the other diagonal matrix \(B\). It changes the trajectory by giving random walk targets different weights. Note that the degree matrix \(D'\) is defined in terms of the biased weight matrix: \(d'_i = \sum_j [W']_{ij} = \sum_j [BWB]_{ij}\), where we set \(W' = BWB\).

To investigate how different dynamic processes effect our perceptions of the graph structure, we study how the results change when \(T\) and \(B\) vary. Recall in Section 2.2, we defined a similar family of Laplacians with the parameter \(r\) (6). As we will show for the following four special cases, they can be formulated as equivalent operators up to a similarity transformation.

* **Normalized Laplacian.** In the case when \(T = B = I\), we obtain the symmetric normalized Laplacian:

\[
L_n = D^{-\frac{1}{2}}(D - W)D^{-\frac{1}{2}}.
\]

The normalized Laplacian is related to the random walk Laplacian \((r = 1)\) by a change of basis: \(L_n = D^{\frac{1}{2}}L_{rw}D^{-\frac{1}{2}}\).

* **Scaled Graph Laplacian.** In the case when \(B = I\) and \(T = d_{max}D^{-1}\) (where \(d_{max}\) represents the maximum degree in the degree matrix), the spreading operator is called the scaled graph Laplacian:

\[
L_{sc1} = \frac{1}{d_{max}}(D - W).
\]

It is just the matrix (6) with \(r = 0\) scaled by \(\frac{1}{d_{max}}\).

* **Unbiased Adjacency Matrix.** If we let \(B = D^{-\frac{1}{2}}\) and \(T = d'_{max}D'^{-1}\) (where \(d'_{max}\) represents the maximum degree in the degree matrix associated with \(W' = D^{-\frac{1}{2}}WD^{-\frac{1}{2}}\)), we obtain the unbiased adjacency matrix:

\[
L_{unb} = \frac{1}{d'_{max}}(D' - D^{-\frac{1}{2}}WD^{-\frac{1}{2}}) = \frac{1}{d'_{max}}(D' - W').
\]

The unbiased adjacency matrix is related to (6) with \(r = 0\) (using a similarity matrix biased by \(D^{-\frac{1}{2}}\)) by a factor of \(\frac{1}{d_{max}}\).
* **Replicator.** If we let $B = V$, where $V$ is a diagonal matrix whose elements are the components of the eigenvector corresponding to the largest eigenvalue of $W$, and $T = I$, we obtain

$$L_{\text{rep}} = I - \frac{1}{\lambda_{\text{max}}} W,$$

where $\lambda_{\text{max}}$ is the largest eigenvalue of $W$. The normalized Laplacian is related to the max entropy random walk by a change of basis: $L_{\text{rep}} = D^{\frac{1}{2}} D^{-1} (D' - V W V) D'^{-\frac{1}{2}}$, where $D'^{-\frac{1}{2}} (D' - V W V)$ is the max entropy random walk.

We utilize these four different formulations of the Laplacian in our paper, and the results are described in the next section.

## 6 Results

In this section, we show results for some benchmark data sets, and then show an application to hyperspectral imagery. We observe that the advantage of these methods is efficiency, while being able to maintain the accuracy of state-of-the-art. Moreover, we note that all the algorithms presented in this paper are alike in terms of efficiency.

The energy minimization proceeds until a steady state condition is reached. Once the change of the norm of the vector field in subsequent iterations falls below a threshold, the system is no longer evolving and the energy decrement is negligible. In the experiments, obtained on a 2.4 GHz Intel Core i2 Quad computer, the calculation is stopped when

$$\|u^{n+1} - u^n\|_2^2 < \eta,$$

where the value of $\eta$ is dependent on the data set.

When choosing a weight function, the goal is to assign a large weight to an edge if the two vertices it is connecting are similar and a small weight otherwise. One popular choice for the weight function is the Gaussian

$$w(x, y) = e^{-\frac{M(x, y)^2}{\sigma^2}},$$

where $M(x, y)$ is some distance measure between the two vertices $x$ and $y$, and $\sigma$ is a parameter to be chosen. For example, if the data set consists of points in $\mathbb{R}^2$, $M(x, y)$ can be chosen to be the Euclidean distance between $x$ and $y$, since points farther away are less likely to belong to the same cluster than points closer together. For images, $M(x, y)$ can be defined as the weighted $L^2$-norm of the difference of the feature vectors of pixels $x$ and $y$, where the feature vector of a pixel can be defined as the set of intensity values in the pixel’s neighborhood, as described in [20].

Another choice for the weight function is the Zelnik-Manor and Perona function [31] for sparse matrices:

$$w(x, y) = e^{-\frac{M(x, y)^2}{\tau(x) \tau(y)}},$$

using $\tau(x) = M(x, z)^2$, where $z$ is the $K^{th}$ closest vertex to vertex $x$. The parameter $K$ is to be chosen.

Note that it is not necessary or even desirable to use a fully connected graph setting, which might be a computational burden. Specifically, the fully connected graph can be approximated by a much smaller graph by only including
Table 1: Modified Cheeger and ratio cut method results and comparison

<table>
<thead>
<tr>
<th>Method</th>
<th>Accuracy</th>
</tr>
</thead>
<tbody>
<tr>
<td>symmetric Laplacian (MRCM- Approach 1)</td>
<td>98.62%</td>
</tr>
<tr>
<td>symmetric Laplacian (MRCM- Approach 2)</td>
<td>98.27%</td>
</tr>
<tr>
<td>symmetric Laplacian (MCCM)</td>
<td>98.33%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 1)</td>
<td>98.45%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 2)</td>
<td>98.12%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MCCM)</td>
<td>98.2%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 1)</td>
<td>98.66%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 2)</td>
<td>98.26%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MCCM)</td>
<td>98.32%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 1)</td>
<td>96.60%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 2)</td>
<td>94.38%</td>
</tr>
<tr>
<td>replicator (MCCM)</td>
<td>94.14%</td>
</tr>
<tr>
<td>method in [3]</td>
<td>98.36%</td>
</tr>
<tr>
<td>Timing for our methods</td>
<td>2 to 15 sec.</td>
</tr>
</tbody>
</table>

Timing for our methods

<table>
<thead>
<tr>
<th>Method</th>
<th>Accuracy</th>
</tr>
</thead>
<tbody>
<tr>
<td>symmetric Laplacian (MRCM- Approach 1)</td>
<td>98.35%</td>
</tr>
<tr>
<td>symmetric Laplacian (MRCM- Approach 2)</td>
<td>97.80%</td>
</tr>
<tr>
<td>symmetric Laplacian (MCCM)</td>
<td>98.00%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 1)</td>
<td>98.35%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 2)</td>
<td>97.45%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MCCM)</td>
<td>97.70%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 1)</td>
<td>98.05%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 2)</td>
<td>97.55%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MCCM)</td>
<td>97.70%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 1)</td>
<td>97.35%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 2)</td>
<td>95.50%</td>
</tr>
<tr>
<td>replicator (MCCM)</td>
<td>95.55%</td>
</tr>
<tr>
<td>method in [35]</td>
<td>95.08%</td>
</tr>
<tr>
<td>method in [8]</td>
<td>93.5%</td>
</tr>
<tr>
<td>method in [23]</td>
<td>95.38%</td>
</tr>
<tr>
<td>method in [5]</td>
<td>91.31%</td>
</tr>
<tr>
<td>method in [3]</td>
<td>95.86%</td>
</tr>
</tbody>
</table>

Two moons

<table>
<thead>
<tr>
<th>Method</th>
<th>Accuracy</th>
</tr>
</thead>
<tbody>
<tr>
<td>symmetric Laplacian (MRCM- Approach 1)</td>
<td>98.35%</td>
</tr>
<tr>
<td>symmetric Laplacian (MRCM- Approach 2)</td>
<td>97.80%</td>
</tr>
<tr>
<td>symmetric Laplacian (MCCM)</td>
<td>98.00%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 1)</td>
<td>98.35%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MRCM- Approach 2)</td>
<td>97.45%</td>
</tr>
<tr>
<td>scaled graph Laplacian (MCCM)</td>
<td>97.70%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 1)</td>
<td>98.05%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MRCM- Approach 2)</td>
<td>97.55%</td>
</tr>
<tr>
<td>unbiased adjacency matrix (MCCM)</td>
<td>97.70%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 1)</td>
<td>97.35%</td>
</tr>
<tr>
<td>replicator (MRCM- Approach 2)</td>
<td>95.50%</td>
</tr>
<tr>
<td>replicator (MCCM)</td>
<td>95.55%</td>
</tr>
<tr>
<td>method in [35]</td>
<td>95.08%</td>
</tr>
<tr>
<td>method in [8]</td>
<td>93.5%</td>
</tr>
<tr>
<td>method in [23]</td>
<td>95.38%</td>
</tr>
<tr>
<td>method in [5]</td>
<td>91.31%</td>
</tr>
<tr>
<td>method in [3]</td>
<td>95.86%</td>
</tr>
</tbody>
</table>

an edge between vertex $x$ and $y$ if $x$ is among the $k$-nearest neighbors of $y$ or vice versa. This is called a $k$-nearest neighbor graph. We sparsify the graph in this way. One can also build a mutual $k$-nearest neighbor graph by only including an edge between $x$ and $y$ if both of them are $k$-nearest neighbors of each other. If two vertices $x$ and $y$ are not connected by an edge, the weight between them is set to 0.

6.1 MNIST Data Set

The MNIST digits data set [26], available at http://yann.lecun.com/exdb/mnist/, is a data set of 70000 $28 \times 28$ images of handwritten digits from 0 – 9 to be classified into ten classes. Since our method is only binary, we obtained a subset of this set to cluster into two classes; in particular, we chose digits 4 and 9 since these digits are sometimes hard to distinguish, if handwritten. This created a set of 13782 digits, each either 4 or 9. We start with a random initialization of the class, and the goal is to classify each image into either a 4 or 9.

The MNIST data set results for both versions of the MCRM (both approaches) and MCCM methods are shown in Table 1. Our methods are italicized, and the best method is written in bold. From the table, we see that all but the replicator achieve accuracy in the 98th percentile. Using the replicator worsens the accuracy by at least one percentile. The results of MRCM (Approach 1) are visualized in Figure 1; those of other algorithms are not shown simply because they are very similar. To compare to some state-of-the-art work involving the cut with normalization (i.e. Cheeger cut,
balanced cut, etc.) or spectral computations, we note the result of 98.36% of Bresson et al. in [3]. We see that our methods achieve accuracy that is very similar to that result, in addition to being very efficient computationally. In fact, after the graph was computed, for all algorithms presented in this paper, the minimization took only around 2 to 15 seconds, depending on the number of iterations. The mean time for the computations recorded in [3] is 92.2 seconds.

![Image of MNIST data set results](image)

(a) ground truth  (b) symmetric graph Laplacian  (c) scaled graph Laplacian

(d) unbiased adjacency matrix  (e) replicator

Figure 1: MNIST data set results

### 6.2 Two Moons Data Set

The data set is constructed from two half circles in $\mathbb{R}^2$ with a radius of one with centers at $(0, 0)$ and $(1, 0.5)$. A thousand uniformly chosen points are sampled from each circle, embedded in $\mathbb{R}^{100}$ and i.i.d. Gaussian noise with standard deviation 0.02 is added to each coordinate, giving a set of two thousand points. Starting from some initial classification of the points, the goal is to segment the two half circles. Here, we use an initialization involving the thresholded second eigenvector of the chosen Laplacian.

The results for both versions of the MRCM (both approaches) and MCCM algorithms are shown in Table 1. The table shows that all but the replicator always achieve accuracies in the 97th or 98th percentile. The results of MRCM (Approach 1) are visualized in Figure 2; those of the rest are not shown because they are very similar.

To compare to some state-of-the-art work involving the cut with normalization (i.e. Cheeger cut, balanced cut, etc.) or spectral computations, we note the 95.08% result of Szlam et al. in [35], the 93.5% result of Buhler et al. in [8], the 95.38% result of Hein et al. in [23], the 91.31% result of Bresson et al. in [5] and the 95.86% result of Bresson et al. in [3]. We see that our method achieves an accuracy that is at least 2% higher.
6.3 Application to Hyperspectral Imagery

We apply our method to hyperspectral data; the goal is to track chemical plumes recorded as a hyperspectral video sequence. The data [7], provided by the Applied Physics Laboratory at John Hopkins University, is a video sequence of plumes released at the Dugway Proving Ground. The images are of dimension 128 × 320 × 129, where the last dimension indicates the number of channels, each depicting a particular frequency from 7,820 nm to 11,700 nm, spaced 30 nm apart. The sets of images were taken from videos captured by three long wave infrared (LWIR) spectrometers, each placed at a different location about 2 km away from the release of plume at an elevation of 1300 feet. One hyperspectral image is captured every five seconds. Other work on this data set can be found in [17], [36] and [34]. Prior work on hyperspectral plume detection using other sensors includes [24] (MWIR) and [28] (HYDICE).

Since our data depends on time, to form a graph on the set, we select \( k \) different video frames and then concatenate the points, allowing for the data to be associated over these \( k \) frames. Although the resulting graph is very large (for 7 frames, the full graph Laplacian is an \( n \times n \) matrix with \( n = 286,720 \)), we use the Nyström extension method to effectively calculate the desired eigenfunctions. The algorithm allows one to calculate an approximation of the eigenfunctions, while only having to compute a small portion of the graph. In this case of a very large graph, we need to use Version 2 of the proposed algorithm, because the goal is to not calculate the Dirichlet’s energy terms using the explicit formula (4), which would be computationally inefficient, but use the eigenfunctions of the Laplacian and (34).

Figure 3 shows a sampling of four different eigenvectors for a subset of 7 frames. Note that each eigenvector highlights a different aspect of the image. To produce a good initialization, we use an operator assisted method involving spectral clustering. In particular, by thresholding appropriately the values of eigenvectors, one obtains information about a particular class. For example, as shown in Figure 3, the third eigenvector highlights the plume.
By thresholding its values, one can find the pixels that are most likely part of the plume.

We tested our method on several video frames. The initialization for 7 of the frames is displayed in Figure 4. The final segmentation results for these frames, after 141 iterations, are shown in Figure 5. The plume is outlined in dark red, and the background shown in light green. This experiment took around 23 seconds after the graph was computed.

7 Conclusion

This paper introduces new methods that utilize the graph Laplacian to solve the fully nonlinear Cheeger and (modified) ratio cut problems, with all algorithms making use of the Ginzburg-Landau functional. To make spectral and graph computations efficient, fast numerical routines are used. The experiments, which include an application to hyperspectral data, show that the algorithms produce results that are comparable with or better than some of the best methods, and very efficiently. However, there is more work to be done; we plan on extending the two procedures to the multiclass case, and finding other interesting applications.
Figure 4: Plume data set initialization

Figure 5: Plume data set results

Acknowledgment

The authors would like to thank Chris Anderson for providing the code for the Rayleigh-Chebyshev procedure of [1], and Fan Chung for her useful suggestions. This work was supported by ONR grant N000141210838, AFOSR MURI grant FA9550-10-1-0569, NSF grants DMS-1417674, DMS-1118971 and CIF-1217605, UC Lab Fees Research grant 12-LR-236660 and DARPA W911NF-12-1-0034. Ekaterina Merkurjev is also supported by the UC President’s Postdoctoral Fellowship.

Appendix

Extensions and Variants

One can derive similar algorithms using volume of a set instead of its size, where volume is defined as \( \text{vol}(S) = \sum_{x \in S} d(x) \), where \( d(x) \) is the degree of vertex \( x \). In Approach I, we can propose a slightly different problem of

\[
\min_{S \subseteq V} \text{cut}(S, \bar{S}) \sqrt{\frac{1}{\text{vol}(S)} + \frac{1}{\text{vol}(\bar{S})}}.
\]
Let us redefine the mean to be \( \text{mean}(u) = \frac{1}{\sum_{x} d(x)} \left( \sum_{x} d(x) u(x) \right) \), and let \( \ell^2 \) be the weighted \( L^2 \) norm with the degrees as the weights. It can be easily shown that

\[
\frac{1}{\| \chi_S - \text{mean}(\chi_S) \|_{\ell^2}} = \sqrt{\frac{1}{\text{vol}(S)} + \frac{1}{\text{vol}(\bar{S})}}.
\]

The rest of the derivation proceeds as before, and we obtain

\[
C_1(u) = \frac{1}{\| u - \text{mean}(u) \|_{\ell}^3}, \quad C_2(u)(x) = -\frac{d(x)(u - \text{mean}(u))}{\| u - \text{mean}(u) \|_{\ell}^3}.
\]

In Approach II, we can propose a slightly different problem of

\[
\min_{S \subseteq V} \text{cut}(S, \bar{S}) \left( \frac{1}{\text{vol}(S)} + \frac{1}{\text{vol}(\bar{S})} \right).
\]

Now, \( C_1 \) and \( C_2 \) would be the following:

\[
C_1(u) = \frac{2}{\sum_{x} d(x) u(x)} + \frac{2}{\sum_{x} d(x)(1 - u(x))},
\]

\[
C_2(u) = \frac{-2}{(\sum_{x} d(x) u(x))^2} + \frac{2}{\sum_{x} d(x)(1 - u(x))^2}.
\]

We have derived the similar algorithms, but the results of the two slightly different formulations are very similar to those of the original versions, and thus we only state the results for the original formulation.

**Calculation of the Eigenvalues/Eigenvalues of the Graph Laplacian**

Our method involves the computation of eigenvalues and associated eigenvectors of graph Laplacians. In practice, one needs to compute only a fraction of the eigenvalues and eigenvectors because only a small portion of the nodes are significant. When the graph is sparse and is of moderate size, around \( 5000 \times 5000 \) or less, we use a Rayleigh-Chebyshev procedure outlined in [1]. It is a modification of an inverse subspace iteration method that uses adaptively determined Chebyshev polynomials. When the graph is very large, the Nyström extension method [15, 16] can be used. It is a matrix completion method, and is often used in many applications, such as kernel principle component analysis [14] and spectral clustering [30]. This procedure is especially efficient because it uses approximations based on calculations on very small submatrices of the original matrix; therefore, it is useful when the graph is very large.

Here, we show how to apply the Nyström extension method in the case of the symmetric Laplacian, but it can be easily applied to other choices. Note that the eigenvectors of the matrix \( \tilde{W} = D^{-\frac{1}{2}} W D^{-\frac{1}{2}} \) are the same as those of the symmetric graph Laplacian, and their eigenvalues have a very simple relationship. Below, we formulate a method to calculate the eigenvectors and eigenvalues of \( \tilde{W} \) and thus of \( L_s \).

Let \( w \) be the weight function, \( \lambda \) be an eigenvalue of \( W \), and \( \phi \) its associated eigenvector. The Nyström method approximates the eigenvalue equation

\[
\int_{\Omega} w(y, x) \phi(x) dx = \lambda \phi(x)
\]

using a quadrature rule, a technique to find weights \( c_j(y) \) and a set of \( l \) interpolation points \( X = \{ x_j \} \) such that

\[
\sum_{j=1}^{l} c_j(y) \phi(x_j) = \int_{\Omega} w(y, x) \phi(x) dx + E(y)
\]
where $E(y)$ represents the error in the approximation. We use $c_j(y) = w(y, x_j)$ and choose the $l$ interpolation points randomly from the vertex set $V$. Denote the set of $l$ randomly chosen points by $X = \{x_i\}_{i=1}^l$ and its complement by $Y$. Partitioning $Z$ into $Z = X \cup Y$ and letting $\phi_k(x)$ be the the $k^{th}$ eigenvector of $W$ and $\lambda_k$ its associated eigenvalue, we obtain the system of equations

$$\sum_{x_j \in X} w(y_i, x_j) \phi_k(x_j) = \lambda_k \phi_k(y_i) \quad \forall y_i \in Y, \quad \forall k \in 1, ..., l.$$  

One cannot solve these equations directly because $\phi_k$ are unknown. However, the eigenvectors of $W$ are approximated using submatrices of $W$. Let $W_{XY}$ be defined as

$$W_{XY} = \begin{bmatrix} w(x_1, y_1) & \cdots & w(x_1, y_{n-l}) \\ \vdots & \ddots & \vdots \\ w(x_l, y_1) & \cdots & w(x_l, y_{n-l}) \end{bmatrix}$$

The matrices $W_{YX}, W_{XX}$ and $W_{YY}$ are defined similarly. The matrix $W$ is then

$$W = \begin{bmatrix} W_{XX} & W_{XY} \\ W_{YX} & W_{YY} \end{bmatrix}$$

We normalize the above weight matrix to calculate the eigenvalues and eigenvectors of $\hat{W}$. Here, $1_K$ is the $K$-dimensional unit vector. The matrices $d_X, d_Y, \hat{W}_{XX}$ and $\hat{W}_{XY}$ are defined as

$$d_X = W_{XX} 1_l + W_{XY} 1_{n-l},$$

$$d_Y = W_{YX} 1_l + (W_{YX} W_{XX}^{-1} W_{XY}) 1_{n-l},$$

$$\hat{W}_{XX} = W_{XX} / (s_X s_X^T),$$

$$\hat{W}_{XY} = W_{XY} / (s_X s_Y^T),$$

where $A / B$ denotes componentwise division between matrices $A$ and $B$, $v^T$ denotes the transpose of vector $v$, $s_X = \sqrt{d_X}$ and $s_Y = \sqrt{d_Y}$. It is shown in [2] that if $\hat{W}_{XX} = B_X D B_X^T$, and if $A$ and $\Gamma$ are matrices such that

$$A^T \Gamma A = \hat{W}_{XX} + \hat{W}_{XX}^{-\frac{1}{2}} \hat{W}_{XY} \hat{W}_{YX} \hat{W}_{XX}^{-\frac{1}{2}},$$

then the eigenvector matrix $V$ consisting of approximations of $L$ eigenvectors of $\hat{W}$, and thus of the symmetric graph Laplacian, is

$$V = \begin{bmatrix} B_X D^\frac{1}{2} B_X^T \Gamma^{-\frac{1}{2}} \\ \hat{W}_{YX} B_X D^{-\frac{1}{2}} B_X^T \Gamma^{-\frac{1}{2}} \end{bmatrix}$$

while $I - \Gamma$ contains the corresponding eigenvalues of the symmetric graph Laplacian in its diagonal entries.

One can see that the Nyström extension method is efficient because it approximates the eigenvalues and eigenvectors of an $n \times n$ matrix ($n$ is large) by calculations using matrices that are much smaller, having dimension no larger than $n \times l$, where $l$ is small. We also note that the method calculates $l$ eigenvalues and corresponding eigenvectors; in practice, only a very small portion (we use around 0.3%) of eigenvalues and eigenvectors of the graph Laplacian are needed to obtain an accurate solution.
References


